

RTEMS POSIX API User's Guide

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On-Line Applications Research Corporation

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Table of Contents

Preface	1
1 Process Creation and Execution Manager	3
1.1 Introduction	3
1.2 Background	3
1.3 Operations	3
1.4 Directives	3
1.4.1 fork - Create a Process	4
1.4.2 execl - Execute a File	5
1.4.3 execv - Execute a File	6
1.4.4 execl - Execute a File	7
1.4.5 execve - Execute a File	8
1.4.6 execlp - Execute a File	9
1.4.7 execvp - Execute a File	10
1.4.8 pthread_atfork - Register Fork Handlers	11
1.4.9 wait - Wait for Process Termination	12
1.4.10 waitpid - Wait for Process Termination	13
1.4.11 _exit - Terminate a Process	14
2 Signal Manager	15
2.1 Introduction	15
2.2 Background	15
2.2.1 Signals	15
2.2.2 Signal Delivery	16
2.3 Operations	16
2.3.1 Signal Set Management	16
2.3.2 Blocking Until Signal Generation	16
2.3.3 Sending a Signal	16
2.4 Directives	16
2.4.1 sigaddset - Add a Signal to a Signal Set	17
2.4.2 sigdelset - Delete a Signal from a Signal Set	18
2.4.3 sigfillset - Fill a Signal Set	19
2.4.4 sigismember - Is Signal a Member of a Signal Set	20
2.4.5 sigemptyset - Empty a Signal Set	21
2.4.6 sigaction - Examine and Change Signal Action	22
2.4.7 pthread_kill - Send a Signal to a Thread	23
2.4.8 sigprocmask - Examine and Change Process Blocked Signals	24
2.4.9 pthread_sigmask - Examine and Change Thread Blocked Signals	25
2.4.10 kill - Send a Signal to a Process	26
2.4.11 sigpending - Examine Pending Signals	27

2.4.12	sigsuspend - Wait for a Signal	28
2.4.13	pause - Suspend Process Execution	29
2.4.14	sigwait - Synchronously Accept a Signal	30
2.4.15	sigwaitinfo - Synchronously Accept a Signal	31
2.4.16	sigtimedwait - Synchronously Accept a Signal with Timeout	32
2.4.17	sigqueue - Queue a Signal to a Process	33
2.4.18	alarm - Schedule Alarm	34
2.4.19	ualarm - Schedule Alarm in Microseconds	35
3	Process Environment Manager	37
3.1	Introduction	37
3.2	Background	37
3.2.1	Users and Groups	37
3.2.2	User and Group Names	38
3.2.3	Environment Variables	38
3.3	Operations	38
3.3.1	Accessing User and Group Ids	38
3.3.2	Accessing Environment Variables	38
3.4	Directives	38
3.4.1	getpid - Get Process ID	39
3.4.2	getppid - Get Parent Process ID	40
3.4.3	getuid - Get User ID	41
3.4.4	geteuid - Get Effective User ID	42
3.4.5	getgid - Get Real Group ID	43
3.4.6	getegid - Get Effective Group ID	44
3.4.7	setuid - Set User ID	45
3.4.8	setgid - Set Group ID	46
3.4.9	getgroups - Get Supplementary Group IDs	47
3.4.10	getlogin - Get User Name	48
3.4.11	getlogin_r - Reentrant Get User Name	49
3.4.12	getpgrp - Get Process Group ID	50
3.4.13	setsid - Create Session and Set Process Group ID	51
3.4.14	setpgid - Set Process Group ID for Job Control	52
3.4.15	uname - Get System Name	53
3.4.16	times - Get process times	54
3.4.17	getenv - Get Environment Variables	55
3.4.18	setenv - Set Environment Variables	56
3.4.19	ctermid - Generate Terminal Pathname	57
3.4.20	ttyname - Determine Terminal Device Name	58
3.4.21	ttyname_r - Reentrant Determine Terminal Device Name	59
3.4.22	isatty - Determine if File Descriptor is Terminal	60
3.4.23	sysconf - Get Configurable System Variables	61

4	Files and Directories Manager	63
4.1	Introduction	63
4.2	Background	64
4.2.1	Path Name Evaluation	64
4.3	Operations	64
4.4	Directives	64
4.4.1	opendir - Open a Directory	65
4.4.2	readdir - Reads a directory	66
4.4.3	rewinddir - Resets the readdir() pointer	67
4.4.4	scandir - Scan a directory for matching entries	68
4.4.5	telldir - Return current location in directory stream	69
4.4.6	closedir - Ends directory read operation	70
4.4.7	chdir - Changes the current working directory	71
4.4.8	fchdir - Changes the current working directory	72
4.4.9	getcwd - Gets current working directory	73
4.4.10	open - Opens a file	74
4.4.11	creat - Create a new file or rewrite an existing one	76
4.4.12	umask - Sets a file creation mask.	77
4.4.13	link - Creates a link to a file	78
4.4.14	symlink - Creates a symbolic link to a file	79
4.4.15	readlink - Obtain the name of a symbolic link destination	80
4.4.16	mkdir - Makes a directory	81
4.4.17	mknfifo - Makes a FIFO special file	82
4.4.18	unlink - Removes a directory entry	83
4.4.19	rmdir - Delete a directory	84
4.4.20	rename - Renames a file	85
4.4.21	stat - Gets information about a file	86
4.4.22	fstat - Gets file status	87
4.4.23	lstat - Gets file status	88
4.4.24	access - Check permissions for a file	89
4.4.25	chmod - Changes file mode.	90
4.4.26	fchmod - Changes permissions of a file	91
4.4.27	getdents - Get directory entries	92
4.4.28	chown - Changes the owner and/or group of a file.	93
4.4.29	utime - Change access and/or modification times of an inode	94
4.4.30	ftruncate - truncate a file to a specified length	95
4.4.31	truncate - truncate a file to a specified length	96
4.4.32	pathconf - Gets configuration values for files	97
4.4.33	fpathconf - Gets configuration values for files	99
4.4.34	mknod - create a directory	101

5 Input and Output Primitives Manager 103

5.1	Introduction	103
5.2	Background	103
5.3	Operations	103
5.4	Directives	103
5.4.1	pipe - Create an Inter-Process Channel	104
5.4.2	dup - Duplicates an open file descriptor	105
5.4.3	dup2 - Duplicates an open file descriptor	106
5.4.4	close - Closes a file	107
5.4.5	read - Reads from a file	108
5.4.6	write - Writes to a file	110
5.4.7	fcntl - Manipulates an open file descriptor	111
5.4.8	lseek - Reposition read/write file offset	113
5.4.9	fsync - Synchronize file complete in-core state with that on disk	114
5.4.10	fdatasync - Synchronize file in-core data with that on disk	115
5.4.11	sync - Schedule file system updates	116
5.4.12	mount - Mount a file system	117
5.4.13	unmount - Unmount file systems	118
5.4.14	readv - Vectored read from a file	119
5.4.15	writew - Vectored write to a file	120
5.4.16	aio_read - Asynchronous Read	121
5.4.17	aio_write - Asynchronous Write	122
5.4.18	lio_listio - List Directed I/O	123
5.4.19	aio_error - Retrieve Error Status of Asynchronous I/O Operation	124
5.4.20	aio_return - Retrieve Return Status Asynchronous I/O Operation	125
5.4.21	aio_cancel - Cancel Asynchronous I/O Request	126
5.4.22	aio_suspend - Wait for Asynchronous I/O Request	127
5.4.23	aio_fsync - Asynchronous File Synchronization	128

6 Device- and Class- Specific Functions Manager 129

6.1	Introduction	129
6.2	Background	129
6.3	Operations	129
6.4	Directives	129
6.4.1	cfgetispeed - Reads terminal input baud rate	130
6.4.2	cfgetospeed - Reads terminal output baud rate	131
6.4.3	cfsetispeed - Sets terminal input baud rate	132
6.4.4	cfsetospeed - Sets terminal output baud rate	133
6.4.5	tcgetattr - Gets terminal attributes	134
6.4.6	tcsetattr - Set terminal attributes	135
6.4.7	tcsendbreak - Sends a break to a terminal	136
6.4.8	tcdrain - Waits for all output to be transmitted to the terminal	137

6.4.9	tcflush - Discards terminal data	138
6.4.10	tcflow - Suspends/restarts terminal output.....	139
6.4.11	tcgetpgrp - Gets foreground process group ID	140
6.4.12	tcsetpgrp - Sets foreground process group ID	141
7	Language-Specific Services for the C Programming Language Manager	143
7.1	Introduction	143
7.2	Background.....	143
7.3	Operations.....	143
7.4	Directives.....	143
7.4.1	setlocale - Set the Current Locale	144
7.4.2	fileno - Obtain File Descriptor Number for this File	145
7.4.3	fdopen - Associate Stream with File Descriptor	146
7.4.4	flockfile - Acquire Ownership of File Stream	147
7.4.5	ftrylockfile - Poll to Acquire Ownership of File Stream...	148
7.4.6	funlockfile - Release Ownership of File Stream	149
7.4.7	getc_unlocked - Get Character without Locking.....	150
7.4.8	getchar_unlocked - Get Character from stdin without Locking	151
7.4.9	putc_unlocked - Put Character without Locking	152
7.4.10	putchar_unlocked - Put Character to stdin without Locking	153
7.4.11	setjmp - Save Context for Non-Local Goto.....	154
7.4.12	longjmp - Non-Local Jump to a Saved Context	155
7.4.13	sigsetjmp - Save Context with Signal Status for Non-Local Goto	156
7.4.14	siglongjmp - Non-Local Jump with Signal Status to a Saved Context	157
7.4.15	tzset - Initialize Time Conversion Information	158
7.4.16	strtok_r - Reentrant Extract Token from String.....	159
7.4.17	asctime_r - Reentrant struct tm to ASCII Time Conversion	160
7.4.18	ctime_r - Reentrant time_t to ASCII Time Conversion ..	161
7.4.19	gmtime_r - Reentrant UTC Time Conversion	162
7.4.20	localtime_r - Reentrant Local Time Conversion	163
7.4.21	rand_r - Reentrant Random Number Generation.....	164
8	System Databases Manager	165
8.1	Introduction	165
8.2	Background.....	165
8.3	Operations.....	165
8.4	Directives.....	165
8.4.1	getgrgid - Get Group File Entry for ID	166
8.4.2	getgrgid_r - Reentrant Get Group File Entry	167
8.4.3	getgrnam - Get Group File Entry for Name.....	168
8.4.4	getgrnam_r - Reentrant Get Group File Entry for Name..	169

8.4.5	getpwuid - Get Password File Entry for UID	170
8.4.6	getpwuid_r - Reentrant Get Password File Entry for UID	171
8.4.7	getpwnam - Password File Entry for Name	172
8.4.8	getpwnam_r - Reentrant Get Password File Entry for Name	173
9	Semaphore Manager	175
9.1	Introduction	175
9.2	Background	175
9.2.1	Theory	175
9.2.2	"sem_t" Structure	175
9.2.3	Building a Semaphore Attribute Set	175
9.3	Operations	175
9.3.1	Using as a Binary Semaphore	176
9.4	Directives	176
9.4.1	sem_init - Initialize an unnamed semaphore	177
9.4.2	sem_destroy - Destroy an unnamed semaphore	178
9.4.3	sem_open - Open a named semaphore	179
9.4.4	sem_close - Close a named semaphore	180
9.4.5	sem_unlink - Unlink a semaphore	181
9.4.6	sem_wait - Wait on a Semaphore	182
9.4.7	sem_trywait - Non-blocking Wait on a Semaphore	183
9.4.8	sem_timedwait - Wait on a Semaphore for a Specified Time	184
9.4.9	sem_post - Unlock a Semaphore	185
9.4.10	sem_getvalue - Get the value of a semaphore	186
10	Mutex Manager	187
10.1	Introduction	187
10.2	Background	187
10.2.1	Mutex Attributes	187
10.2.2	PTHREAD_MUTEX_INITIALIZER	187
10.3	Operations	188
10.4	Services	188
10.4.1	pthread_mutexattr_init - Initialize a Mutex Attribute Set	189
10.4.2	pthread_mutexattr_destroy - Destroy a Mutex Attribute Set	190
10.4.3	pthread_mutexattr_setprotocol - Set the Blocking Protocol	191
10.4.4	pthread_mutexattr_getprotocol - Get the Blocking Protocol	192
10.4.5	pthread_mutexattr_setprioceiling - Set the Priority Ceiling	193
10.4.6	pthread_mutexattr_getprioceiling - Get the Priority Ceiling	194
10.4.7	pthread_mutexattr_setpshared - Set the Visibility	195

10.4.8	pthread_mutexattr_getpshared - Get the Visibility	196
10.4.9	pthread_mutex_init - Initialize a Mutex.....	197
10.4.10	pthread_mutex_destroy - Destroy a Mutex.....	198
10.4.11	pthread_mutex_lock - Lock a Mutex.....	199
10.4.12	pthread_mutex_trylock - Poll to Lock a Mutex	200
10.4.13	pthread_mutex_timedlock - Lock a Mutex with Timeout	201
10.4.14	pthread_mutex_unlock - Unlock a Mutex.....	202
10.4.15	pthread_mutex_setprioceiling - Dynamically Set the Priority Ceiling.....	203
10.4.16	pthread_mutex_getprioceiling - Get the Current Priority Ceiling.....	204
11	Condition Variable Manager.....	205
11.1	Introduction.....	205
11.2	Background.....	205
11.3	Operations	205
11.4	Directives.....	205
11.4.1	pthread_condattr_init - Initialize a Condition Variable Attribute Set.....	206
11.4.2	pthread_condattr_destroy - Destroy a Condition Variable Attribute Set.....	207
11.4.3	pthread_condattr_setpshared - Set Process Shared Attribute	208
11.4.4	pthread_condattr_getpshared - Get Process Shared Attribute	209
11.4.5	pthread_cond_init - Initialize a Condition Variable.....	210
11.4.6	pthread_cond_destroy - Destroy a Condition Variable...	211
11.4.7	pthread_cond_signal - Signal a Condition Variable.....	212
11.4.8	pthread_cond_broadcast - Broadcast a Condition Variable	213
11.4.9	pthread_cond_wait - Wait on a Condition Variable.....	214
11.4.10	pthread_cond_timedwait - Wait with Timeout a Condition Variable.....	215
12	Memory Management Manager.....	217
12.1	Introduction.....	217
12.2	Background.....	217
12.3	Operations	217
12.4	Directives.....	217
12.4.1	mlockall - Lock the Address Space of a Process.....	218
12.4.2	munlockall - Unlock the Address Space of a Process	219
12.4.3	mlock - Lock a Range of the Process Address Space	220
12.4.4	munlock - Unlock a Range of the Process Address Space	221
12.4.5	mmap - Map Process Addresses to a Memory Object ...	222
12.4.6	munmap - Unmap Previously Mapped Addresses	223
12.4.7	mprotect - Change Memory Protection	224

12.4.8	msync - Memory Object Synchronization	225
12.4.9	shm_open - Open a Shared Memory Object.....	226
12.4.10	shm_unlink - Remove a Shared Memory Object	227
13	Scheduler Manager	229
13.1	Introduction	229
13.2	Background.....	229
13.2.1	Priority.....	229
13.2.2	Scheduling Policies	229
13.3	Operations	229
13.4	Directives.....	230
13.4.1	sched_get_priority_min - Get Minimum Priority Value ..	231
13.4.2	sched_get_priority_max - Get Maximum Priority Value..	232
13.4.3	sched_rr_get_interval - Get Timeslicing Quantum	233
13.4.4	sched_yield - Yield the Processor	234
14	Clock Manager	235
14.1	Introduction	235
14.2	Background.....	235
14.3	Operations	235
14.4	Directives.....	235
14.4.1	clock_gettime - Obtain Time of Day	235
14.4.2	clock_settime - Set Time of Day	237
14.4.3	clock_getres - Get Clock Resolution	238
14.4.4	sleep - Delay Process Execution	239
14.4.5	usleep - Delay Process Execution in Microseconds	240
14.4.6	nanosleep - Delay with High Resolution	241
14.4.7	gettimeofday - Get the Time of Day	242
14.4.8	time - Get time in seconds	243
15	Timer Manager	245
15.1	Introduction	245
15.2	Background.....	245
15.3	Operations	245
15.4	System Calls.....	245
15.4.1	timer_create - Create a Per-Process Timer.....	246
15.4.2	timer_delete - Delete a Per-Process Timer	247
15.4.3	timer_settime - Set Next Timer Expiration	248
15.4.4	timer_gettime - Get Time Remaining on Timer	249
15.4.5	timer_getoverrun - Get Timer Overrun Count	250

16	Message Passing Manager	251
16.1	Introduction	251
16.2	Background	251
16.2.1	Theory	251
16.2.2	Messages	251
16.2.3	Message Queues	251
16.2.4	Building a Message Queue Attribute Set	252
16.2.5	Notification of a Message on the Queue	252
16.2.6	POSIX Interpretation Issues	253
16.3	Operations	253
16.3.1	Opening or Creating a Message Queue	253
16.3.2	Closing a Message Queue	253
16.3.3	Removing a Message Queue	253
16.3.4	Sending a Message to a Message Queue	253
16.3.5	Receiving a Message from a Message Queue	254
16.3.6	Notification of Receipt of a Message on an Empty Queue	254
16.3.7	Setting the Attributes of a Message Queue	254
16.3.8	Getting the Attributes of a Message Queue	255
16.4	Directives	255
16.4.1	mq_open - Open a Message Queue	256
16.4.2	mq_close - Close a Message Queue	258
16.4.3	mq_unlink - Remove a Message Queue	259
16.4.4	mq_send - Send a Message to a Message Queue	260
16.4.5	mq_receive - Receive a Message from a Message Queue	261
16.4.6	mq_notify - Notify Process that a Message is Available	262
16.4.7	mq_setattr - Set Message Queue Attributes	263
16.4.8	mq_getattr - Get Message Queue Attributes	264
17	Thread Manager	265
17.1	Introduction	265
17.2	Background	265
17.2.1	Thread Attributes	266
17.3	Operations	266
17.4	Services	266
17.4.1	pthread_attr_init - Initialize a Thread Attribute Set	267
17.4.2	pthread_attr_destroy - Destroy a Thread Attribute Set	268
17.4.3	pthread_attr_setdetachstate - Set Detach State	269
17.4.4	pthread_attr_getdetachstate - Get Detach State	270
17.4.5	pthread_attr_setstacksize - Set Thread Stack Size	271
17.4.6	pthread_attr_getstacksize - Get Thread Stack Size	272
17.4.7	pthread_attr_setstackaddr - Set Thread Stack Address	273
17.4.8	pthread_attr_getstackaddr - Get Thread Stack Address	274
17.4.9	pthread_attr_setscope - Set Thread Scheduling Scope	275
17.4.10	pthread_attr_getscope - Get Thread Scheduling Scope	276

17.4.11	pthread_attr_setinheritsched - Set Inherit Scheduler Flag	277
17.4.12	pthread_attr_getinheritsched - Get Inherit Scheduler Flag	278
17.4.13	pthread_attr_setschedpolicy - Set Scheduling Policy	279
17.4.14	pthread_attr_getschedpolicy - Get Scheduling Policy	280
17.4.15	pthread_attr_setschedparam - Set Scheduling Parameters	281
17.4.16	pthread_attr_getschedparam - Get Scheduling Parameters	282
17.4.17	pthread_create - Create a Thread	283
17.4.18	pthread_exit - Terminate the Current Thread	285
17.4.19	pthread_detach - Detach a Thread	286
17.4.20	pthread_join - Wait for Thread Termination	287
17.4.21	pthread_self - Get Thread ID	288
17.4.22	pthread_equal - Compare Thread IDs	289
17.4.23	pthread_once - Dynamic Package Initialization	290
17.4.24	pthread_setschedparam - Set Thread Scheduling Parameters	291
17.4.25	pthread_getschedparam - Get Thread Scheduling Parameters	292
18	Key Manager	293
18.1	Introduction	293
18.2	Background	293
18.3	Operations	293
18.4	Directives	293
18.4.1	pthread_key_create - Create Thread Specific Data Key	294
18.4.2	pthread_key_delete - Delete Thread Specific Data Key	295
18.4.3	pthread_setspecific - Set Thread Specific Key Value	296
18.4.4	pthread_getspecific - Get Thread Specific Key Value	297
19	Thread Cancellation Manager	299
19.1	Introduction	299
19.2	Background	299
19.3	Operations	299
19.4	Directives	299
19.4.1	pthread_cancel - Cancel Execution of a Thread	300
19.4.2	pthread_setcancelstate - Set Cancelability State	301
19.4.3	pthread_setcanceltype - Set Cancelability Type	302
19.4.4	pthread_testcancel - Create Cancellation Point	303
19.4.5	pthread_cleanup_push - Establish Cancellation Handler	304
19.4.6	pthread_cleanup_pop - Remove Cancellation Handler	305

20	Services Provided by C Library (libc)	307
20.1	Introduction	307
20.2	Standard Utility Functions (stdlib.h)	307
20.3	Character Type Macros and Functions (ctype.h)	308
20.4	Input and Output (stdio.h)	308
20.5	Strings and Memory (string.h)	309
20.6	Signal Handling (signal.h)	310
20.7	Time Functions (time.h)	310
20.8	Locale (locale.h)	310
20.9	Reentrant Versions of Functions	310
20.10	Miscellaneous Macros and Functions	312
20.11	Variable Argument Lists	313
20.12	Reentrant System Calls	313
21	Services Provided by the Math Library (libm)	
	315
21.1	Introduction	315
21.2	Standard Math Functions (math.h)	315
22	Status of Implementation	317
	Command and Variable Index	319
	Concept Index	323

Preface

This is the User's Guide for the POSIX API support provided in RTEMS.

The functionality described in this document is based on the following standards:

- POSIX 1003.1b-1993.
- POSIX 1003.1h/D3.
- Open Group Single UNIX Specification.

Much of the POSIX API standard is actually implemented in the Cygnus Newlib ANSI C Library. Please refer to documentation on Newlib for more information on the functionality it supplies.

This manual is still under construction and improvements are welcomed from users.

1 Process Creation and Execution Manager

1.1 Introduction

The process creation and execution manager provides the functionality associated with the creation and termination of processes.

The directives provided by the process creation and execution manager are:

- `fork` - Create a Process
- `execl` - Execute a File
- `execv` - Execute a File
- `execle` - Execute a File
- `execve` - Execute a File
- `execlp` - Execute a File
- `execvp` - Execute a File
- `pthread_atfork` - Register Fork Handlers
- `wait` - Wait for Process Termination
- `waitpid` - Wait for Process Termination
- `_exit` - Terminate a Process

1.2 Background

POSIX process functionality can not be completely supported by RTEMS. This is because RTEMS provides no memory protection and implements a *single process, multi-threaded execution model*. In this light, RTEMS provides none of the routines that are associated with the creation of new processes. However, since the entire RTEMS application (e.g. executable) is logically a single POSIX process, RTEMS is able to provide implementations of many operations on processes. The rule of thumb is that those routines provide a meaningful result. For example, `getpid()` returns the node number.

1.3 Operations

The only functionality method defined by this manager which is supported by RTEMS is the `_exit` service. The implementation of `_exit` shuts the application down and is equivalent to invoking either `exit` or `rtems_shutdown_executive`.

1.4 Directives

This section details the process creation and execution manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

1.4.1 fork - Create a Process

CALLING SEQUENCE:

```
#include <sys/types.h>
```

```
int fork( void );
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.2 execl - Execute a File

CALLING SEQUENCE:

```
int execl(  
    const char *path,  
    const char *arg,  
    ...  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.3 `execv` - Execute a File

CALLING SEQUENCE:

```
int execv(  
    const char *path,  
    char const *argv[],  
    ...  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.4 execl - Execute a File

CALLING SEQUENCE:

```
int execl(  
    const char *path,  
    const char *arg,  
    ...  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.5 `execve` - Execute a File

CALLING SEQUENCE:

```
int execve(  
    const char *path,  
    char *const argv[],  
    char *const envp[]  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.6 execlp - Execute a File

CALLING SEQUENCE:

```
int execlp(  
    const char *file,  
    const char *arg,  
    ...  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.7 execvp - Execute a File

CALLING SEQUENCE:

```
int execvp(  
    const char *file,  
    char *const argv[]  
    ...  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.8 pthread_atfork - Register Fork Handlers

CALLING SEQUENCE:

```
#include <sys/types.h>

int pthread_atfork(
    void (*prepare)(void),
    void (*parent)(void),
    void (*child)(void)
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.9 wait - Wait for Process Termination

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/wait.h>

int wait(
    int *stat_loc
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.10 waitpid - Wait for Process Termination

CALLING SEQUENCE:

```
int wait(  
    pid_t pid,  
    int *stat_loc,  
    int options  
);
```

STATUS CODES:

ENOSYS This routine is not supported by RTEMS.

DESCRIPTION:

This routine is not supported by RTEMS.

NOTES:

NONE

1.4.11 `_exit` - Terminate a Process

CALLING SEQUENCE:

```
void _exit(  
    int status  
);
```

STATUS CODES:

NONE

DESCRIPTION:

The `_exit()` function terminates the calling process.

NOTES:

In RTEMS, a process is equivalent to the entire application on a single processor. Invoking this service terminates the application.

2 Signal Manager

2.1 Introduction

The signal manager provides the functionality associated with the generation, delivery, and management of process-oriented signals.

The directives provided by the signal manager are:

- `sigaddset` - Add a Signal to a Signal Set
- `sigdelset` - Delete a Signal from a Signal Set
- `sigfillset` - Fill a Signal Set
- `sigismember` - Is Signal a Member of a Signal Set
- `sigemptyset` - Empty a Signal Set
- `sigaction` - Examine and Change Signal Action
- `pthread_kill` - Send a Signal to a Thread
- `sigprocmask` - Examine and Change Process Blocked Signals
- `pthread_sigmask` - Examine and Change Thread Blocked Signals
- `kill` - Send a Signal to a Process
- `sigpending` - Examine Pending Signals
- `sigsuspend` - Wait for a Signal
- `pause` - Suspend Process Execution
- `sigwait` - Synchronously Accept a Signal
- `sigwaitinfo` - Synchronously Accept a Signal
- `sigtimedwait` - Synchronously Accept a Signal with Timeout
- `sigqueue` - Queue a Signal to a Process
- `alarm` - Schedule Alarm
- `ualarm` - Schedule Alarm in Microseconds

2.2 Background

2.2.1 Signals

POSIX signals are an asynchronous event mechanism. Each process and thread has a set of signals associated with it. Individual signals may be enabled (e.g. unmasked) or blocked (e.g. ignored) on both a per-thread and process level. Signals which are enabled have a signal handler associated with them. When the signal is generated and conditions are met, then the signal handler is invoked in the proper process or thread context asynchronous relative to the logical thread of execution.

If a signal has been blocked when it is generated, then it is queued and kept pending until the thread or process unblocks the signal or explicitly checks for it. Traditional, non-real-time POSIX signals do not queue. Thus if a process or thread has blocked a particular signal, then multiple occurrences of that signal are recorded as a single occurrence of that signal.

One can check for the set of outstanding signals that have been blocked. Services are provided to check for outstanding process or thread directed signals.

2.2.2 Signal Delivery

Signals which are directed at a thread are delivered to the specified thread.

Signals which are directed at a process are delivered to a thread which is selected based on the following algorithm:

1. If the action for this signal is currently `SIG_IGN`, then the signal is simply ignored.
2. If the currently executing thread has the signal unblocked, then the signal is delivered to it.
3. If any threads are currently blocked waiting for this signal (`sigwait()`), then the signal is delivered to the highest priority thread waiting for this signal.
4. If any other threads are willing to accept delivery of the signal, then the signal is delivered to the highest priority thread of this set. In the event, multiple threads of the same priority are willing to accept this signal, then priority is given first to ready threads, then to threads blocked on calls which may be interrupted, and finally to threads blocked on non-interruptible calls.
5. In the event the signal still can not be delivered, then it is left pending. The first thread to unblock the signal (`sigprocmask()` or `pthread_sigprocmask()`) or to wait for this signal (`sigwait()`) will be the recipient of the signal.

2.3 Operations

2.3.1 Signal Set Management

Each process and each thread within that process has a set of individual signals and handlers associated with it. Services are provided to construct signal sets for the purposes of building signal sets – type `sigset_t` – that are used to provide arguments to the services that mask, unmask, and check on pending signals.

2.3.2 Blocking Until Signal Generation

A thread may block until receipt of a signal. The "sigwait" and "pause" families of services block until the requested signal is received or if using `sigtimedwait()` until the specified timeout period has elapsed.

2.3.3 Sending a Signal

This is accomplished via one of a number of services that sends a signal to either a process or thread. Signals may be directed at a process by the service `kill()` or at a thread by the service `pthread_kill()`

2.4 Directives

This section details the signal manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

2.4.1 sigaddset - Add a Signal to a Signal Set

CALLING SEQUENCE:

```
#include <signal.h>

int sigaddset(
    sigset_t *set,
    int      signo
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function adds the `signo` to the specified signal `set`.

NOTES:

NONE

2.4.2 sigdelset - Delete a Signal from a Signal Set

CALLING SEQUENCE:

```
#include <signal.h>

int sigdelset(
    sigset_t *set,
    int      signo
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function deletes the `signo` to the specified signal `set`.

NOTES:

NONE

2.4.3 sigfillset - Fill a Signal Set

CALLING SEQUENCE:

```
#include <signal.h>
```

```
int sigfillset(  
    sigset_t *set  
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function fills the specified signal **set** such that all signals are set.

NOTES:

NONE

2.4.4 sigismember - Is Signal a Member of a Signal Set

CALLING SEQUENCE:

```
#include <signal.h>

int sigismember(
    const sigset_t *set,
    int             signo
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function returns returns 1 if **signo** is a member of **set** and 0 otherwise.

NOTES:

NONE

2.4.5 sigemptyset - Empty a Signal Set

CALLING SEQUENCE:

```
#include <signal.h>

int sigemptyset(
    sigset_t *set
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function fills the specified signal **set** such that all signals are cleared.

NOTES:

NONE

2.4.6 sigaction - Examine and Change Signal Action

CALLING SEQUENCE:

```
#include <signal.h>

int sigaction(
    int                sig,
    const struct sigaction *act,
    struct sigaction   *oact
);
```

STATUS CODES:

EINVAL Invalid argument passed.

ENOTSUP Realtime Signals Extension option not supported.

DESCRIPTION:

This function is used to change the action taken by a process on receipt of the specific signal **sig**. The new action is specified by **act** and the previous action is returned via **oact**.

NOTES:

The signal number cannot be **SIGKILL**.

2.4.7 pthread_kill - Send a Signal to a Thread

CALLING SEQUENCE:

```
#include <signal.h>

int pthread_kill(
    pthread_t thread,
    int      sig
);
```

STATUS CODES:

ESRCH The thread indicated by the parameter `thread` is invalid.
EINVAL Invalid argument passed.

DESCRIPTION:

This function sends the specified signal `sig` to `thread`.

NOTES:

NONE

2.4.8 sigprocmask - Examine and Change Process Blocked Signals

CALLING SEQUENCE:

```
#include <signal.h>

int sigprocmask(
    int          how,
    const sigset_t *set,
    sigset_t     *oset
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function is used to alter the set of currently blocked signals on a process wide basis. A blocked signal will not be received by the process. The behavior of this function is dependent on the value of **how** which may be one of the following:

SIG_BLOCK	The set of blocked signals is set to the union of set and those signals currently blocked.
SIG_UNBLOCK	The signals specific in set are removed from the currently blocked set.
SIG_SETMASK	The set of currently blocked signals is set to set .

If **oset** is not **NULL**, then the set of blocked signals prior to this call is returned in **oset**.

NOTES:

It is not an error to unblock a signal which is not blocked.

2.4.9 pthread_sigmask - Examine and Change Thread Blocked Signals

CALLING SEQUENCE:

```
#include <signal.h>

int pthread_sigmask(
    int          how,
    const sigset_t *set,
    sigset_t     *oset
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

This function is used to alter the set of currently blocked signals for the calling thread. A blocked signal will not be received by the process. The behavior of this function is dependent on the value of `how` which may be one of the following:

SIG_BLOCK	The set of blocked signals is set to the union of <code>set</code> and those signals currently blocked.
SIG_UNBLOCK	The signals specific in <code>set</code> are removed from the currently blocked set.
SIG_SETMASK	The set of currently blocked signals is set to <code>set</code> .

If `oset` is not `NULL`, then the set of blocked signals prior to this call is returned in `oset`.

NOTES:

It is not an error to unblock a signal which is not blocked.

2.4.10 kill - Send a Signal to a Process

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <signal.h>

int kill(
    pid_t pid,
    int sig
);
```

STATUS CODES:

EINVAL	Invalid argument passed.
EPERM	Process does not have permission to send the signal to any receiving process.
ESRCH	The process indicated by the parameter <code>pid</code> is invalid.

DESCRIPTION:

This function sends the signal `sig` to the process `pid`.

NOTES:

NONE

2.4.11 sigpending - Examine Pending Signals

CALLING SEQUENCE:

```
#include <signal.h>

int sigpending(
    const sigset_t *set
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** to one of the following:

EFAULT Invalid address for set.

DESCRIPTION:

This function allows the caller to examine the set of currently pending signals. A pending signal is one which has been raised but is currently blocked. The set of pending signals is returned in **set**.

NOTES:

NONE

2.4.12 sigsuspend - Wait for a Signal

CALLING SEQUENCE:

```
#include <signal.h>

int sigsuspend(
    const sigset_t *sigmask
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** to one of the following:

EINTR Signal interrupted this function.

DESCRIPTION:

This function temporarily replaces the signal mask for the process with that specified by **sigmask** and blocks the calling thread until the signal is raised.

NOTES:

NONE

2.4.13 pause - Suspend Process Execution

CALLING SEQUENCE:

```
#include <signal.h>
```

```
int pause( void );
```

STATUS CODES:

On error, this routine returns -1 and sets `errno` to one of the following:

EINTR Signal interrupted this function.

DESCRIPTION:

This function causes the calling thread to be blocked until an unblocked signal is received.

NOTES:

NONE

2.4.14 sigwait - Synchronously Accept a Signal

CALLING SEQUENCE:

```
#include <signal.h>

int sigwait(
    const sigset_t *set,
    int             *sig
);
```

STATUS CODES:

EINVAL Invalid argument passed.
EINTR Signal interrupted this function.

DESCRIPTION:

This function selects a pending signal based on the set specified in **set**, atomically clears it from the set of pending signals, and returns the signal number for that signal in **sig**.

NOTES:

NONE

2.4.15 sigwaitinfo - Synchronously Accept a Signal

CALLING SEQUENCE:

```
#include <signal.h>

int sigwaitinfo(
    const sigset_t *set,
    siginfo_t      *info
);
```

STATUS CODES:

EINTR Signal interrupted this function.

DESCRIPTION:

This function selects a pending signal based on the set specified in **set**, atomically clears it from the set of pending signals, and returns information about that signal in **info**.

NOTES:

NONE

2.4.16 sigtimedwait - Synchronously Accept a Signal with Timeout

CALLING SEQUENCE:

```
#include <signal.h>

int sigtimedwait(
    const sigset_t      *set,
    siginfo_t           *info,
    const struct timespec *timeout
);
```

STATUS CODES:

EAGAIN	Timed out while waiting for the specified signal set.
EINVAL	Nanoseconds field of the timeout argument is invalid.
EINTR	Signal interrupted this function.

DESCRIPTION:

This function selects a pending signal based on the set specified in **set**, atomically clears it from the set of pending signals, and returns information about that signal in **info**. The calling thread will block up to **timeout** waiting for the signal to arrive.

NOTES:

If **timeout** is NULL, then the calling thread will wait forever for the specified signal set.

2.4.17 sigqueue - Queue a Signal to a Process

CALLING SEQUENCE:

```
#include <signal.h>

int sigqueue(
    pid_t          pid,
    int            signo,
    const union sigval value
);
```

STATUS CODES:

EAGAIN	No resources available to queue the signal. The process has already queued SIGQUEUE_MAX signals that are still pending at the receiver or the systemwide resource limit has been exceeded.
EINVAL	The value of the signo argument is an invalid or unsupported signal number.
EPERM	The process does not have the appropriate privilege to send the signal to the receiving process.
ESRCH	The process pid does not exist.

DESCRIPTION:

This function sends the signal specified by **signo** to the process **pid**

NOTES:

NONE

2.4.18 alarm - Schedule Alarm

CALLING SEQUENCE:

```
#include <unistd.h>

unsigned int alarm(
    unsigned int seconds
);
```

STATUS CODES:

This call always succeeds.

If there was a previous `alarm()` request with time remaining, then this routine returns the number of seconds until that outstanding alarm would have fired. If no previous `alarm()` request was outstanding, then zero is returned.

DESCRIPTION:

The `alarm()` service causes the `SIGALRM` signal to be generated after the number of seconds specified by `seconds` has elapsed.

NOTES:

Alarm requests do not queue. If `alarm` is called while a previous request is outstanding, the call will result in rescheduling the time at which the `SIGALRM` signal will be generated.

If the notification signal, `SIGALRM`, is not caught or ignored, the calling process is terminated.

2.4.19 ualarm - Schedule Alarm in Microseconds

CALLING SEQUENCE:

```
#include <unistd.h>

useconds_t ualarm(
    useconds_t useconds,
    useconds_t interval
);
```

STATUS CODES:

This call always succeeds.

If there was a previous `ualarm()` request with time remaining, then this routine returns the number of seconds until that outstanding alarm would have fired. If no previous `alarm()` request was outstanding, then zero is returned.

DESCRIPTION:

The `ualarm()` service causes the `SIGALRM` signal to be generated after the number of microseconds specified by `useconds` has elapsed.

When `interval` is non-zero, repeated timeout notification occurs with a period in microseconds specified by `interval`.

NOTES:

Alarm requests do not queue. If `alarm` is called while a previous request is outstanding, the call will result in rescheduling the time at which the `SIGALRM` signal will be generated.

If the notification signal, `SIGALRM`, is not caught or ignored, the calling process is terminated.

3 Process Environment Manager

3.1 Introduction

The process environment manager is responsible for providing the functions related to user and group Id management.

The directives provided by the process environment manager are:

- `getpid` - Get Process ID
- `getppid` - Get Parent Process ID
- `getuid` - Get User ID
- `geteuid` - Get Effective User ID
- `getgid` - Get Real Group ID
- `getegid` - Get Effective Group ID
- `setuid` - Set User ID
- `setgid` - Set Group ID
- `getgroups` - Get Supplementary Group IDs
- `getlogin` - Get User Name
- `getlogin_r` - Reentrant Get User Name
- `getpgrp` - Get Process Group ID
- `setsid` - Create Session and Set Process Group ID
- `setpgid` - Set Process Group ID for Job Control
- `uname` - Get System Name
- `times` - Get Process Times
- `getenv` - Get Environment Variables
- `setenv` - Set Environment Variables
- `ctermid` - Generate Terminal Pathname
- `ttyname` - Determine Terminal Device Name
- `ttyname_r` - Reentrant Determine Terminal Device Name
- `isatty` - Determine if File Descriptor is Terminal
- `sysconf` - Get Configurable System Variables

3.2 Background

3.2.1 Users and Groups

RTEMS provides a single process, multi-threaded execution environment. In this light, the notion of user and group is somewhat without meaning. But RTEMS does provide services to provide a synthetic version of user and group. By default, a single user and group is associated with the application. Thus unless special actions are taken, every thread in the application shares the same user and group Id. The initial rationale for providing user and group Id functionality in RTEMS was for the filesystem infrastructure to implement file

permission checks. The effective user/group Id capability has since been used to implement permissions checking by the `ftpd` server.

In addition to the "real" user and group Ids, a process may have an effective user/group Id. This allows a process to function using a more limited permission set for certain operations.

3.2.2 User and Group Names

POSIX considers user and group Ids to be a unique integer that may be associated with a name. This is usually accomplished via a file named `/etc/passwd` for user Id mapping and `/etc/groups` for group Id mapping. Again, although RTEMS is effectively a single process and thus single user system, it provides limited support for user and group names. When configured with an appropriate filesystem, RTEMS will access the appropriate files to map user and group Ids to names.

If these files do not exist, then RTEMS will synthesize a minimal version so this family of services return without error. It is important to remember that a design goal of the RTEMS POSIX services is to provide useable and meaningful results even though a full process model is not available.

3.2.3 Environment Variables

POSIX allows for variables in the run-time environment. These are name/value pairs that make be dynamically set and obtained by programs. In a full POSIX environment with command line shell and multiple processes, environment variables may be set in one process – such as the shell – and inherited by child processes. In RTEMS, there is only one process and thus only one set of environment variables across all processes.

3.3 Operations

3.3.1 Accessing User and Group Ids

The user Id associated with the current thread may be obtain using the `getuid()` service. Similarly, the group Id may be obtained using the `getgid()` service.

3.3.2 Accessing Environment Variables

The value associated with an environment variable may be obtained using the `getenv()` service and set using the `putenv()` service.

3.4 Directives

This section details the process environment manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

3.4.1 getpid - Get Process ID

CALLING SEQUENCE:

```
int getpid( void );
```

STATUS CODES:

The process Id is returned.

DESCRIPTION:

This service returns the process Id.

NOTES:

NONE

3.4.2 getppid - Get Parent Process ID

CALLING SEQUENCE:

```
int getppid( void );
```

STATUS CODES:

The parent process Id is returned.

DESCRIPTION:

This service returns the parent process Id.

NOTES:

NONE

3.4.3 getuid - Get User ID

CALLING SEQUENCE:

```
int getuid( void );
```

STATUS CODES:

The effective user Id is returned.

DESCRIPTION:

This service returns the effective user Id.

NOTES:

NONE

3.4.4 geteuid - Get Effective User ID

CALLING SEQUENCE:

```
int geteuid( void );
```

STATUS CODES:

The effective group Id is returned.

DESCRIPTION:

This service returns the effective group Id.

NOTES:

NONE

3.4.5 getgid - Get Real Group ID

CALLING SEQUENCE:

```
int getgid( void );
```

STATUS CODES:

The group Id is returned.

DESCRIPTION:

This service returns the group Id.

NOTES:

NONE

3.4.6 getegid - Get Effective Group ID

CALLING SEQUENCE:

```
int getegid( void );
```

STATUS CODES:

The effective group Id is returned.

DESCRIPTION:

This service returns the effective group Id.

NOTES:

NONE

3.4.7 setuid - Set User ID

CALLING SEQUENCE:

```
int setuid(  
    uid_t uid  
);
```

STATUS CODES:

This service returns 0.

DESCRIPTION:

This service sets the user Id to uid.

NOTES:

NONE

3.4.8 setgid - Set Group ID

CALLING SEQUENCE:

```
int setgid(  
    gid_t gid  
);
```

STATUS CODES:

This service returns 0.

DESCRIPTION:

This service sets the group Id to `gid`.

NOTES:

NONE

3.4.9 getgroups - Get Supplementary Group IDs

CALLING SEQUENCE:

```
int getgroups(  
    int    gidsetsize,  
    gid_t  grouplist[]  
);
```

STATUS CODES:

NA

DESCRIPTION:

This service is not implemented as RTEMS has no notion of supplemental groups.

NOTES:

If supported, this routine would only be allowed for the super-user.

3.4.10 `getlogin` - Get User Name

CALLING SEQUENCE:

```
char *getlogin( void );
```

STATUS CODES:

Returns a pointer to a string containing the name of the current user.

DESCRIPTION:

This routine returns the name of the current user.

NOTES:

This routine is not reentrant and subsequent calls to `getlogin()` will overwrite the same buffer.

3.4.11 getlogin_r - Reentrant Get User Name

CALLING SEQUENCE:

```
int getlogin_r(  
    char    *name,  
    size_t  namesize  
);
```

STATUS CODES:

EINVAL The arguments were invalid.

DESCRIPTION:

This is a reentrant version of the `getlogin()` service. The caller specified their own buffer, `name`, as well as the length of this buffer, `namesize`.

NOTES:

NONE

3.4.12 getpgrp - Get Process Group ID

CALLING SEQUENCE:

```
pid_t getpgrp( void );
```

STATUS CODES:

The process group Id is returned.

DESCRIPTION:

This service returns the current process group Id.

NOTES:

This routine is implemented in a somewhat meaningful way for RTEMS but is truly not functional.

3.4.13 setsid - Create Session and Set Process Group ID

CALLING SEQUENCE:

```
pid_t setsid( void );
```

STATUS CODES:

EPERM The application does not have permission to create a process group.

DESCRIPTION:

This routine always returns **EPERM** as RTEMS has no way to create new processes and thus no way to create a new process group.

NOTES:

NONE

3.4.14 setpgid - Set Process Group ID for Job Control

CALLING SEQUENCE:

```
int setpgid(
    pid_t pid,
    pid_t pgid
);
```

STATUS CODES:

ENOSYS The routine is not implemented.

DESCRIPTION:

This service is not implemented for RTEMS as process groups are not supported.

NOTES:

NONE

3.4.15 uname - Get System Name

CALLING SEQUENCE:

```
int uname(  
    struct utsname *name  
);
```

STATUS CODES:

EPERM The provided structure pointer is invalid.

DESCRIPTION:

This service returns system information to the caller. It does this by filling in the `struct utsname` format structure for the caller.

NOTES:

The information provided includes the operating system (RTEMS in all configurations), the node number, the release as the RTEMS version, and the CPU family and model. The CPU model name will indicate the multilib executive variant being used.

3.4.16 times - Get process times

CALLING SEQUENCE:

```
#include <sys/time.h>

clock_t times(
    struct tms *ptms
);
```

STATUS CODES:

This routine returns the number of clock ticks that have elapsed since the system was initialized (e.g. the application was started).

DESCRIPTION:

`times` stores the current process times in `ptms`. The format of `struct tms` is as defined in `<sys/times.h>`. RTEMS fills in the field `tms_utime` with the number of ticks that the calling thread has executed and the field `tms_stime` with the number of clock ticks since system boot (also returned). All other fields in the `ptms` are left zero.

NOTES:

RTEMS has no way to distinguish between user and system time so this routine returns the most meaningful information possible.

3.4.17 `getenv` - Get Environment Variables

CALLING SEQUENCE:

```
char *getenv(  
    const char *name  
);
```

STATUS CODES:

NULL when no match
pointer to value when successful

DESCRIPTION:

This service searches the set of environment variables for a string that matches the specified **name**. If found, it returns the associated value.

NOTES:

The environment list consists of name value pairs that are of the form *name = value*.

3.4.18 setenv - Set Environment Variables

CALLING SEQUENCE:

```
int setenv(  
    const char *name,  
    const char *value,  
    int overwrite  
);
```

STATUS CODES:

Returns 0 if successful and -1 otherwise.

DESCRIPTION:

This service adds the variable `name` to the environment with `value`. If `name` is not already exist, then it is created. If `name` exists and `overwrite` is zero, then the previous value is not overwritten.

NOTES:

NONE

3.4.19 ctermid - Generate Terminal Pathname

CALLING SEQUENCE:

```
char *ctermid(  
    char *s  
);
```

STATUS CODES:

Returns a pointer to a string indicating the pathname for the controlling terminal.

DESCRIPTION:

This service returns the name of the terminal device associated with this process. If **s** is NULL, then a pointer to a static buffer is returned. Otherwise, **s** is assumed to have a buffer of sufficient size to contain the name of the controlling terminal.

NOTES:

By default on RTEMS systems, the controlling terminal is `/dev/console`. Again this implementation is of limited meaning, but it provides true and useful results which should be sufficient to ease porting applications from a full POSIX implementation to the reduced profile supported by RTEMS.

3.4.20 ttyname - Determine Terminal Device Name

CALLING SEQUENCE:

```
char *ttyname(  
    int fd  
);
```

STATUS CODES:

Pointer to a string containing the terminal device name or NULL is returned on any error.

DESCRIPTION:

This service returns a pointer to the pathname of the terminal device that is open on the file descriptor `fd`. If `fd` is not a valid descriptor for a terminal device, then NULL is returned.

NOTES:

This routine uses a static buffer.

3.4.21 `ttynam_r` - Reentrant Determine Terminal Device Name

CALLING SEQUENCE:

```
int ttynam_r(  
    int fd,  
    char *name,  
    int namesize  
);
```

STATUS CODES:

This routine returns -1 and sets `errno` as follows:

EBADF	If not a valid descriptor for a terminal device.
EINVAL	If <code>name</code> is NULL or <code>namesize</code> are insufficient.

DESCRIPTION:

This service the pathname of the terminal device that is open on the file descriptor `fd`.

NOTES:

NONE

3.4.22 isatty - Determine if File Descriptor is Terminal

CALLING SEQUENCE:

```
int isatty(  
    int fd  
);
```

STATUS CODES:

Returns 1 if `fd` is a terminal device and 0 otherwise.

DESCRIPTION:

This service returns 1 if `fd` is an open file descriptor connected to a terminal and 0 otherwise.

NOTES:

3.4.23 sysconf - Get Configurable System Variables

CALLING SEQUENCE:

```
long sysconf(  
    int name  
);
```

STATUS CODES:

The value returned is the actual value of the system resource. If the requested configuration name is a feature flag, then 1 is returned if the available and 0 if it is not. On any other error condition, -1 is returned.

DESCRIPTION:

This service is the mechanism by which an application determines values for system limits or options at runtime.

NOTES:

Much of the information that may be obtained via `sysconf` has equivalent macros in `<unistd.h>`. However, those macros reflect conservative limits which may have been altered by application configuration.

4 Files and Directories Manager

4.1 Introduction

The files and directories manager is ...

The directives provided by the files and directories manager are:

- `opendir` - Open a Directory
- `readdir` - Reads a directory
- `rewinddir` - Resets the `readdir()` pointer
- `scandir` - Scan a directory for matching entries
- `telldir` - Return current location in directory stream
- `closedir` - Ends directory read operation
- `getdents` - Get directory entries
- `chdir` - Changes the current working directory
- `fchdir` - Changes the current working directory
- `getcwd` - Gets current working directory
- `open` - Opens a file
- `creat` - Create a new file or rewrite an existing one
- `umask` - Sets a file creation mask
- `link` - Creates a link to a file
- `symlink` - Creates a symbolic link to a file
- `readlink` - Obtain the name of the link destination
- `mkdir` - Makes a directory
- `mkfifo` - Makes a FIFO special file
- `unlink` - Removes a directory entry
- `rmdir` - Delete a directory
- `rename` - Renames a file
- `stat` - Gets information about a file.
- `fstat` - Gets file status
- `lstat` - Gets file status
- `access` - Check permissions for a file.
- `chmod` - Changes file mode
- `fchmod` - Changes permissions of a file
- `chown` - Changes the owner and/ or group of a file
- `utime` - Change access and/or modification times of an inode
- `ftruncate` - Truncate a file to a specified length
- `truncate` - Truncate a file to a specified length
- `pathconf` - Gets configuration values for files
- `fpathconf` - Get configuration values for files
- `mknod` - Create a directory

4.2 Background

4.2.1 Path Name Evaluation

A pathname is a string that consists of no more than `PATH_MAX` bytes, including the terminating null character. A pathname has an optional beginning slash, followed by zero or more filenames separated by slashes. If the pathname refers to a directory, it may also have one or more trailing slashes. Multiple successive slashes are considered to be the same as one slash.

POSIX allows a pathname that begins with precisely two successive slashes to be interpreted in an implementation-defined manner. RTEMS does not currently recognize this as a special condition. Any number of successive slashes is treated the same as a single slash. POSIX requires that an implementation treat more than two leading slashes as a single slash.

4.3 Operations

There is currently no text in this section.

4.4 Directives

This section details the files and directories manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

4.4.1 opendir - Open a Directory

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <dirent.h>

int opendir(
    const char *dirname
);
```

STATUS CODES:

EACCES	Search permission was denied on a component of the path prefix of <code>dirname</code> , or read permission is denied
EMFILE	Too many file descriptors in use by process
ENFILE	Too many files are currently open in the system.
ENOENT	Directory does not exist, or <code>name</code> is an empty string.
ENOMEM	Insufficient memory to complete the operation.
ENOTDIR	<code>name</code> is not a directory.

DESCRIPTION:

This routine opens a directory stream corresponding to the directory specified by the `dirname` argument. The directory stream is positioned at the first entry.

NOTES:

The routine is implemented in Cygnus newlib.

4.4.2 readdir - Reads a directory

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <dirent.h>

int readdir(
    DIR *dirp
);
```

STATUS CODES:

EBADF Invalid file descriptor

DESCRIPTION:

The `readdir()` function returns a pointer to a structure `dirent` representing the next directory entry from the directory stream pointed to by `dirp`. On end-of-file, `NULL` is returned.

The `readdir()` function may (or may not) return entries for `.` or `..`. Your program should tolerate reading dot and dot-dot but not require them.

The data pointed to by `readdir()` may be overwritten by another call to `readdir()` for the same directory stream. It will not be overwritten by a call for another directory.

NOTES:

If `ptr` is not a pointer returned by `malloc()`, `calloc()`, or `realloc()` or has been deallocated with `free()` or `realloc()`, the results are not portable and are probably disastrous.

The routine is implemented in Cygnus newlib.

4.4.3 rewinddir - Resets the readdir() pointer

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <dirent.h>

void rewinddir(
    DIR *dirp
);
```

STATUS CODES:

No value is returned.

DESCRIPTION:

The `rewinddir()` function resets the position associated with the directory stream pointed to by `dirp`. It also causes the directory stream to refer to the current state of the directory.

NOTES:

NONE

If `dirp` is not a pointer by `opendir()`, the results are undefined.

The routine is implemented in Cygnus newlib.

4.4.4 scandir - Scan a directory for matching entries

CALLING SEQUENCE:

```
#include <dirent.h>

int scandir(
    const char      *dir,
    struct dirent  ***namelist,
    int  (*select)(const struct dirent *),
    int  (*compar)(const struct dirent **, const struct dirent **)
);
```

STATUS CODES:

ENOMEM Insufficient memory to complete the operation.

DESCRIPTION:

The `scandir()` function scans the directory `dir`, calling `select()` on each directory entry. Entries for which `select()` returns non-zero are stored in strings allocated via `malloc()`, sorted using `qsort()` with the comparison function `compar()`, and collected in array `namelist` which is allocated via `malloc()`. If `select` is `NULL`, all entries are selected.

NOTES:

The routine is implemented in Cygnus newlib.

4.4.5 telldir - Return current location in directory stream

CALLING SEQUENCE:

```
#include <dirent.h>
```

```
off_t telldir(  
    DIR *dir  
);
```

STATUS CODES:

EBADF Invalid directory stream descriptor `dir`.

DESCRIPTION:

The `telldir()` function returns the current location associated with the directory stream `dir`.

NOTES:

The routine is implemented in Cygnus newlib.

4.4.6 closedir - Ends directory read operation

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <dirent.h>

int closedir(
    DIR *dirp
);
```

STATUS CODES:

EBADF Invalid file descriptor

DESCRIPTION:

The directory stream associated with `dirp` is closed. The value in `dirp` may not be usable after a call to `closedir()`.

NOTES:

NONE

The argument to `closedir()` must be a pointer returned by `opendir()`. If it is not, the results are not portable and most likely unpleasant.

The routine is implemented in Cygnus newlib.

4.4.7 chdir - Changes the current working directory

CALLING SEQUENCE:

```
#include <unistd.h>

int chdir(
    const char *path
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** to one of the following:

EACCES	Search permission is denied for a directory in a file's path prefix.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified pathname was not a directory when directory was expected.

DESCRIPTION:

The `chdir()` function causes the directory named by `path` to become the current working directory; that is, the starting point for searches of pathnames not beginning with a slash.

If `chdir()` detects an error, the current working directory is not changed.

NOTES:

NONE

4.4.8 fchdir - Changes the current working directory

CALLING SEQUENCE:

```
#include <unistd.h>

int fchdir(
    int fd
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** to one of the following:

EACCES	Search permission is denied for a directory in a file's path prefix.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified pathname was not a directory when directory was expected.

DESCRIPTION:

The `fchdir()` function causes the directory named by `fd` to become the current working directory; that is, the starting point for searches of pathnames not beginning with a slash.

If `fchdir()` detects an error, the current working directory is not changed.

NOTES:

NONE

4.4.9 `getcwd` - Gets current working directory

CALLING SEQUENCE:

```
#include <unistd.h>
```

```
int getcwd( void );
```

STATUS CODES:

EINVAL	Invalid argument
ERANGE	Result is too large
EACCES	Search permission is denied for a directory in a file's path prefix.

DESCRIPTION:

The `getcwd()` function copies the absolute pathname of the current working directory to the character array pointed to by `buf`. The `size` argument is the number of bytes available in `buf`.

NOTES:

There is no way to determine the maximum string length that `getcwd()` may need to return. Applications should tolerate getting **ERANGE** and allocate a larger buffer.

It is possible for `getcwd()` to return **EACCES** if, say, `login` puts the process into a directory without read access.

The 1988 standard uses `int` instead of `size_t` for the second parameter.

4.4.10 open - Opens a file

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>

int open(
    const char *path,
    int         oflag,
    mode_t      mode
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix.
EEXIST	The named file already exists.
EINTR	Function was interrupted by a signal.
EISDIR	Attempt to open a directory for writing or to rename a file to be a directory.
EMFILE	Too many file descriptors are in use by this process.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENFILE	Too many files are currently open in the system.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
ENXIO	No such device. This error may also occur when a device is not ready, for example, a tape drive is off-line.
EROFS	Read-only file system.

DESCRIPTION:

The `open` function establishes a connection between a file and a file descriptor. The file descriptor is a small integer that is used by I/O functions to reference the file. The `path` argument points to the pathname for the file.

The `oflag` argument is the bitwise inclusive OR of the values of symbolic constants. The programmer must specify exactly one of the following three symbols:

O_RDONLY	Open for reading only.
O_WRONLY	Open for writing only.
O_RDWR	Open for reading and writing.

Any combination of the following symbols may also be used.

O_APPEND	Set the file offset to the end-of-file prior to each write.
O_CREAT	If the file does not exist, allow it to be created. This flag indicates that the <code>mode</code> argument is present in the call to <code>open</code> .
O_EXCL	This flag may be used only if <code>O_CREAT</code> is also set. It causes the call to <code>open</code> to fail if the file already exists.
O_NOCTTY	If <code>path</code> identifies a terminal, this flag prevents that terminal from becoming the controlling terminal for this process. See Chapter 8 for a description of terminal I/O.
O_NONBLOCK	Do not wait for the device or file to be ready or available. After the file is open, the <code>read</code> and <code>write</code> calls return immediately. If the process would be delayed in the read or write operation, <code>-1</code> is returned and <code>errno</code> is set to <code>EAGAIN</code> instead of blocking the caller.
O_TRUNC	This flag should be used only on ordinary files opened for writing. It causes the file to be truncated to zero length.

Upon successful completion, `open` returns a non-negative file descriptor.

NOTES:

NONE

4.4.11 creat - Create a new file or rewrite an existing one

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>

int creat(
    const char *path,
    mode_t      mode
);
```

STATUS CODES:

EEXIST	<code>path</code> already exists and <code>O_CREAT</code> and <code>O_EXCL</code> were used.
EISDIR	<code>path</code> refers to a directory and the access requested involved writing
ETXTBSY	<code>path</code> refers to an executable image which is currently being executed and write access was requested
EFAULT	<code>path</code> points outside your accessible address space
EACCES	The requested access to the file is not allowed, or one of the directories in <code>path</code> did not allow search (execute) permission.
ENAMETOOLONG	<code>path</code> was too long.
ENOENT	A directory component in <code>path</code> does not exist or is a dangling symbolic link.
ENOTDIR	A component used as a directory in <code>path</code> is not, in fact, a directory.
EMFILE	The process already has the maximum number of files open.
ENFILE	The limit on the total number of files open on the system has been reached.
ENOMEM	Insufficient kernel memory was available.
EROFS	<code>path</code> refers to a file on a read-only filesystem and write access was requested

DESCRIPTION:

`creat` attempts to create a file and return a file descriptor for use in read, write, etc.

NOTES:

NONE

The routine is implemented in Cygnus newlib.

4.4.12 umask - Sets a file creation mask.

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

mode_t umask(
    mode_t cmask
);
```

STATUS CODES:

DESCRIPTION:

The `umask()` function sets the process file creation mask to `cmask`. The file creation mask is used during `open()`, `creat()`, `mkdir()`, `mkfifo()` calls to turn off permission bits in the `mode` argument. Bit positions that are set in `cmask` are cleared in the mode of the created file.

NOTES:

NONE

The `cmask` argument should have only permission bits set. All other bits should be zero.

In a system which supports multiple processes, the file creation mask is inherited across `fork()` and `exec()` calls. This makes it possible to alter the default permission bits of created files. RTEMS does not support multiple processes so this behavior is not possible.

4.4.13 link - Creates a link to a file

CALLING SEQUENCE:

```
#include <unistd.h>

int link(
    const char *existing,
    const char *new
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EEXIST	The named file already exists.
EMLINK	The number of links would exceed LINK_MAX.
ENAMETOOLONG	Length of a filename string exceeds PATH_MAX and _POSIX_NO_TRUNC is in effect.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EPERM	Operation is not permitted. Process does not have the appropriate privileges or permissions to perform the requested operations.
EROFS	Read-only file system.
EXDEV	Attempt to link a file to another file system.

DESCRIPTION:

The `link()` function atomically creates a new link for an existing file and increments the link count for the file.

If the `link()` function fails, no directories are modified.

The `existing` argument should not be a directory.

The caller may (or may not) need permission to access the existing file.

NOTES:

NONE

4.4.14 symlink - Creates a symbolic link to a file

CALLING SEQUENCE:

```
#include <unistd.h>

int symlink(
    const char *topath,
    const char *frompath
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EEXIST	The named file already exists.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EPERM	Operation is not permitted. Process does not have the appropriate privileges or permissions to perform the requested operations.
EROFS	Read-only file system.

DESCRIPTION:

The `symlink()` function creates a symbolic link from the `frompath` to the `topath`. The symbolic link will be interpreted at run-time.

If the `symlink()` function fails, no directories are modified.

The caller may (or may not) need permission to access the existing file.

NOTES:

NONE

4.4.15 readlink - Obtain the name of a symbolic link destination

CALLING SEQUENCE:

```
#include <unistd.h>

int readlink(
    const char *path,
    char      *buf,
    size_t    bufsize
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the prefix pathname was not a directory when a directory was expected.
ELOOP	Too many symbolic links were encountered in the pathname.
EINVAL	The pathname does not refer to a symbolic link
EFAULT	An invalid pointer was passed into the <code>readlink()</code> routine.

DESCRIPTION:

The `readlink()` function places the symbolic link destination into `buf` argument and returns the number of characters copied.

If the symbolic link destination is longer than `bufsize` characters the name will be truncated.

NOTES:

NONE

4.4.16 mkdir - Makes a directory

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int mkdir(
    const char *path,
    mode_t     mode
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EEXIST	The name file already exist.
EMLINK	The number of links would exceed LINK_MAX
ENAMETOOLONG	Length of a filename string exceeds PATH_MAX and _POSIX_NO_TRUNC is in effect.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EROFS	Read-only file system.

DESCRIPTION:

The `mkdir()` function creates a new diectory named `path`. The permission bits (modified by the file creation mask) are set from `mode`. The owner and group IDs for the directory are set from the effective user ID and group ID.

The new directory may (or may not) contain entries for `..` and `..` but is otherwise empty.

NOTES:

NONE

4.4.17 mkfifo - Makes a FIFO special file

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>
```

```
int mkfifo(
    const char *path,
    mode_t      mode
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EEXIST	The named file already exists.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified <code>path</code> was not a directory when a directory was expected.
EROFS	Read-only file system.

DESCRIPTION:

The `mkfifo()` function creates a new FIFO special file named `path`. The permission bits (modified by the file creation mask) are set from `mode`. The owner and group IDs for the FIFO are set from the effective user ID and group ID.

NOTES:

NONE

4.4.18 unlink - Removes a directory entry

CALLING SEQUENCE:

```
#include <unistd.h>

int unlink(
    const char path
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EBUSY	The directory is in use.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified <code>path</code> was not a directory when a directory was expected.
EPERM	Operation is not permitted. Process does not have the appropriate privileges or permissions to perform the requested operations.
EROFS	Read-only file system.

DESCRIPTION:

The `unlink` function removes the link named by `path` and decrements the link count of the file referenced by the link. When the link count goes to zero and no process has the file open, the space occupied by the file is freed and the file is no longer accessible.

NOTES:

NONE

4.4.19 rmdir - Delete a directory

CALLING SEQUENCE:

```
#include <unistd.h>

int rmdir(
    const char *pathname
);
```

STATUS CODES:

EPERM	The filesystem containing <code>pathname</code> does not support the removal of directories.
EFAULT	<code>pathname</code> points outside your accessible address space.
EACCES	Write access to the directory containing <code>pathname</code> was not allowed for the process's effective uid, or one of the directories in <code>pathname</code> did not allow search (execute) permission.
EPERM	The directory containing <code>pathname</code> has the stickybit (S_ISVTX) set and the process's effective uid is neither the uid of the file to be deleted nor that of the director containing it.
ENAMETOOLONG	<code>pathname</code> was too long.
ENOENT	A directory component in <code>pathname</code> does not exist or is a dangling symbolic link.
ENOTDIR	<code>pathname</code> , or a component used as a directory in <code>pathname</code> , is not, in fact, a directory.
ENOTEMPTY	<code>pathname</code> contains entries other than <code>.</code> and <code>..</code> .
EBUSY	<code>pathname</code> is the current working directory or root directory of some process
EBUSY	<code>pathname</code> is the current directory or root directory of some process.
ENOMEM	Insufficient kernel memory was available
EROFS	<code>pathname</code> refers to a file on a read-only filesystem.
ELOOP	<code>pathname</code> contains a reference to a circular symbolic link

DESCRIPTION:

`rmdir` deletes a directory, which must be empty

NOTES:

NONE

4.4.20 rename - Renames a file

CALLING SEQUENCE:

```
#include <unistd.h>

int rename(
    const char *old,
    const char *new
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix.
EBUSY	The directory is in use.
EEXIST	The named file already exists.
EINVAL	Invalid argument.
EISDIR	Attempt to open a directory for writing or to rename a file to be a directory.
EMLINK	The number of links would exceed LINK_MAX.
ENAMETOOLONG	Length of a filename string exceeds PATH_MAX and _POSIX_NO_TRUNC is in effect.
ENOENT	A file or directory does not exist.
ENOSPC	No space left on disk.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
ENOTEMPTY	Attempt to delete or rename a non-empty directory.
EROFS	Read-only file system
EXDEV	Attempt to link a file to another file system.

DESCRIPTION:

The `rename()` function causes the file known as `old` to now be known as `new`.

Ordinary files may be renamed to ordinary files, and directories may be renamed to directories; however, files cannot be converted using `rename()`. The `new` pathname may not contain a path prefix of `old`.

NOTES:

If a file already exists by the name `new`, it is removed. The `rename()` function is atomic. If the `rename()` detects an error, no files are removed. This guarantees that the `rename("x", "x")` does not remove `x`.

You may not rename dot or dot-dot.

The routine is implemented in Cygnus newlib using `link()` and `unlink()`.

4.4.21 stat - Gets information about a file

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int stat(
    const char *path,
    struct stat *buf
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix.
EBADF	Invalid file descriptor.
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.

DESCRIPTION:

The `path` argument points to a pathname for a file. Read, write, or execute permission for the file is not required, but all directories listed in `path` must be searchable. The `stat()` function obtains information about the named file and writes it to the area pointed to by `buf`.

NOTES:

NONE

4.4.22 `fstat` - Gets file status

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int fstat(
    int      fildes,
    struct stat *buf
);
```

STATUS CODES:

EBADF Invalid file descriptor

DESCRIPTION:

The `fstat()` function obtains information about the file associated with `fildes` and writes it to the area pointed to by the `buf` argument.

NOTES:

If the filesystem object referred to by `fildes` is a link, then the information returned in `buf` refers to the destination of that link. This is in contrast to `lstat()` which does not follow the link.

4.4.23 lstat - Gets file status

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int lstat(
    int      fildes,
    struct stat *buf
);
```

STATUS CODES:

EBADF Invalid file descriptor

DESCRIPTION:

The `lstat()` function obtains information about the file associated with `fildes` and writes it to the area pointed to by the `buf` argument.

NOTES:

If the filesystem object referred to by `fildes` is a link, then the information returned in `buf` refers to the link itself. This is in contrast to `fstat()` which follows the link.

The `lstat()` routine is defined by BSD 4.3 and SVR4 and not included in POSIX 1003.1b-1996.

4.4.24 access - Check permissions for a file

CALLING SEQUENCE:

```
#include <unistd.h>

int access(
    const char *pathname,
    int         mode
);
```

STATUS CODES:

EACCES	The requested access would be denied, either to the file itself or one of the directories in <code>pathname</code> .
EFAULT	<code>pathname</code> points outside your accessible address space.
EINVAL	Mode was incorrectly specified.
ENAMETOOLONG	<code>pathname</code> is too long.
ENOENT	A directory component in <code>pathname</code> would have been accessible but does not exist or was a dangling symbolic link.
ENOTDIR	A component used as a directory in <code>pathname</code> is not, in fact, a directory.
ENOMEM	Insufficient kernel memory was available.

DESCRIPTION:

`Access` checks whether the process would be allowed to read, write or test for existence of the file (or other file system object) whose name is `pathname`. If `pathname` is a symbolic link permissions of the file referred by this symbolic link are tested.

Mode is a mask consisting of one or more of `R_OK`, `W_OK`, `X_OK` and `F_OK`.

NOTES:

NONE

4.4.25 chmod - Changes file mode.

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int chmod(
    const char *path,
    mode_t      mode
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EPERM	Operation is not permitted. Process does not have the appropriate privileges or permissions to perform the requested operations.
EROFS	Read-only file system.

DESCRIPTION:

Set the file permission bits, the set user ID bit, and the set group ID bit for the file named by `path` to `mode`. If the effective user ID does not match the owner of the file and the calling process does not have the appropriate privileges, `chmod()` returns -1 and sets `errno` to `EPERM`.

NOTES:

NONE

4.4.26 fchmod - Changes permissions of a file

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <sys/stat.h>

int fchmod(
    int    fildes,
    mode_t mode
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix.
EBADF	The descriptor is not valid.
EFAULT	<code>path</code> points outside your accessible address space.
EIO	A low-level I/o error occurred while modifying the inode.
ELOOP	<code>path</code> contains a circular reference
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does no exist.
ENOMEM	Insufficient kernel memory was available.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EPERM	The effective UID does not match the owner of the file, and is not zero
EROFS	Read-only file system

DESCRIPTION:

The mode of the file given by `path` or referenced by `fildes` is changed.

NOTES:

NONE

4.4.27 getdents - Get directory entries

CALLING SEQUENCE:

```
#include <unistd.h>
#include <linux/dirent.h>
#include <linux/unistd.h>

long getdents(
    int    dd_fd,
    char *dd_buf,
    int    dd_len
);
```

STATUS CODES:

A successful call to `getdents` returns the number of bytes read. On end of directory, 0 is returned. When an error occurs, -1 is returned, and `errno` is set appropriately.

EBADF	Invalid file descriptor <code>fd</code> .
EFAULT	Argument points outside the calling process's address space.
EINVAL	Result buffer is too small.
ENOENT	No such directory.
ENOTDIR	File descriptor does not refer to a directory.

DESCRIPTION:

`getdents` reads several `dirent` structures from the directory pointed by `fd` into the memory area pointed to by `dirp`. The parameter `count` is the size of the memory area.

NOTES:

NONE

4.4.28 chown - Changes the owner and/or group of a file.

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <unistd.h>

int chown(
    const char *path,
    uid_t      owner,
    gid_t      group
);
```

STATUS CODES:

EACCES	Search permission is denied for a directory in a file's path prefix
EINVAL	Invalid argument
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist.
ENOTDIR	A component of the specified pathname was not a directory when a directory was expected.
EPERM	Operation is not permitted. Process does not have the appropriate privileges or permissions to perform the requested operations.
EROFS	Read-only file system.

DESCRIPTION:

The user ID and group ID of the file named by `path` are set to `owner` and `path`, respectively.

For regular files, the set group ID (`S_ISGID`) and set user ID (`S_ISUID`) bits are cleared.

Some systems consider it a security violation to allow the owner of a file to be changed, If users are billed for disk space usage, loaning a file to another user could result in incorrect billing. The `chown()` function may be restricted to privileged users for some or all files. The group ID can still be changed to one of the supplementary group IDs.

NOTES:

This function may be restricted for some file. The `pathconf` function can be used to test the `_PC_CHOWN_RESTRICTED` flag.

4.4.29 utime - Change access and/or modification times of an inode

CALLING SEQUENCE:

```
#include <sys/types.h>

int utime(
    const char    *filename,
    struct utimbuf *buf
);
```

STATUS CODES:

EACCES	Permission to write the file is denied
ENOENT	Filename does not exist

DESCRIPTION:

Utime changes the access and modification times of the inode specified by `filename` to the `actime` and `modtime` fields of `buf` respectively. If `buf` is `NULL`, then the access and modification times of the file are set to the current time.

NOTES:

NONE

4.4.30 `ftruncate` - truncate a file to a specified length

CALLING SEQUENCE:

```
#include <unistd.h>

int ftruncate(
    int    fd,
    size_t length
);
```

STATUS CODES:

ENOTDIR	A component of the path prefix is not a directory.
EINVAL	The pathname contains a character with the high-order bit set.
ENAMETOOLONG	A component of a pathname exceeded 255 characters, or an entire path name exceeded 1023 characters.
ENOENT	The named file does not exist.
EACCES	The named file is not writable by the user.
EACCES	Search permission is denied for a component of the path prefix.
ELOOP	Too many symbolic links were encountered in translating the path-name
EISDIR	The named file is a directory.
EROFS	The named file resides on a read-only file system
ETXTBSY	The file is a pure procedure (shared text) file that is being executed
EIO	An I/O error occurred updating the inode.
EFAULT	<code>Path</code> points outside the process's allocated address space.
EBADF	The <code>fd</code> is not a valid descriptor.

DESCRIPTION:

`truncate()` causes the file named by `path` or referenced by `fd` to be truncated to at most `length` bytes in size. If the file previously was larger than this size, the extra data is lost. With `ftruncate()`, the file must be open for writing.

NOTES:

NONE

4.4.31 truncate - truncate a file to a specified length

CALLING SEQUENCE:

```
#include <unistd.h>

int truncate(
    const char *path,
    size_t      length
);
```

STATUS CODES:

ENOTDIR	A component of the path prefix is not a directory.
EINVAL	The pathname contains a character with the high-order bit set.
ENAMETOOLONG	A component of a pathname exceeded 255 characters, or an entire path name exceeded 1023 characters.
ENOENT	The named file does not exist.
EACCES	The named file is not writable by the user.
EACCES	Search permission is denied for a component of the path prefix.
ELOOP	Too many symbolic links were encountered in translating the path-name
EISDIR	The named file is a directory.
EROFS	The named file resides on a read-only file system
ETXTBSY	The file is a pure procedure (shared text) file that is being executed
EIO	An I/O error occurred updating the inode.
EFAULT	Path points outside the process's allocated address space.
EBADF	The fd is not a valid descriptor.

DESCRIPTION:

`truncate()` causes the file named by `path` or referenced by `fd` to be truncated to at most `length` bytes in size. If the file previously was larger than this size, the extra data is lost. With `ftruncate()`, the file must be open for writing.

NOTES:

NONE

4.4.32 pathconf - Gets configuration values for files

CALLING SEQUENCE:

```
#include <unistd.h>

int pathconf(
    const char *path,
    int        name
);
```

STATUS CODES:

EINVAL	Invalid argument
EACCES	Permission to write the file is denied
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist
ENOTDIR	A component of the specified <code>path</code> was not a directory when a directory was expected.

DESCRIPTION:

`pathconf()` gets a value for the configuration option `name` for the open file descriptor `filedes`.

The possible values for `name` are:

_PC_LINK_MAX	returns the maximum number of links to the file. If <code>filedes</code> or <code>path</code> refer to a directory, then the value applies to the whole directory. The corresponding macro is <code>_POSIX_LINK_MAX</code> .
_PC_MAX_CANON	returns the maximum length of a formatted input line, where <code>filedes</code> or <code>path</code> must refer to a terminal. The corresponding macro is <code>_POSIX_MAX_CANON</code> .
_PC_MAX_INPUT	returns the maximum length of an input line, where <code>filedes</code> or <code>path</code> must refer to a terminal. The corresponding macro is <code>_POSIX_MAX_INPUT</code> .
_PC_NAME_MAX	returns the maximum length of a filename in the directory <code>path</code> or <code>filedes</code> . The process is allowed to create. The corresponding macro is <code>_POSIX_NAME_MAX</code> .
_PC_PATH_MAX	returns the maximum length of a relative pathname when <code>path</code> or <code>filedes</code> is the current working directory. The corresponding macro is <code>_POSIX_PATH_MAX</code> .
_PC_PIPE_BUF	returns the size of the pipe buffer, where <code>filedes</code> must refer to a pipe or FIFO and <code>path</code> must refer to a FIFO. The corresponding macro is <code>_POSIX_PIPE_BUF</code> .

_PC_CHOWN_RESTRICTED

returns nonzero if the `chown(2)` call may not be used on this file. If `filedes` or `path` refer to a directory, then this applies to all files in that directory. The corresponding macro is `_POSIX_CHOWN_RESTRICTED`.

NOTES:

Files with name lengths longer than the value returned for `name` equal `_PC_NAME_MAX` may exist in the given directory.

4.4.33 fpathconf - Gets configuration values for files

CALLING SEQUENCE:

```
#include <unistd.h>

int fpathconf(
    int filedes,
    int name
);
```

STATUS CODES:

EINVAL	Invalid argument
EACCES	Permission to write the file is denied
ENAMETOOLONG	Length of a filename string exceeds <code>PATH_MAX</code> and <code>_POSIX_NO_TRUNC</code> is in effect.
ENOENT	A file or directory does not exist
ENOTDIR	A component of the specified <code>path</code> was not a directory when a directory was expected.

DESCRIPTION:

`pathconf()` gets a value for the configuration option `name` for the open file descriptor `filedes`.

The possible values for `name` are:

_PC_LINK_MAX	returns the maximum number of links to the file. If <code>filedes</code> or <code>path</code> refer to a directory, then the value applies to the whole directory. The corresponding macro is <code>_POSIX_LINK_MAX</code> .
_PC_MAX_CANON	returns the maximum length of a formatted input line, where <code>filedes</code> or <code>path</code> must refer to a terminal. The corresponding macro is <code>_POSIX_MAX_CANON</code> .
_PC_MAX_INPUT	returns the maximum length of an input line, where <code>filedes</code> or <code>path</code> must refer to a terminal. The corresponding macro is <code>_POSIX_MAX_INPUT</code> .
_PC_NAME_MAX	returns the maximum length of a filename in the directory <code>path</code> or <code>filedes</code> . The process is allowed to create. The corresponding macro is <code>_POSIX_NAME_MAX</code> .
_PC_PATH_MAX	returns the maximum length of a relative pathname when <code>path</code> or <code>filedes</code> is the current working directory. The corresponding macro is <code>_POSIX_PATH_MAX</code> .
_PC_PIPE_BUF	returns the size of the pipe buffer, where <code>filedes</code> must refer to a pipe or FIFO and <code>path</code> must refer to a FIFO. The corresponding macro is <code>_POSIX_PIPE_BUF</code> .

_PC_CHOWN_RESTRICTED

returns nonzero if the `chown()` call may not be used on this file. If `filedes` or `path` refer to a directory, then this applies to all files in that directory. The corresponding macro is `_POSIX_CHOWN_RESTRICTED`.

NOTES:

NONE

4.4.34 mknod - create a directory

CALLING SEQUENCE:

```
#include <unistd.h>
#include <fcntl.h>
#include <sys/types.h>
#include <sys/stat.h>

long mknod(
    const char *pathname,
    mode_t      mode,
    dev_t       dev
);
```

STATUS CODES:

`mknod` returns zero on success, or -1 if an error occurred (in which case, `errno` is set appropriately).

ENAMETOOLONG	<code>pathname</code> was too long.
ENOENT	A directory component in <code>pathname</code> does not exist or is a dangling symbolic link.
ENOTDIR	A component used in the directory <code>pathname</code> is not, in fact, a directory.
ENOMEM	Insufficient kernel memory was available
EROFS	<code>pathname</code> refers to a file on a read-only filesystem.
ELOOP	<code>pathname</code> contains a reference to a circular symbolic link, ie a symbolic link whose expansion contains a reference to itself.
ENOSPC	The device containing <code>pathname</code> has no room for the new node.

DESCRIPTION:

`mknod` attempts to create a filesystem node (file, device special file or named pipe) named `pathname`, specified by `mode` and `dev`.

`mode` specifies both the permissions to use and the type of node to be created.

It should be a combination (using bitwise OR) of one of the file types listed below and the permissions for the new node.

The permissions are modified by the process's `umask` in the usual way: the permissions of the created node are `(mode & ~umask)`.

The file type should be one of `S_IFREG`, `S_IFCHR`, `S_IFBLK` and `S_IFIFO` to specify a normal file (which will be created empty), character special file, block special file or FIFO (named pipe), respectively, or zero, which will create a normal file.

If the file type is `S_IFCHR` or `S_IFBLK` then `dev` specifies the major and minor numbers of the newly created device special file; otherwise it is ignored.

The newly created node will be owned by the effective uid of the process. If the directory containing the node has the set group id bit set, or if the filesystem is mounted with BSD group semantics, the new node will inherit the group ownership from its parent directory; otherwise it will be owned by the effective gid of the process.

NOTES:

NONE

5 Input and Output Primitives Manager

5.1 Introduction

The input and output primitives manager is ...

The directives provided by the input and output primitives manager are:

- `pipe` - Create an Inter-Process Channel
- `dup` - Duplicates an open file descriptor
- `dup2` - Duplicates an open file descriptor
- `close` - Closes a file
- `read` - Reads from a file
- `write` - Writes to a file
- `fcntl` - Manipulates an open file descriptor
- `lseek` - Reposition read/write file offset
- `fsync` - Synchronize file complete in-core state with that on disk
- `fdatasync` - Synchronize file in-core data with that on disk
- `sync` - Schedule file system updates
- `mount` - Mount a file system
- `unmount` - Unmount file systems
- `readv` - Vectored read from a file
- `writew` - Vectored write to a file
- `aio_read` - Asynchronous Read
- `aio_write` - Asynchronous Write
- `lio_listio` - List Directed I/O
- `aio_error` - Retrieve Error Status of Asynchronous I/O Operation
- `aio_return` - Retrieve Return Status Asynchronous I/O Operation
- `aio_cancel` - Cancel Asynchronous I/O Request
- `aio_suspend` - Wait for Asynchronous I/O Request
- `aio_fsync` - Asynchronous File Synchronization

5.2 Background

There is currently no text in this section.

5.3 Operations

There is currently no text in this section.

5.4 Directives

This section details the input and output primitives manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

5.4.1 pipe - Create an Inter-Process Channel

CALLING SEQUENCE:

```
int pipe(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.2 dup - Duplicates an open file descriptor

CALLING SEQUENCE:

```
#include <unistd.h>

int dup(
    int fildes
);
```

STATUS CODES:

EBADF	Invalid file descriptor.
EINTR	Function was interrupted by a signal.
EMFILE	The process already has the maximum number of file descriptors open and tried to open a new one.

DESCRIPTION:

The `dup` function returns the lowest numbered available file descriptor. This new descriptor refers to the same open file as the original descriptor and shares any locks.

NOTES:

NONE

5.4.3 dup2 - Duplicates an open file descriptor

CALLING SEQUENCE:

```
#include <unistd.h>

int dup2(
    int fildes,
    int fildes2
);
```

STATUS CODES:

EBADF	Invalid file descriptor.
EINTR	Function was interrupted by a signal.
EMFILE	The process already has the maximum number of file descriptors open and tried to open a new one.

DESCRIPTION:

`dup2` creates a copy of the file descriptor `oldfd`.

The old and new descriptors may be used interchangeably. They share locks, file position pointers and flags; for example, if the file position is modified by using `lseek` on one of the descriptors, the position is also changed for the other.

NOTES:

NONE

5.4.4 close - Closes a file

CALLING SEQUENCE:

```
#include <unistd.h>

int close(
    int fildes
);
```

STATUS CODES:

EBADF Invalid file descriptor
EINTR Function was interrupted by a signal.

DESCRIPTION:

The `close()` function deallocates the file descriptor named by `fildes` and makes it available for reuse. All outstanding record locks owned by this process for the file are unlocked.

NOTES:

A signal can interrupt the `close()` function. In that case, `close()` returns -1 with `errno` set to `EINTR`. The file may or may not be closed.

5.4.5 read - Reads from a file

CALLING SEQUENCE:

```
#include <unistd.h>

int read(
    int          fildes,
    void         *buf,
    unsigned int nbyte
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** to one of the following:

EAGAIN	The O_NONBLOCK flag is set for a file descriptor and the process would be delayed in the I/O operation.
EBADF	Invalid file descriptor
EINTR	Function was interrupted by a signal.
EIO	Input or output error
EINVAL	Bad buffer pointer

DESCRIPTION:

The `read()` function reads `nbyte` bytes from the file associated with `fildes` into the buffer pointed to by `buf`.

The `read()` function returns the number of bytes actually read and placed in the buffer. This will be less than `nbyte` if:

- The number of bytes left in the file is less than `nbyte`.
- The `read()` request was interrupted by a signal.
- The file is a pipe or FIFO or special file with less than `nbytes` immediately available for reading.

When attempting to read from any empty pipe or FIFO:

- If no process has the pipe open for writing, zero is returned to indicate end-of-file.
- If some process has the pipe open for writing and O_NONBLOCK is set, -1 is returned and **errno** is set to EAGAIN.
- If some process has the pipe open for writing and O_NONBLOCK is clear, `read()` waits for some data to be written or the pipe to be closed.

When attempting to read from a file other than a pipe or FIFO and no data is available.

- If O_NONBLOCK is set, -1 is returned and **errno** is set to EAGAIN.
- If O_NONBLOCK is clear, `read()` waits for some data to become available.
- The O_NONBLOCK flag is ignored if data is available.

NOTES:

NONE

5.4.6 write - Writes to a file

CALLING SEQUENCE:

```
#include <unistd.h>

int write(
    int          fildes,
    const void   *buf,
    unsigned int  nbytes
);
```

STATUS CODES:

EAGAIN	The O_NONBLOCK flag is set for a file descriptor and the process would be delayed in the I/O operation.
EBADF	Invalid file descriptor
EFBIG	An attempt was made to write to a file that exceeds the maximum file size
EINTR	The function was interrupted by a signal.
EIO	Input or output error.
ENOSPC	No space left on disk.
EPIPE	Attempt to write to a pipe or FIFO with no reader.
EINVAL	Bad buffer pointer

DESCRIPTION:

The `write()` function writes `nbyte` from the array pointed to by `buf` into the file associated with `fildes`.

If `nbyte` is zero and the file is a regular file, the `write()` function returns zero and has no other effect. If `nbyte` is zero and the file is a special file, the results are not portable.

The `write()` function returns the number of bytes written. This number will be less than `nbytes` if there is an error. It will never be greater than `nbytes`.

NOTES:

NONE

5.4.7 fcntl - Manipulates an open file descriptor

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <fcntl.h>
#include <unistd.h>

int fcntl(
    int fildes,
    int cmd
);
```

STATUS CODES:

EACCESS	Search permission is denied for a directory in a file's path prefix.
EAGAIN	The O_NONBLOCK flag is set for a file descriptor and the process would be delayed in the I/O operation.
EBADF	Invalid file descriptor
EDEADLK	An fcntl with function F_SETLK would cause a deadlock.
EINTR	The function was interrupted by a signal.
EINVAL	Invalid argument
EMFILE	Too many file descriptors or in use by the process.
ENOLCK	No locks available

DESCRIPTION:

fcntl() performs one of various miscellaneous operations on fd. The operation in question is determined by cmd:

F_DUPFD	Makes arg be a copy of fd, closing fd first if necessary. The same functionality can be more easily achieved by using dup2(). The old and new descriptors may be used interchangeably. They share locks, file position pointers and flags; for example, if the file position is modified by using lseek() on one of the descriptors, the position is also changed for the other. The two descriptors do not share the close-on-exec flag, however. The close-on-exec flag of the copy is off, meaning that it will be closed on exec. On success, the new descriptor is returned.
F_GETFD	Read the close-on-exec flag. If the low-order bit is 0, the file will remain open across exec, otherwise it will be closed.
F_SETFD	Set the close-on-exec flag to the value specified by arg (only the least significant bit is used).

- F_GETFL** Read the descriptor's flags (all flags (as set by `open()`) are returned).
- F_SETFL** Set the descriptor's flags to the value specified by `arg`. Only `O_APPEND` and `O_NONBLOCK` may be set.
The flags are shared between copies (made with `dup()` etc.) of the same file descriptor.
The flags and their semantics are described in `open()`.
- F_GETLK, F_SETLK and F_SETLKW** Manage discretionary file locks. The third argument `arg` is a pointer to a struct flock (that may be overwritten by this call).
- F_GETLK** Return the flock structure that prevents us from obtaining the lock, or set the `l_type` field of the lock to `F_UNLCK` if there is no obstruction.
- F_SETLK** The lock is set (when `l_type` is `F_RDLCK` or `F_WRLCK`) or cleared (when it is `F_UNLCK`. If lock is held by someone else, this call returns -1 and sets `errno` to `EACCES` or `EAGAIN`.
- F_SETLKW** Like `F_SETLK`, but instead of returning an error we wait for the lock to be released.
- F_GETOWN** Get the process ID (or process group) of the owner of a socket.
Process groups are returned as negative values.
- F_SETOWN** Set the process or process group that owns a socket.
For these commands, ownership means receiving `SIGIO` or `SIGURG` signals.
Process groups are specified using negative values.

NOTES:

The errors returned by `dup2` are different from those returned by `F_DUPFD`.

5.4.8 lseek - Reposition read/write file offset

CALLING SEQUENCE:

```
#include <sys/types.h>
#include <unistd.h>

int lseek(
    int    fildes,
    off_t  offset,
    int    whence
);
```

STATUS CODES:

EBADF	<code>fildes</code> is not an open file descriptor.
ESPIPE	<code>fildes</code> is associated with a pipe, socket or FIFO.
EINVAL	<code>whence</code> is not a proper value.

DESCRIPTION:

The `lseek` function repositions the offset of the file descriptor `fildes` to the argument `offset` according to the directive `whence`. The argument `fildes` must be an open file descriptor. `Lseek` repositions the file pointer `fildes` as follows:

- If `whence` is `SEEK_SET`, the offset is set to `offset` bytes.
- If `whence` is `SEEK_CUR`, the offset is set to its current location plus `offset` bytes.
- If `whence` is `SEEK_END`, the offset is set to the size of the file plus `offset` bytes.

The `lseek` function allows the file offset to be set beyond the end of the existing end-of-file of the file. If data is later written at this point, subsequent reads of the data in the gap return bytes of zeros (until data is actually written into the gap).

Some devices are incapable of seeking. The value of the pointer associated with such a device is undefined.

NOTES:

NONE

5.4.9 fsync - Synchronize file complete in-core state with that on disk

CALLING SEQUENCE:

```
int fsync(  
    int fd  
);
```

STATUS CODES:

On success, zero is returned. On error, -1 is returned, and `errno` is set appropriately.

EBADF	<code>fd</code> is not a valid descriptor open for writing
EINVAL	<code>fd</code> is bound to a special file which does not support synchronization
EROFS	<code>fd</code> is bound to a special file which does not support synchronization
EIO	An error occurred during synchronization

DESCRIPTION:

`fsync` copies all in-core parts of a file to disk.

NOTES:

NONE

5.4.10 `fdatasync` - Synchronize file in-core data with that on disk

CALLING SEQUENCE:

```
int fdatasync(  
    int fd  
);
```

STATUS CODES:

On success, zero is returned. On error, -1 is returned, and `errno` is set appropriately.

EBADF	<code>fd</code> is not a valid file descriptor open for writing.
EINVAL	<code>fd</code> is bound to a special file which does not support synchronization.
EIO	An error occurred during synchronization.
EROFS	<code>fd</code> is bound to a special file which does not support synchronization.

DESCRIPTION:

`fdatasync` flushes all data buffers of a file to disk (before the system call returns). It resembles `fsync` but is not required to update the metadata such as access time.

Applications that access databases or log files often write a tiny data fragment (e.g., one line in a log file) and then call `fsync` immediately in order to ensure that the written data is physically stored on the harddisk. Unfortunately, `fsync` will always initiate two write operations: one for the newly written data and another one in order to update the modification time stored in the inode. If the modification time is not a part of the transaction concept `fdatasync` can be used to avoid unnecessary inode disk write operations.

NOTES:

NONE

5.4.11 sync - Schedule file system updates

CALLING SEQUENCE:

```
void sync(void);
```

STATUS CODES:

NONE

DESCRIPTION:

The `sync` service causes all information in memory that updates file systems to be scheduled for writing out to all file systems.

NOTES:

The writing of data to the file systems is only guaranteed to be scheduled upon return. It is not necessarily complete upon return from `sync`.

5.4.12 mount - Mount a file system

CALLING SEQUENCE:

```
#include <libio.h>

int mount(
    rtems_filesystem_mount_table_entry_t **mt_entry,
    rtems_filesystem_operations_table    *fs_ops,
    rtems_filesystem_options_t          fsoptions,
    char                                 *device,
    char                                 *mount_point
);
```

STATUS CODES:

EXXX

DESCRIPTION:

The `mount` routines mounts the filesystem class which uses the filesystem operations specified by `fs_ops` and `fsoptions`. The filesystem is mounted at the directory `mount_point` and the mode of the mounted filesystem is specified by `fsoptions`. If this filesystem class requires a device, then the name of the device must be specified by `device`.

If this operation succeeds, the mount table entry for the mounted filesystem is returned in `mt_entry`.

NOTES:

NONE

5.4.13 unmount - Unmount file systems

CALLING SEQUENCE:

```
#include <libio.h>

int unmount(
    const char *mount_path
);
```

STATUS CODES:

EXXX

DESCRIPTION:

The `unmount` routine removes the attachment of the filesystem specified by `mount_path`.

NOTES:

NONE

5.4.14 readv - Vectored read from a file

CALLING SEQUENCE:

```
#include <sys/uio.h>

ssize_t readv(
    int          fd,
    const struct iovec *iov,
    int          iovcnt
);
```

STATUS CODES:

In addition to the errors detected by Input and Output Primitives Manager `read` - Reads from a file, `read()`, this routine may return -1 and sets `errno` based upon the following errors:

EINVAL	The sum of the <code>iov_len</code> values in the <code>iov</code> array overflowed an <code>ssize_t</code> .
EINVAL	The <code>iovcnt</code> argument was less than or equal to 0, or greater than <code>IOV_MAX</code> .

DESCRIPTION:

The `readv()` function is equivalent to `read()` except as described here. The `readv()` function shall place the input data into the `iovcnt` buffers specified by the members of the `iov` array: `iov[0]`, `iov[1]`, ..., `iov[iovcnt-1]`.

Each `iovec` entry specifies the base address and length of an area in memory where data should be placed. The `readv()` function always fills an area completely before proceeding to the next.

NOTES:

NONE

5.4.15 writev - Vectored write to a file

CALLING SEQUENCE:

```
#include <sys/uio.h>

ssize_t writev(
    int          fd,
    const struct iovec *iov,
    int          iovcnt
);
```

STATUS CODES:

In addition to the errors detected by Input and Output Primitives Manager write - Write to a file, `write()`, this routine may return -1 and sets `errno` based upon the following errors:

- | | |
|---------------|---|
| EINVAL | The sum of the <code>iov_len</code> values in the <code>iov</code> array overflowed an <code>ssize_t</code> . |
| EINVAL | The <code>iovcnt</code> argument was less than or equal to 0, or greater than <code>IOV_MAX</code> . |

DESCRIPTION:

The `writev()` function is equivalent to `write()`, except as noted here. The `writev()` function gathers output data from the `iovcnt` buffers specified by the members of the `iov` array: `iov[0]`, `iov[1]`, ..., `iov[iovcnt-1]`. The `iovcnt` argument is valid if greater than 0 and less than or equal to `IOV_MAX`.

Each `iovec` entry specifies the base address and length of an area in memory from which data should be written. The `writev()` function always writes a complete area before proceeding to the next.

If `fd` refers to a regular file and all of the `iov_len` members in the array pointed to by `iov` are 0, `writev()` returns 0 and has no other effect. For other file types, the behavior is unspecified by POSIX.

NOTES:

NONE

5.4.16 aio_read - Asynchronous Read

CALLING SEQUENCE:

```
int aio_read(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.17 aio_write - Asynchronous Write

CALLING SEQUENCE:

```
int aio_write(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.18 lio_listio - List Directed I/O

CALLING SEQUENCE:

```
int lio_listio(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.19 aio_error - Retrieve Error Status of Asynchronous I/O Operation

CALLING SEQUENCE:

```
int aio_error(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.20 aio_return - Retrieve Return Status Asynchronous I/O Operation

CALLING SEQUENCE:

```
int aio_return(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.21 aio_cancel - Cancel Asynchronous I/O Request

CALLING SEQUENCE:

```
int aio_cancel(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.22 aio_suspend - Wait for Asynchronous I/O Request

CALLING SEQUENCE:

```
int aio_suspend(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

5.4.23 aio_fsync - Asynchronous File Synchronization

CALLING SEQUENCE:

```
int aio_fsync(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

6 Device- and Class- Specific Functions Manager

6.1 Introduction

The device- and class- specific functions manager is ...

The directives provided by the device- and class- specific functions manager are:

- `cfgetispeed` - Reads terminal input baud rate
- `cfgetospeed` - Reads terminal output baud rate
- `cfsetispeed` - Sets terminal input baud rate
- `cfsetospeed` - Set terminal output baud rate
- `tcgetattr` - Gets terminal attributes
- `tcsetattr` - Set terminal attributes
- `tcsendbreak` - Sends a break to a terminal
- `tcdrain` - Waits for all output to be transmitted to the terminal
- `tcflush` - Discards terminal data
- `tcflow` - Suspends/restarts terminal output
- `tcgetpgrp` - Gets foreground process group ID
- `tcsetpgrp` - Sets foreground process group ID

6.2 Background

There is currently no text in this section.

6.3 Operations

There is currently no text in this section.

6.4 Directives

This section details the device- and class- specific functions manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

6.4.1 cfgetispeed - Reads terminal input baud rate

CALLING SEQUENCE:

```
#include <termios.h>

int cfgetispeed(
    const struct termios *p
);
```

STATUS CODES:

The `cfgetispeed()` function returns a code for baud rate.

DESCRIPTION:

The `cfsetispeed()` function stores a code for the terminal speed stored in a struct `termios`. The codes are defined in `<termios.h>` by the macros `BO`, `B50`, `B75`, `B110`, `B134`, `B150`, `B200`, `B300`, `B600`, `B1200`, `B1800`, `B2400`, `B4800`, `B9600`, `B19200`, and `B38400`.

The `cfsetispeed()` function does not do anything to the hardware. It merely stores a value for use by `tcsetattr()`.

NOTES:

Baud rates are defined by symbols, such as `B110`, `B1200`, `B2400`. The actual number returned for any given speed may change from system to system.

6.4.2 cfgetospeed - Reads terminal output baud rate

CALLING SEQUENCE:

```
#include <termios.h>

int cfgetospeed(
    const struct termios *p
);
```

STATUS CODES:

The `cfgetospeed()` function returns the `termios` code for the baud rate.

DESCRIPTION:

The `cfgetospeed()` function returns a code for the terminal speed stored in a `struct termios`. The codes are defined in `<termios.h>` by the macros `BO`, `B50`, `B75`, `B110`, `B134`, `B150`, `B200`, `B300`, `B600`, `B1200`, `B1800`, `B2400`, `B4800`, `B9600`, `B19200`, and `B38400`.

The `cfgetospeed()` function does not do anything to the hardware. It merely returns the value stored by a previous call to `tcgetattr()`.

NOTES:

Baud rates are defined by symbols, such as `B110`, `B1200`, `B2400`. The actual number returned for any given speed may change from system to system.

6.4.3 cfsetispeed - Sets terminal input baud rate

CALLING SEQUENCE:

```
#include <termios.h>

int cfsetispeed(
    struct termios *p,
    speed_t        speed
);
```

STATUS CODES:

The `cfsetispeed()` function returns a zero when successful and returns -1 when an error occurs.

DESCRIPTION:

The `cfsetispeed()` function stores a code for the terminal speed stored in a `struct termios`. The codes are defined in `<termios.h>` by the macros `B0`, `B50`, `B75`, `B110`, `B134`, `B150`, `B200`, `B300`, `B600`, `B1200`, `B1800`, `B2400`, `B4800`, `B9600`, `B19200`, and `B38400`.

NOTES:

This function merely stores a value in the `termios` structure. It does not change the terminal speed until a `tcsetattr()` is done. It does not detect impossible terminal speeds.

6.4.4 cfsetospeed - Sets terminal output baud rate

CALLING SEQUENCE:

```
#include <termios.h>

int cfsetospeed(
    struct termios *p,
    speed_t        speed
);
```

STATUS CODES:

The `cfsetospeed()` function returns a zero when successful and returns -1 when an error occurs.

DESCRIPTION:

The `cfsetospeed()` function stores a code for the terminal speed stored in a struct `termios`. The codes are defined in `<termios.h>` by the macros `B0`, `B50`, `B75`, `B110`, `B134`, `B150`, `B200`, `B300`, `B600`, `B1200`, `B1800`, `B2400`, `B4800`, `B9600`, `B19200`, and `B38400`.

The `cfsetospeed()` function does not do anything to the hardware. It merely stores a value for use by `tcsetattr()`.

NOTES:

This function merely stores a value in the `termios` structure. It does not change the terminal speed until a `tcsetattr()` is done. It does not detect impossible terminal speeds.

6.4.5 tcgetattr - Gets terminal attributes

CALLING SEQUENCE:

```
#include <termios.h>
#include <unistd.h>

int tcgetattr(
    int          fildes,
    struct termios *p
);
```

STATUS CODES:

EBADF Invalid file descriptor

ENOTTY Terminal control function attempted for a file that is not a terminal.

DESCRIPTION:

The `tcgetattr()` gets the parameters associated with the terminal referred to by `fildes` and stores them into the `termios()` structure pointed to by `termios_p`.

NOTES:

NONE

6.4.6 tcsetattr - Set terminal attributes

CALLING SEQUENCE:

```
#include <termios.h>
#include <unistd.h>

int tcsetattr(
    int          fildes,
    int          options,
    const struct termios *tp
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

6.4.7 tcseendbreak - Sends a break to a terminal

CALLING SEQUENCE:

```
int tcseendbreak(  
    int fd  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

6.4.8 tcdrain - Waits for all output to be transmitted to the terminal.

CALLING SEQUENCE:

```
#include <termios.h>
#include <unistd.h>

int tcdrain(
    int fildes
);
```

STATUS CODES:

EBADF	Invalid file descriptor
EINTR	Function was interrupted by a signal
ENOTTY	Terminal control function attempted for a file that is not a terminal.

DESCRIPTION:

The `tcdrain()` function waits until all output written to `fildes` has been transmitted.

NOTES:

NONE

6.4.9 tcflush - Discards terminal data

CALLING SEQUENCE:

```
int tcflush(  
    int fd  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

6.4.10 tcflow - Suspends/restarts terminal output.**CALLING SEQUENCE:**

```
int tcflow(  
    int fd  
);
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

This routine is not currently supported by RTEMS but could be in a future version.

6.4.11 tcgetpgrp - Gets foreground process group ID

CALLING SEQUENCE:

```
int tcgetpgrp(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

This routine is not currently supported by RTEMS but could be in a future version.

6.4.12 tcsetpgrp - Sets foreground process group ID**CALLING SEQUENCE:**

```
int tcsetpgrp(  
);
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

This routine is not currently supported by RTEMS but could be in a future version.

7 Language-Specific Services for the C Programming Language Manager

7.1 Introduction

The language-specific services for the C programming language manager is ...

The directives provided by the language-specific services for the C programming language manager are:

- `setlocale` - Set the Current Locale
- `fileno` - Obtain File Descriptor Number for this File
- `fdopen` - Associate Stream with File Descriptor
- `flockfile` - Acquire Ownership of File Stream
- `ftrylockfile` - Poll to Acquire Ownership of File Stream
- `funlockfile` - Release Ownership of File Stream
- `getc_unlocked` - Get Character without Locking
- `getchar_unlocked` - Get Character from stdin without Locking
- `putc_unlocked` - Put Character without Locking
- `putchar_unlocked` - Put Character to stdin without Locking
- `setjmp` - Save Context for Non-Local Goto
- `longjmp` - Non-Local Jump to a Saved Context
- `sigsetjmp` - Save Context with Signal Status for Non-Local Goto
- `siglongjmp` - Non-Local Jump with Signal Status to a Saved Context
- `tzset` - Initialize Time Conversion Information
- `strtok_r` - Reentrant Extract Token from String
- `asctime_r` - Reentrant struct tm to ASCII Time Conversion
- `ctime_r` - Reentrant time_t to ASCII Time Conversion
- `gmtime_r` - Reentrant UTC Time Conversion
- `localtime_r` - Reentrant Local Time Conversion
- `rand_r` - Reentrant Random Number Generation

7.2 Background

There is currently no text in this section.

7.3 Operations

There is currently no text in this section.

7.4 Directives

This section details the language-specific services for the C programming language manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

7.4.1 setlocale - Set the Current Locale

CALLING SEQUENCE:

```
int setlocale(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.2 fileno - Obtain File Descriptor Number for this File

CALLING SEQUENCE:

```
int fileno(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.3 fdopen - Associate Stream with File Descriptor

CALLING SEQUENCE:

```
int fdopen(  
);
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

7.4.4 flockfile - Acquire Ownership of File Stream

CALLING SEQUENCE:

```
int flockfile(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.5 ftrylockfile - Poll to Acquire Ownership of File Stream

CALLING SEQUENCE:

```
int ftrylockfile(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.6 funlockfile - Release Ownership of File Stream

CALLING SEQUENCE:

```
int funlockfile(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.7 `getc_unlocked` - Get Character without Locking

CALLING SEQUENCE:

```
int getc_unlocked(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.8 `getchar_unlocked` - Get Character from `stdin` without Locking

CALLING SEQUENCE:

```
int getchar_unlocked(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.9 `putc_unlocked` - Put Character without Locking

CALLING SEQUENCE:

```
int putc_unlocked(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.10 putchar_unlocked - Put Character to stdin without Locking

CALLING SEQUENCE:

```
int putchar_unlocked(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.11 setjmp - Save Context for Non-Local Goto

CALLING SEQUENCE:

```
int setjmp(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.12 longjmp - Non-Local Jump to a Saved Context

CALLING SEQUENCE:

```
int longjmp(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.13 sigsetjmp - Save Context with Signal Status for Non-Local Goto

CALLING SEQUENCE:

```
int sigsetjmp(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.14 siglongjmp - Non-Local Jump with Signal Status to a Saved Context

CALLING SEQUENCE:

```
int siglongjmp(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.15 tzset - Initialize Time Conversion Information

CALLING SEQUENCE:

```
int tzset(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.16 strtok_r - Reentrant Extract Token from String

CALLING SEQUENCE:

```
int strtok_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.17 asctime_r - Reentrant struct tm to ASCII Time Conversion

CALLING SEQUENCE:

```
int asctime_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.18 `ctime_r` - Reentrant `time_t` to ASCII Time Conversion

CALLING SEQUENCE:

```
int ctime_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.19 gmtime_r - Reentrant UTC Time Conversion

CALLING SEQUENCE:

```
int gmtime_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.20 localtime_r - Reentrant Local Time Conversion

CALLING SEQUENCE:

```
int localtime_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

7.4.21 rand_r - Reentrant Random Number Generation

CALLING SEQUENCE:

```
int rand_r(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8 System Databases Manager

8.1 Introduction

The system databases manager is ...

The directives provided by the system databases manager are:

- `getgrgid` - Get Group File Entry for ID
- `getgrgid_r` - Reentrant Get Group File Entry
- `getgrnam` - Get Group File Entry for Name
- `getgrnam_r` - Reentrant Get Group File Entry for Name
- `getpwuid` - Get Password File Entry for UID
- `getpwuid_r` - Reentrant Get Password File Entry for UID
- `getpwnam` - Get Password File Entry for Name
- `getpwnam_r` - Reentrant Get Password File Entry for Name

8.2 Background

There is currently no text in this section.

8.3 Operations

There is currently no text in this section.

8.4 Directives

This section details the system databases manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

8.4.1 getgrgid - Get Group File Entry for ID

CALLING SEQUENCE:

```
int getgrgid(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8.4.2 getgrgid_r - Reentrant Get Group File Entry

CALLING SEQUENCE:

```
int getgrgid_r(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8.4.3 getgrnam - Get Group File Entry for Name

CALLING SEQUENCE:

```
int getgrnam(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8.4.4 getgrnam_r - Reentrant Get Group File Entry for Name

CALLING SEQUENCE:

```
int getgrnam_r(  
);
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

8.4.5 getpwuid - Get Password File Entry for UID

CALLING SEQUENCE:

```
int getpwuid(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8.4.6 getpwuid_r - Reentrant Get Password File Entry for UID

CALLING SEQUENCE:

```
int getpwuid_r(  
);
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

8.4.7 getpwnam - Password File Entry for Name

CALLING SEQUENCE:

```
int getpwnam(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

8.4.8 getpwnam_r - Reentrant Get Password File Entry for Name

CALLING SEQUENCE:

```
int getpwnam_r(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

9 Semaphore Manager

9.1 Introduction

The semaphore manager provides functions to allocate, delete, and control semaphores. This manager is based on the POSIX 1003.1 standard.

The directives provided by the semaphore manager are:

- `sem_init` - Initialize an unnamed semaphore
- `sem_destroy` - Destroy an unnamed semaphore
- `sem_open` - Open a named semaphore
- `sem_close` - Close a named semaphore
- `sem_unlink` - Remove a named semaphore
- `sem_wait` - Lock a semaphore
- `sem_trywait` - Lock a semaphore
- `sem_timedwait` - Wait on a Semaphore for a Specified Time
- `sem_post` - Unlock a semaphore
- `sem_getvalue` - Get the value of a semaphore

9.2 Background

9.2.1 Theory

Semaphores are used for synchronization and mutual exclusion by indicating the availability and number of resources. The task (the task which is returning resources) notifying other tasks of an event increases the number of resources held by the semaphore by one. The task (the task which will obtain resources) waiting for the event decreases the number of resources held by the semaphore by one. If the number of resources held by a semaphore is insufficient (namely 0), the task requiring resources will wait until the next time resources are returned to the semaphore. If there is more than one task waiting for a semaphore, the tasks will be placed in the queue.

9.2.2 "sem_t" Structure

The `sem_t` structure is used to represent semaphores. It is passed as an argument to the semaphore directives and is defined as follows:

```
typedef int sem_t;
```

9.2.3 Building a Semaphore Attribute Set

9.3 Operations

9.3.1 Using as a Binary Semaphore

Although POSIX supports mutexes, they are only visible between threads. To work between processes, a binary semaphore must be used.

Creating a semaphore with a limit on the count of 1 effectively restricts the semaphore to being a binary semaphore. When the binary semaphore is available, the count is 1. When the binary semaphore is unavailable, the count is 0.

Since this does not result in a true binary semaphore, advanced binary features like the Priority Inheritance and Priority Ceiling Protocols are not available.

There is currently no text in this section.

9.4 Directives

This section details the semaphore manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

9.4.1 sem_init - Initialize an unnamed semaphore

CALLING SEQUENCE:

```
int sem_init(  
    sem_t      *sem,  
    int        pshared,  
    unsigned int value  
);
```

STATUS CODES:

EINVAL	The value argument exceeds SEM_VALUE_MAX
ENOSPC	A resource required to initialize the semaphore has been exhausted The limit on semaphores (SEM_VALUE_MAX) has been reached
ENOSYS	The function sem_init is not supported by this implementation
EPERM	The process lacks appropriate privileges to initialize the semaphore

DESCRIPTION:

The sem_init function is used to initialize the unnamed semaphore referred to by "sem". The value of the initialized semaphore is the parameter "value". The semaphore remains valid until it is destroyed.

ADD MORE HERE XXX

NOTES:

If the functions completes successfully, it shall return a value of zero. Otherwise, it shall return a value of -1 and set "errno" to specify the error that occurred.

Multiprocessing is currently not supported in this implementation.

9.4.2 sem_destroy - Destroy an unnamed semaphore

CALLING SEQUENCE:

```
int sem_destroy(  
    sem_t *sem  
);
```

STATUS CODES:

EINVAL	The value argument exceeds SEM_VALUE_MAX
ENOSYS	The function sem_init is not supported by this implementation
EBUSY	There are currently processes blocked on the semaphore

DESCRIPTION:

The sem_destroy function is used to destroy an unnamed semaphore referred to by "sem". sem_destroy can only be used on a semaphore that was created using sem_init.

NOTES:

If the function completes successfully, it shall return a value of zero. Otherwise, it shall return a value of -1 and set "errno" to specify the error that occurred.

Multiprocessing is currently not supported in this implementation.

9.4.3 sem_open - Open a named semaphore

CALLING SEQUENCE:

```
int sem_open(
    const char *name,
    int        oflag
);
```

ARGUMENTS:

The following flag bit may be set in oflag:

O_CREAT - Creates the semaphore if it does not already exist. If **O_CREAT** is set and the semaphore already exists then **O_CREAT** has no effect. Otherwise, `sem_open()` creates a semaphore. The **O_CREAT** flag requires the third and fourth argument: mode and value of type `mode_t` and unsigned int, respectively.

O_EXCL - If **O_EXCL** and **O_CREAT** are set, all call to `sem_open()` shall fail if the semaphore name exists

STATUS CODES:

EACCES	Valid name specified but oflag permissions are denied, or the semaphore name specified does not exist and permission to create the named semaphore is denied.
EEXIST	O_CREAT and O_EXCL are set and the named semaphore already exists.
EINTR	The <code>sem_open()</code> operation was interrupted by a signal.
EINVAL	The <code>sem_open()</code> operation is not supported for the given name.
EMFILE	Too many semaphore descriptors or file descriptors in use by this process.
ENAMETOOLONG	The length of the name exceed <code>PATH_MAX</code> or name component is longer than <code>NAME_MAX</code> while <code>POSIX_NO_TRUNC</code> is in effect.
ENOENT	O_CREAT is not set and the named semaphore does not exist.
ENOSPC	There is insufficient space for the creation of a new named semaphore.
ENOSYS	The function <code>sem_open()</code> is not supported by this implementation.

DESCRIPTION:

The `sem_open()` function establishes a connection between a specified semaphore and a process. After a call to `sem_open` with a specified semaphore name, a process can reference to semaphore by the associated name using the address returned by the call. The oflag arguments listed above control the state of the semaphore by determining if the semaphore is created or accessed by a call to `sem_open()`.

NOTES:

9.4.4 sem_close - Close a named semaphore

CALLING SEQUENCE:

```
int sem_close(  
    sem_t *sem_close  
);
```

STATUS CODES:

EACCES The semaphore argument is not a valid semaphore descriptor.

ENOSYS The function sem_close is not supported by this implementation.

DESCRIPTION:

The sem_close() function is used to indicate that the calling process is finished using the named semaphore indicated by sem. The function sem_close deallocates any system resources that were previously allocated by a sem_open system call. If sem_close() completes successfully it returns a 1, otherwise a value of -1 is return and errno is set.

NOTES:

9.4.5 sem_unlink - Unlink a semaphore

CALLING SEQUENCE:

```
int sem_unlink(  
    const char *name  
);
```

STATUS CODES:

EACCESS	Permission is denied to unlink a semaphore.
ENAMETOOLONG	The length of the strong name exceed NAME_MAX while POSIX_NO_TRUNC is in effect.
ENOENT	The name of the semaphore does not exist.
ENOSPC	There is insufficient space for the creation of a new named semaphore.
ENOSYS	The function sem_unlink is not supported by this implementation.

DESCRIPTION:

The `sem_unlink()` function shall remove the semaphore name by the string name. If a process is currently accessing the name semaphore, the `sem_unlink` command has no effect. If one or more processes have the semaphore open when the `sem_unlink` function is called, the destruction of semaphores shall be postponed until all reference to semaphore are destroyed by calls to `sem_close`, `_exit()`, or `exec`. After all references have been destroyed, it returns immediately.

If the termination is successful, the function shall return 0. Otherwise, a -1 is returned and the `errno` is set.

NOTES:

9.4.6 `sem_wait` - Wait on a Semaphore

CALLING SEQUENCE:

```
int sem_wait(  
    sem_t *sem  
);
```

STATUS CODES:

EINVAL The "sem" argument does not refer to a valid semaphore

DESCRIPTION:

This function attempts to lock a semaphore specified by `sem`. If the semaphore is available, then the semaphore is locked (i.e., the semaphore value is decremented). If the semaphore is unavailable (i.e., the semaphore value is zero), then the function will block until the semaphore becomes available. It will then successfully lock the semaphore. The semaphore remains locked until released by a `sem_post()` call.

If the call is unsuccessful, then the function returns -1 and sets `errno` to the appropriate error code.

NOTES:

Multiprocessing is not supported in this implementation.

9.4.7 sem_trywait - Non-blocking Wait on a Semaphore

CALLING SEQUENCE:

```
int sem_trywait(  
    sem_t *sem  
);
```

STATUS CODES:

EAGAIN The semaphore is not available (i.e., the semaphore value is zero), so the semaphore could not be locked.

EINVAL The `sem` argument does not refer to a valid semaphore

DESCRIPTION:

This function attempts to lock a semaphore specified by `sem`. If the semaphore is available, then the semaphore is locked (i.e., the semaphore value is decremented) and the function returns a value of 0. The semaphore remains locked until released by a `sem_post()` call. If the semaphore is unavailable (i.e., the semaphore value is zero), then the function will return a value of -1 immediately and set `errno` to `EAGAIN`.

If the call is unsuccessful, then the function returns -1 and sets `errno` to the appropriate error code.

NOTES:

Multiprocessing is not supported in this implementation.

9.4.8 `sem_timedwait` - Wait on a Semaphore for a Specified Time

CALLING SEQUENCE:

```
int sem_timedwait(  
    sem_t          *sem,  
    const struct timespec *abstime  
);
```

STATUS CODES:

EAGAIN The semaphore is not available (i.e., the semaphore value is zero), so the semaphore could not be locked.

EINVAL The `sem` argument does not refer to a valid semaphore

DESCRIPTION:

This function attempts to lock a semaphore specified by `sem`, and will wait for the semaphore until the absolute time specified by `abstime`. If the semaphore is available, then the semaphore is locked (i.e., the semaphore value is decremented) and the function returns a value of 0. The semaphore remains locked until released by a `sem_post()` call. If the semaphore is unavailable, then the function will wait for the semaphore to become available for the amount of time specified by `timeout`.

If the semaphore does not become available within the interval specified by `timeout`, then the function returns -1 and sets `errno` to `EAGAIN`. If any other error occurs, the function returns -1 and sets `errno` to the appropriate error code.

NOTES:

Multiprocessing is not supported in this implementation.

9.4.9 sem_post - Unlock a Semaphore

CALLING SEQUENCE:

```
int sem_post(  
    sem_t *sem  
);
```

STATUS CODES:

EINVAL The `sem` argument does not refer to a valid semaphore

DESCRIPTION:

This function attempts to release the semaphore specified by `sem`. If other tasks are waiting on the semaphore, then one of those tasks (which one depends on the scheduler being used) is allowed to lock the semaphore and return from its `sem_wait()`, `sem_trywait()`, or `sem_timedwait()` call. If there are no other tasks waiting on the semaphore, then the semaphore value is simply incremented. `sem_post()` returns 0 upon successful completion.

If an error occurs, the function returns -1 and sets `errno` to the appropriate error code.

NOTES:

Multiprocessing is not supported in this implementation.

9.4.10 sem_getvalue - Get the value of a semaphore

CALLING SEQUENCE:

```
int sem_getvalue(  
    sem_t *sem,  
    int *sval  
);
```

STATUS CODES:

- EINVAL** The "sem" argument does not refer to a valid semaphore
- ENOSYS** The function sem_getvalue is not supported by this implementation

DESCRIPTION:

The sem_getvalue functions sets the location referenced by the "sval" argument to the value of the semaphore without affecting the state of the semaphore. The updated value represents a semaphore value that occurred at some point during the call, but is not necessarily the actual value of the semaphore when it returns to the calling process.

If "sem" is locked, the value returned by sem_getvalue will be zero or a negative number whose absolute value is the number of processes waiting for the semaphore at some point during the call.

NOTES:

If the functions completes successfully, it shall return a value of zero. Otherwise, it shall return a value of -1 and set "errno" to specify the error that occurred.

10 Mutex Manager

10.1 Introduction

The mutex manager implements the functionality required of the mutex manager as defined by POSIX 1003.1b-1996. This standard requires that a compliant operating system provide the facilities to ensure that threads can operate with mutual exclusion from one another and defines the API that must be provided.

The services provided by the mutex manager are:

- `pthread_mutexattr_init` - Initialize a Mutex Attribute Set
- `pthread_mutexattr_destroy` - Destroy a Mutex Attribute Set
- `pthread_mutexattr_setprotocol` - Set the Blocking Protocol
- `pthread_mutexattr_getprotocol` - Get the Blocking Protocol
- `pthread_mutexattr_setprioceiling` - Set the Priority Ceiling
- `pthread_mutexattr_getprioceiling` - Get the Priority Ceiling
- `pthread_mutexattr_setpshared` - Set the Visibility
- `pthread_mutexattr_getpshared` - Get the Visibility
- `pthread_mutex_init` - Initialize a Mutex
- `pthread_mutex_destroy` - Destroy a Mutex
- `pthread_mutex_lock` - Lock a Mutex
- `pthread_mutex_trylock` - Poll to Lock a Mutex
- `pthread_mutex_timedlock` - Lock a Mutex with Timeout
- `pthread_mutex_unlock` - Unlock a Mutex
- `pthread_mutex_setprioceiling` - Dynamically Set the Priority Ceiling
- `pthread_mutex_getprioceiling` - Dynamically Get the Priority Ceiling

10.2 Background

10.2.1 Mutex Attributes

Mutex attributes are utilized only at mutex creation time. A mutex attribute structure may be initialized and passed as an argument to the `mutex_init` routine. Note that the priority ceiling of a mutex may be set at run-time.

blocking protocol is the XXX

priority ceiling is the XXX

pshared is the XXX

10.2.2 PTHREAD_MUTEX_INITIALIZER

This is a special value that a variable of type `pthread_mutex_t` may be statically initialized to as shown below:

```
pthread_mutex_t my_mutex = PTHREAD_MUTEX_INITIALIZER;
```

This indicates that `my_mutex` will be automatically initialized by an implicit call to `pthread_mutex_init` the first time the mutex is used.

Note that the mutex will be initialized with default attributes.

10.3 Operations

There is currently no text in this section.

10.4 Services

This section details the mutex manager's services. A subsection is dedicated to each of this manager's services and describes the calling sequence, related constants, usage, and status codes.

10.4.1 pthread_mutexattr_init - Initialize a Mutex Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_init(
    pthread_mutexattr_t *attr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

DESCRIPTION:

The `pthread_mutexattr_init` routine initializes the mutex attributes object specified by `attr` with the default value for all of the individual attributes.

NOTES:

XXX insert list of default attributes here.

10.4.2 pthread_mutexattr_destroy - Destroy a Mutex Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_destroy(
    pthread_mutexattr_t *attr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.
EINVAL The attribute set is not initialized.

DESCRIPTION:

The `pthread_mutex_attr_destroy` routine is used to destroy a mutex attributes object. The behavior of using an attributes object after it is destroyed is implementation dependent.

NOTES:

NONE

10.4.3 pthread_mutexattr_setprotocol - Set the Blocking Protocol

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_setprotocol(
    pthread_mutexattr_t *attr,
    int                  protocol
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.
EINVAL The attribute set is not initialized.
EINVAL The protocol argument is invalid.

DESCRIPTION:

The `pthread_mutexattr_setprotocol` routine is used to set value of the `protocol` attribute. This attribute controls the order in which threads waiting on this mutex will receive it.

The `protocol` can be one of the following:

`PTHREAD_PRIO_NONE` in which case blocking order is FIFO.

`PTHREAD_PRIO_INHERIT`
in which case blocking order is priority with the priority inheritance protocol in effect.

`PTHREAD_PRIO_PROTECT`
in which case blocking order is priority with the priority ceiling protocol in effect.

NOTES:

There is currently no way to get simple priority blocking ordering with POSIX mutexes even though this could easily be supported by RTEMS.

10.4.4 pthread_mutexattr_getprotocol - Get the Blocking Protocol

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_getprotocol(
    pthread_mutexattr_t *attr,
    int *protocol
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The protocol pointer argument is invalid.

DESCRIPTION:

The `pthread_mutexattr_getprotocol` routine is used to obtain the value of the `protocol` attribute. This attribute controls the order in which threads waiting on this mutex will receive it.

NOTES:

NONE

10.4.5 pthread_mutexattr_setprioceiling - Set the Priority Ceiling

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_setprioceiling(
    pthread_mutexattr_t *attr,
    int prioceiling
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The prioceiling argument is invalid.

DESCRIPTION:

The `pthread_mutexattr_setprioceiling` routine is used to set value of the `prioceiling` attribute. This attribute specifies the priority that is the ceiling for threads obtaining this mutex. Any task obtaining this mutex may not be of greater priority than the ceiling. If it is of lower priority, then its priority will be elevated to `prioceiling`.

NOTES:

NONE

10.4.6 pthread_mutexattr_getprioceiling - Get the Priority Ceiling

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_getprioceiling(
    const pthread_mutexattr_t *attr,
    int *prioceiling
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The prioceiling pointer argument is invalid.

DESCRIPTION:

The `pthread_mutexattr_getprioceiling` routine is used to obtain the value of the `prioceiling` attribute. This attribute specifies the priority ceiling for this mutex.

NOTES:

NONE

10.4.7 pthread_mutexattr_setpshared - Set the Visibility

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_setpshared(
    pthread_mutexattr_t *attr,
    int pshared
);
```

STATUS CODES:

EINVAL	The attribute pointer argument is invalid.
EINVAL	The attribute set is not initialized.
EINVAL	The pshared argument is invalid.

DESCRIPTION:

NOTES:

10.4.8 pthread_mutexattr_getpshared - Get the Visibility

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutexattr_getpshared(
    const pthread_mutexattr_t *attr,
    int *pshared
);
```

STATUS CODES:

EINVAL	The attribute pointer argument is invalid.
EINVAL	The attribute set is not initialized.
EINVAL	The pshared pointer argument is invalid.

DESCRIPTION:

NOTES:

10.4.9 pthread_mutex_init - Initialize a Mutex

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_init(
    pthread_mutex_t      *mutex,
    const pthread_mutexattr_t *attr
);
```

STATUS CODES:

EINVAL	The attribute set is not initialized.
EINVAL	The specified protocol is invalid.
EAGAIN	The system lacked the necessary resources to initialize another mutex.
ENOMEM	Insufficient memory exists to initialize the mutex.
EBUSY	Attempted to reinitialize the object reference by mutex, a previously initialized, but not yet destroyed.

DESCRIPTION:

NOTES:

10.4.10 pthread_mutex_destroy - Destroy a Mutex

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_destroy(
    pthread_mutex_t *mutex
);
```

STATUS CODES:

EINVAL	The specified mutex is invalid.
EBUSY	Attempted to destroy the object reference by mutex, while it is locked or referenced by another thread.

DESCRIPTION:

NOTES:

10.4.11 pthread_mutex_lock - Lock a Mutex

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_lock(
    pthread_mutex_t *mutex
);
```

STATUS CODES:

EINVAL The specified mutex is invalid.

EINVAL The mutex has the protocol attribute of `PTHREAD_PRIO_PROTECT` and the priority of the calling thread is higher than the current priority ceiling.

EDEADLK The current thread already owns the mutex.

DESCRIPTION:

NOTES:

10.4.12 pthread_mutex_trylock - Poll to Lock a Mutex

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_trylock(
    pthread_mutex_t *mutex
);
```

STATUS CODES:

- EINVAL** The specified mutex is invalid.
- EINVAL** The mutex has the protocol attribute of `PTHREAD_PRIO_PROTECT` and the priority of the calling thread is higher than the current priority ceiling.
- EDEADLK** The current thread already owns the mutex.

DESCRIPTION:

NOTES:

10.4.13 pthread_mutex_timedlock - Lock a Mutex with Timeout

CALLING SEQUENCE:

```
#include <pthread.h>
#include <time.h>

int pthread_mutex_timedlock(
    pthread_mutex_t      *mutex,
    const struct timespec *timeout
);
```

STATUS CODES:

EINVAL	The specified mutex is invalid.
EINVAL	The nanoseconds field of timeout is invalid.
EINVAL	The mutex has the protocol attribute of PTHREAD_PRIO_PROTECT and the priority of the calling thread is higher than the current priority ceiling.
EDEADLK	The current thread already owns the mutex.
ETIMEDOUT	The calling thread was unable to obtain the mutex within the specified timeout period.

DESCRIPTION:

NOTES:

10.4.14 pthread_mutex_unlock - Unlock a Mutex

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_unlock(
    pthread_mutex_t *mutex
);
```

STATUS CODES:

EINVAL The specified mutex is invalid.

DESCRIPTION:

NOTES:

10.4.15 pthread_mutex_setprioceiling - Dynamically Set the Priority Ceiling

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_setprioceiling(
    pthread_mutex_t *mutex,
    int prioceiling,
    int *oldceiling
);
```

STATUS CODES:

EINVAL	The oldceiling pointer parameter is invalid.
EINVAL	The prioceiling parameter is an invalid priority.
EINVAL	The specified mutex is invalid.

DESCRIPTION:

NOTES:

10.4.16 pthread_mutex_getprioceiling - Get the Current Priority Ceiling

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_mutex_getprioceiling(
    pthread_mutex_t *mutex,
    int *prioceiling
);
```

STATUS CODES:

EINVAL The prioceiling pointer parameter is invalid.
EINVAL The specified mutex is invalid.

DESCRIPTION:

NOTES:

11 Condition Variable Manager

11.1 Introduction

The condition variable manager ...

The directives provided by the condition variable manager are:

- `pthread_condattr_init` - Initialize a Condition Variable Attribute Set
- `pthread_condattr_destroy` - Destroy a Condition Variable Attribute Set
- `pthread_condattr_setpshared` - Set Process Shared Attribute
- `pthread_condattr_getpshared` - Get Process Shared Attribute
- `pthread_cond_init` - Initialize a Condition Variable
- `pthread_cond_destroy` - Destroy a Condition Variable
- `pthread_cond_signal` - Signal a Condition Variable
- `pthread_cond_broadcast` - Broadcast a Condition Variable
- `pthread_cond_wait` - Wait on a Condition Variable
- `pthread_cond_timedwait` - With with Timeout a Condition Variable

11.2 Background

There is currently no text in this section.

11.3 Operations

There is currently no text in this section.

11.4 Directives

This section details the condition variable manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

11.4.1 pthread_condattr_init - Initialize a Condition Variable Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_condattr_init(
    pthread_condattr_t *attr
);
```

STATUS CODES:

ENOMEM Insufficient memory is available to initialize the condition variable attributes object.

DESCRIPTION:

NOTES:

11.4.2 pthread_condattr_destroy - Destroy a Condition Variable Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_condattr_destroy(
    pthread_condattr_t *attr
);
```

STATUS CODES:

EINVAL The attribute object specified is invalid.

DESCRIPTION:

NOTES:

11.4.3 pthread_condattr_setpshared - Set Process Shared Attribute

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_condattr_setpshared(
    pthread_condattr_t *attr,
    int pshared
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

NOTES:

11.4.4 pthread_condattr_getpshared - Get Process Shared Attribute

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_condattr_getpshared(
    const pthread_condattr_t *attr,
    int *pshared
);
```

STATUS CODES:

EINVAL Invalid argument passed.

DESCRIPTION:

NOTES:

11.4.5 pthread_cond_init - Initialize a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_init(
    pthread_cond_t      *cond,
    const pthread_condattr_t *attr
);
```

STATUS CODES:

EAGAIN	The system lacked a resource other than memory necessary to create the initialize the condition variable object.
ENOMEM	Insufficient memory is available to initialize the condition variable object.
EBUSY	The specified condition variable has already been initialized.
EINVAL	The specified attribute value is invalid.

DESCRIPTION:

NOTES:

11.4.6 pthread_cond_destroy - Destroy a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_destroy(
    pthread_cond_t *cond
);
```

STATUS CODES:

EINVAL The specified condition variable is invalid.
EBUSY The specified condition variable is currently in use.

DESCRIPTION:

NOTES:

11.4.7 pthread_cond_signal - Signal a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_signal(
    pthread_cond_t *cond
);
```

STATUS CODES:

EINVAL The specified condition variable is not valid.

DESCRIPTION:

NOTES:

This routine should not be invoked from a handler from an asynchronous signal handler or an interrupt service routine.

11.4.8 pthread_cond_broadcast - Broadcast a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_broadcast(
    pthread_cond_t *cond
);
```

STATUS CODES:

EINVAL The specified condition variable is not valid.

DESCRIPTION:

NOTES:

This routine should not be invoked from a handler from an asynchronous signal handler or an interrupt service routine.

11.4.9 pthread_cond_wait - Wait on a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_wait(
    pthread_cond_t *cond,
    pthread_mutex_t *mutex
);
```

STATUS CODES:

EINVAL The specified condition variable or mutex is not initialized OR different mutexes were specified for concurrent pthread_cond_wait() and pthread_cond_timedwait() operations on the same condition variable OR the mutex was not owned by the current thread at the time of the call.

DESCRIPTION:

NOTES:

11.4.10 pthread_cond_timedwait - Wait with Timeout a Condition Variable

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_cond_timedwait(
    pthread_cond_t      *cond,
    pthread_mutex_t     *mutex,
    const struct timespec *abstime
);
```

STATUS CODES:

- EINVAL** The specified condition variable or mutex is not initialized OR different mutexes were specified for concurrent pthread_cond_wait() and pthread_cond_timedwait() operations on the same condition variable OR the mutex was not owned by the current thread at the time of the call.
- ETIMEDOUT** The specified time has elapsed without the condition variable being satisfied.

DESCRIPTION:

NOTES:

12 Memory Management Manager

12.1 Introduction

The memory management manager is ...

The directives provided by the memory management manager are:

- `mlockall` - Lock the Address Space of a Process
- `munlockall` - Unlock the Address Space of a Process
- `mlock` - Lock a Range of the Process Address Space
- `munlock` - Unlock a Range of the Process Address Space
- `mmap` - Map Process Addresses to a Memory Object
- `munmap` - Unmap Previously Mapped Addresses
- `mprotect` - Change Memory Protection
- `msync` - Memory Object Synchronization
- `shm_open` - Open a Shared Memory Object
- `shm_unlink` - Remove a Shared Memory Object

12.2 Background

There is currently no text in this section.

12.3 Operations

There is currently no text in this section.

12.4 Directives

This section details the memory management manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

12.4.1 mlockall - Lock the Address Space of a Process

CALLING SEQUENCE:

```
int mlockall(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.2 munlockall - Unlock the Address Space of a Process

CALLING SEQUENCE:

```
int munlockall(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.3 mlock - Lock a Range of the Process Address Space

CALLING SEQUENCE:

```
int mlock(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.4 munlock - Unlock a Range of the Process Address Space

CALLING SEQUENCE:

```
int munlock(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.5 mmap - Map Process Addresses to a Memory Object

CALLING SEQUENCE:

```
int mmap(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.6 munmap - Unmap Previously Mapped Addresses

CALLING SEQUENCE:

```
int munmap(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.7 mprotect - Change Memory Protection

CALLING SEQUENCE:

```
int mprotect(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.8 msync - Memory Object Synchronization

CALLING SEQUENCE:

```
int msync(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.9 shm_open - Open a Shared Memory Object

CALLING SEQUENCE:

```
int shm_open(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

12.4.10 shm_unlink - Remove a Shared Memory Object

CALLING SEQUENCE:

```
int shm_unlink(  
);
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

13 Scheduler Manager

13.1 Introduction

The scheduler manager ...

The directives provided by the scheduler manager are:

- `sched_get_priority_min` - Get Minimum Priority Value
- `sched_get_priority_max` - Get Maximum Priority Value
- `sched_rr_get_interval` - Get Timeslicing Quantum
- `sched_yield` - Yield the Processor

13.2 Background

13.2.1 Priority

In the RTEMS implementation of the POSIX API, the priorities range from the low priority of `sched_get_priority_min()` to the highest priority of `sched_get_priority_max()`. Numerically higher values represent higher priorities.

13.2.2 Scheduling Policies

The following scheduling policies are available:

SCHED_FIFO	Priority-based, preemptive scheduling with no timeslicing. This is equivalent to what is called "manual round-robin" scheduling.
SCHED_RR	Priority-based, preemptive scheduling with timeslicing. Time quanta are maintained on a per-thread basis and are not reset at each context switch. Thus, a thread which is preempted and subsequently resumes execution will attempt to complete the unused portion of its time quantum.
SCHED_OTHER	Priority-based, preemptive scheduling with timeslicing. Time quanta are maintained on a per-thread basis and are reset at each context switch.
SCHED_SPORADIC	Priority-based, preemptive scheduling utilizing three additional parameters: budget, replenishment period, and low priority. Under this policy, the thread is allowed to execute for "budget" amount of time before its priority is lowered to "low priority". At the end of each replenishment period, the thread resumes its initial priority and has its budget replenished.

13.3 Operations

There is currently no text in this section.

13.4 Directives

This section details the scheduler manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

13.4.1 sched_get_priority_min - Get Minimum Priority Value

CALLING SEQUENCE:

```
#include <sched.h>

int sched_get_priority_min(
    int policy
);
```

STATUS CODES:

On error, this routine returns -1 and sets errno to one of the following:

EINVAL The indicated policy is invalid.

DESCRIPTION:

This routine return the minimum (numerically and logically lowest) priority for the specified policy.

NOTES:

NONE

13.4.2 sched_get_priority_max - Get Maximum Priority Value

CALLING SEQUENCE:

```
#include <sched.h>

int sched_get_priority_max(
    int policy
);
```

STATUS CODES:

On error, this routine returns -1 and sets errno to one of the following:

EINVAL The indicated policy is invalid.

DESCRIPTION:

This routine return the maximum (numerically and logically highest) priority for the specified policy.

NOTES:

NONE

13.4.3 sched_rr_get_interval - Get Timeslicing Quantum

CALLING SEQUENCE:

```
#include <sched.h>

int sched_rr_get_interval(
    pid_t          pid,
    struct timespec *interval
);
```

STATUS CODES:

On error, this routine returns -1 and sets `errno` to one of the following:

ESRCH	The indicated process id is invalid.
EINVAL	The specified interval pointer parameter is invalid.

DESCRIPTION:

This routine returns the length of the timeslice quantum in the `interval` parameter for the specified `pid`.

NOTES:

The `pid` argument should be 0 to indicate the calling process.

13.4.4 sched_yield - Yield the Processor

CALLING SEQUENCE:

```
#include <sched.h>
```

```
int sched_yield( void );
```

STATUS CODES:

This routine always returns zero to indicate success.

DESCRIPTION:

This call forces the calling thread to yield the processor to another thread. Normally this is used to implement voluntary round-robin task scheduling.

NOTES:

NONE

14 Clock Manager

14.1 Introduction

The clock manager provides services two primary classes of services. The first focuses on obtaining and setting the current date and time. The other category of services focus on allowing a thread to delay for a specific length of time.

The directives provided by the clock manager are:

- `clock_gettime` - Obtain Time of Day
- `clock_settime` - Set Time of Day
- `clock_getres` - Get Clock Resolution
- `sleep` - Delay Process Execution
- `usleep` - Delay Process Execution in Microseconds
- `nanosleep` - Delay with High Resolution
- `gettimeofday` - Get the Time of Day
- `time` - Get time in seconds

14.2 Background

There is currently no text in this section.

14.3 Operations

There is currently no text in this section.

14.4 Directives

This section details the clock manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

14.4.1 `clock_gettime` - Obtain Time of Day

CALLING SEQUENCE:

```
#include <time.h>

int clock_gettime(
    clockid_t      clock_id,
    struct timespec *tp
);
```

STATUS CODES:

On error, this routine returns -1 and sets `errno` to one of the following:

- | | |
|---------------|---|
| EINVAL | The <code>tp</code> pointer parameter is invalid. |
| EINVAL | The <code>clock_id</code> specified is invalid. |

DESCRIPTION:

NOTES:

NONE

14.4.2 clock_settime - Set Time of Day

CALLING SEQUENCE:

```
#include <time.h>

int clock_settime(
    clockid_t          clock_id,
    const struct timespec *tp
);
```

STATUS CODES:

On error, this routine returns -1 and sets errno to one of the following:

EINVAL	The tp pointer parameter is invalid.
EINVAL	The clock_id specified is invalid.
EINVAL	The contents of the tp structure are invalid.

DESCRIPTION:

NOTES:

NONE

14.4.3 clock_getres - Get Clock Resolution

CALLING SEQUENCE:

```
#include <time.h>

int clock_getres(
    clockid_t      clock_id,
    struct timespec *res
);
```

STATUS CODES:

On error, this routine returns -1 and sets errno to one of the following:

EINVAL The res pointer parameter is invalid.

EINVAL The clock_id specified is invalid.

DESCRIPTION:

NOTES:

If res is NULL, then the resolution is not returned.

14.4.4 sleep - Delay Process Execution

CALLING SEQUENCE:

```
#include <unistd.h>

unsigned int sleep(
    unsigned int seconds
);
```

STATUS CODES:

This routine returns the number of unslept seconds.

DESCRIPTION:

The `sleep()` function delays the calling thread by the specified number of `seconds`.

NOTES:

This call is interruptible by a signal.

14.4.5 `usleep` - Delay Process Execution in Microseconds

CALLING SEQUENCE:

```
#include <time.h>

useconds_t usleep(
    useconds_t useconds
);
```

STATUS CODES:

This routine returns the number of unslept seconds.

DESCRIPTION:

The `sleep()` function delays the calling thread by the specified number of `seconds`.

The `usleep()` function suspends the calling thread from execution until either the number of microseconds specified by the `useconds` argument has elapsed or a signal is delivered to the calling thread and its action is to invoke a signal-catching function or to terminate the process.

Because of other activity, or because of the time spent in processing the call, the actual length of time the thread is blocked may be longer than the amount of time specified.

NOTES:

This call is interruptible by a signal.

The Single UNIX Specification allows this service to be implemented using the same timer as that used by the `alarm()` service. This is **NOT** the case for **RTEMS** and this call has no interaction with the `SIGALRM` signal.

14.4.6 nanosleep - Delay with High Resolution

CALLING SEQUENCE:

```
#include <time.h>

int nanosleep(
    const struct timespec *rqtp,
    struct timespec      *rmtp
);
```

STATUS CODES:

On error, this routine returns -1 and sets errno to one of the following:

EINTR	The routine was interrupted by a signal.
EAGAIN	The requested sleep period specified negative seconds or nanoseconds.
EINVAL	The requested sleep period specified an invalid number for the nanoseconds field.

DESCRIPTION:

NOTES:

This call is interruptible by a signal.

14.4.7 gettimeofday - Get the Time of Day

CALLING SEQUENCE:

```
#include <sys/time.h>
#include <unistd.h>

int gettimeofday(
    struct timeval *tp,
    struct timezone *tzp
);
```

STATUS CODES:

On error, this routine returns -1 and sets **errno** as appropriate.

EPerm	settimeofday is called by someone other than the superuser.
EINVAL	Timezone (or something else) is invalid.
EFAULT	One of tv or tz pointed outside your accessible address space

DESCRIPTION:

This routine returns the current time of day in the **tp** structure.

NOTES:

Currently, the timezone information is not supported. The **tzp** argument is ignored.

14.4.8 time - Get time in seconds

CALLING SEQUENCE:

```
#include <time.h>

int time(
    time_t *tloc
);
```

STATUS CODES:

This routine returns the number of seconds since the Epoch.

DESCRIPTION:

`time` returns the time since 00:00:00 GMT, January 1, 1970, measured in seconds. If `tloc` is non null, the return value is also stored in the memory pointed to by `t`.

NOTES:

NONE

15 Timer Manager

15.1 Introduction

The timer manager is ...

The services provided by the timer manager are:

- `timer_create` - Create a Per-Process Timer
- `timer_delete` - Delete a Per-Process Timer
- `timer_settime` - Set Next Timer Expiration
- `timer_gettime` - Get Time Remaining on Timer
- `timer_getoverrun` - Get Timer Overrun Count

15.2 Background

15.3 Operations

15.4 System Calls

This section details the timer manager's services. A subsection is dedicated to each of this manager's services and describes the calling sequence, related constants, usage, and status codes.

15.4.1 timer_create - Create a Per-Process Timer

CALLING SEQUENCE:

```
#include <time.h>
#include <signal.h>

int timer_create(
    clockid_t      clock_id,
    struct sigevent *evp,
    timer_t        *timerid
);
```

STATUS CODES:

EXXX -

DESCRIPTION:

NOTES:

15.4.2 timer_delete - Delete a Per-Process Timer

CALLING SEQUENCE:

```
#include <time.h>

int timer_delete(
    timer_t timerid
);
```

STATUS CODES:

EINVAL -

DESCRIPTION:

NOTES:

15.4.3 timer_settime - Set Next Timer Expiration

CALLING SEQUENCE:

```
#include <time.h>

int timer_settime(
    timer_t          timerid,
    int              flags,
    const struct itimerspec *value,
    struct itimerspec *ovalue
);
```

STATUS CODES:

EXXX -

DESCRIPTION:

NOTES:

15.4.4 timer_gettime - Get Time Remaining on Timer

CALLING SEQUENCE:

```
#include <time.h>

int timer_gettime(
    timer_t          timerid,
    struct itimerspec *value
);
```

STATUS CODES:

EXXX -

DESCRIPTION:

NOTES:

15.4.5 timer_getoverrun - Get Timer Overrun Count

CALLING SEQUENCE:

```
#include <time.h>

int timer_getoverrun(
    timer_t timerid
);
```

STATUS CODES:

EINVAL -

DESCRIPTION:

NOTES:

16 Message Passing Manager

16.1 Introduction

The message passing manager is the means to provide communication and synchronization capabilities using POSIX message queues.

The directives provided by the message passing manager are:

- `mq_open` - Open a Message Queue
- `mq_close` - Close a Message Queue
- `mq_unlink` - Remove a Message Queue
- `mq_send` - Send a Message to a Message Queue
- `mq_receive` - Receive a Message from a Message Queue
- `mq_notify` - Notify Process that a Message is Available
- `mq_setattr` - Set Message Queue Attributes
- `mq_getattr` - Get Message Queue Attributes

16.2 Background

16.2.1 Theory

Message queues are named objects that operate with readers and writers. In addition, a message queue is a priority queue of discrete messages. POSIX message queues offer a certain, basic amount of application access to, and control over, the message queue geometry that can be changed.

16.2.2 Messages

A message is a variable length buffer where information can be stored to support communication. The length of the message and the information stored in that message are user-defined and can be actual data, pointer(s), or empty. There is a maximum acceptable length for a message that is associated with each message queue.

16.2.3 Message Queues

Message queues are named objects similar to the pipes of POSIX. They are a means of communicating data between multiple processes and for passing messages among tasks and ISRs. Message queues can contain a variable number of messages from 0 to an upper limit that is user defined. The maximum length of the message can be set on a per message queue basis. Normally messages are sent and received from the message queue in FIFO order. However, messages can also be prioritized and a priority queue established for the passing of messages. Synchronization is needed when a task waits for a message to arrive at a queue. Also, a task may poll a queue for the arrival of a message.

The message queue descriptor `mqd_t` represents the message queue. It is passed as an argument to all of the message queue functions.

16.2.4 Building a Message Queue Attribute Set

The `mq_attr` structure is used to define the characteristics of the message queue.

```
typedef struct mq_attr{
    long mq_flags;
    long mq_maxmsg;
    long mq_msgsize;
    long mq_curmsgs;
};
```

All of these attributes are set when the message queue is created using `mq_open`. The `mq_flags` field is not used in the creation of a message queue, it is only used by `mq_setattr` and `mq_getattr`. The structure `mq_attr` is passed as an argument to `mq_setattr` and `mq_getattr`.

The `mq_flags` contain information affecting the behavior of the message queue. The `O_NONBLOCK` `mq_flag` is the only flag that is defined. In `mq_setattr`, the `mq_flag` can be set to dynamically change the blocking and non-blocking behavior of the message queue. If the non-block flag is set then the message queue is non-blocking, and requests to send and receive messages do not block waiting for resources. For a blocking message queue, a request to send might have to wait for an empty message queue, and a request to receive might have to wait for a message to arrive on the queue. Both `mq_maxmsg` and `mq_msgsize` affect the sizing of the message queue. `mq_maxmsg` specifies how many messages the queue can hold at any one time. `mq_msgsize` specifies the size of any one message on the queue. If either of these limits is exceeded, an error message results.

Upon return from `mq_getattr`, the `mq_curmsgs` is set according to the current state of the message queue. This specifies the number of messages currently on the queue.

16.2.5 Notification of a Message on the Queue

Every message queue has the ability to notify one (and only one) process whenever the queue's state changes from empty (0 messages) to nonempty. This means that the process does not have to block or constantly poll while it waits for a message. By calling `mq_notify`, you can attach a notification request to a message queue. When a message is received by an empty queue, if there are no processes blocked and waiting for the message, then the queue notifies the requesting process of a message arrival. There is only one signal sent by the message queue, after that the notification request is de-registered and another process can attach its notification request. After receipt of a notification, a process must re-register if it wishes to be notified again.

If there is a process blocked and waiting for the message, that process gets the message, and notification is not sent. It is also possible for another process to receive the message after the notification is sent but before the notified process has sent its receive request.

Only one process can have a notification request attached to a message queue at any one time. If another process attempts to register a notification request, it fails. You can de-register for a message queue by passing a `NULL` to `mq_notify`, this removes any notification request attached to the queue. Whenever the message queue is closed, all notification attachments are removed.

16.2.6 POSIX Interpretation Issues

There is one significant point of interpretation related to the RTEMS implementation of POSIX message queues:

What happens to threads already blocked on a message queue when the mode of that same message queue is changed from blocking to non-blocking?

The RTEMS POSIX implementation decided to unblock all waiting tasks with an **EAGAIN** status just as if a non-blocking version of the same operation had returned unsatisfied. This case is not discussed in the POSIX standard and other implementations may have chosen alternative behaviors.

16.3 Operations

16.3.1 Opening or Creating a Message Queue

If the message queue already exists, `mq_open()` opens it, if the message queue does not exist, `mq_open()` creates it. When a message queue is created, the geometry of the message queue is contained in the attribute structure that is passed in as an argument. This includes `mq_msgsize` that dictates the maximum size of a single message, and the `mq_maxmsg` that dictates the maximum number of messages the queue can hold at one time. The blocking or non-blocking behavior of the queue can also be specified.

16.3.2 Closing a Message Queue

The `mq_close()` function is used to close the connection made to a message queue that was made during `mq_open`. The message queue itself and the messages on the queue are persistent and remain after the queue is closed.

16.3.3 Removing a Message Queue

The `mq_unlink()` function removes the named message queue. If the message queue is not open when `mq_unlink` is called, then the queue is immediately eliminated. Any messages that were on the queue are lost, and the queue can not be opened again. If processes have the queue open when `mq_unlink` is called, the removal of the queue is delayed until the last process using the queue has finished. However, the name of the message queue is removed so that no other process can open it.

16.3.4 Sending a Message to a Message Queue

The `mq_send()` function adds the message in priority order to the message queue. Each message has an assigned priority. The highest priority message is at the front of the queue.

The maximum number of messages that a message queue may accept is specified at creation by the `mq_maxmsg` field of the attribute structure. If this amount is exceeded, the behavior of the process is determined according to what `oflag` was used when the message queue was opened. If the queue was opened with `O_NONBLOCK` flag set, the process does not block, and an error is returned. If the `O_NONBLOCK` flag was not set, the process does block and wait for space on the queue.

16.3.5 Receiving a Message from a Message Queue

The `mq_receive()` function is used to receive the oldest of the highest priority message(s) from the message queue specified by `mqdes`. The messages are received in FIFO order within the priorities. The received message's priority is stored in the location referenced by the `msg_prio`. If the `msg_prio` is a NULL, the priority is discarded. The message is removed and stored in an area pointed to by `msg_ptr` whose length is of `msg_len`. The `msg_len` must be at least equal to the `mq_msgsize` attribute of the message queue.

The blocking behavior of the message queue is set by `O_NONBLOCK` at `mq_open` or by setting `O_NONBLOCK` in `mq_flags` in a call to `mq_setattr`. If this is a blocking queue, the process does block and wait on an empty queue. If this a non-blocking queue, the process does not block. Upon successful completion, `mq_receive` returns the length of the selected message in bytes and the message is removed from the queue.

16.3.6 Notification of Receipt of a Message on an Empty Queue

The `mq_notify()` function registers the calling process to be notified of message arrival at an empty message queue. Every message queue has the ability to notify one (and only one) process whenever the queue's state changes from empty (0 messages) to nonempty. This means that the process does not have to block or constantly poll while it waits for a message. By calling `mq_notify`, a notification request is attached to a message queue. When a message is received by an empty queue, if there are no processes blocked and waiting for the message, then the queue notifies the requesting process of a message arrival. There is only one signal sent by the message queue, after that the notification request is de-registered and another process can attach its notification request. After receipt of a notification, a process must re-register if it wishes to be notified again.

If there is a process blocked and waiting for the message, that process gets the message, and notification is not sent. Only one process can have a notification request attached to a message queue at any one time. If another process attempts to register a notification request, it fails. You can de-register for a message queue by passing a NULL to `mq_notify`, this removes any notification request attached to the queue. Whenever the message queue is closed, all notification attachments are removed.

16.3.7 Setting the Attributes of a Message Queue

The `mq_setattr()` function is used to set attributes associated with the open message queue description referenced by the message queue descriptor specified by `mqdes`. The `*omqstat` represents the old or previous attributes. If `omqstat` is non-NULL, the function `mq_setattr()` stores, in the location referenced by `omqstat`, the previous message queue attributes and the current queue status. These values are the same as would be returned by a call to `mq_getattr()` at that point.

There is only one `mq_attr.mq_flag` that can be altered by this call. This is the flag that deals with the blocking and non-blocking behavior of the message queue. If the flag is set then the message queue is non-blocking, and requests to send or receive do not block while waiting for resources. If the flag is not set, then message send and receive may involve waiting for an empty queue or waiting for a message to arrive.

16.3.8 Getting the Attributes of a Message Queue

The `mq_getattr()` function is used to get status information and attributes of the message queue associated with the message queue descriptor. The results are returned in the `mq_attr` structure referenced by the `mqstat` argument. All of these attributes are set at create time, except the blocking/non-blocking behavior of the message queue which can be dynamically set by using `mq_setattr`. The attribute `mq_curmsg` is set to reflect the number of messages on the queue at the time that `mq_getattr` was called.

16.4 Directives

This section details the message passing manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

16.4.1 mq_open - Open a Message Queue

CALLING SEQUENCE:

```
#include <mqqueue.h>

mqd_t mq_open(
    const char    *name,
    int           oflag,
    mode_t        mode,
    struct mq_attr *attr
);
```

STATUS CODES:

EACCES - Either the message queue exists and the permissions requested in oflags were denied, or the message does not exist and permission to create one is denied.

EEXIST - You tried to create a message queue that already exists.

EINVAL - An inappropriate name was given for the message queue, or the values of mq_maxmsg or mq_msgsize were less than 0.

ENOENT - The message queue does not exist, and you did not specify to create it.

EINTR - The call to mq_open was interrupted by a signal.

EMFILE - The process has too many files or message queues open. This is a process limit error.

ENFILE - The system has run out of resources to support more open message queues. This is a system error.

ENAMETOOLONG - mq_name is too long.

DESCRIPTION:

The mq_open () function establishes the connection between a process and a message queue with a message queue descriptor. If the message queue already exists, mq_open opens it, if the message queue does not exist, mq_open creates it. Message queues can have multiple senders and receivers. If mq_open is successful, the function returns a message queue descriptor. Otherwise, the function returns a -1 and sets 'errno' to indicate the error.

The name of the message queue is used as an argument. For the best of portability, the name of the message queue should begin with a "/" and no other "/" should be in the name. Different systems interpret the name in different ways.

The oflags contain information on how the message is opened if the queue already exists. This may be O_RDONLY for read only, O_WRONLY for write only, of O_RDWR, for read and write.

In addition, the oflags contain information needed in the creation of a message queue. O_NONBLOCK - If the non-block flag is set then the message queue is non-blocking, and requests to send and receive messages do not block waiting for resources. If the flag is not set then the message queue is blocking, and a request to send might have to wait for an empty

message queue. Similarly, a request to receive might have to wait for a message to arrive on the queue. **O_CREAT** - This call specifies that the call the `mq_open` is to create a new message queue. In this case the mode and attribute arguments of the function call are utilized. The message queue is created with a mode similar to the creation of a file, read and write permission creator, group, and others.

The geometry of the message queue is contained in the attribute structure. This includes `mq_msgsize` that dictates the maximum size of a single message, and the `mq_maxmsg` that dictates the maximum number of messages the queue can hold at one time. If a `NULL` is used in the `mq_attr` argument, then the message queue is created with implementation defined defaults. **O_EXCL** - is always set if **O_CREAT** flag is set. If the message queue already exists, **O_EXCL** causes an error message to be returned, otherwise, the new message queue fails and appends to the existing one.

NOTES:

The `mq_open ()` function does not add or remove messages from the queue. When a new message queue is being created, the `mq_flag` field of the attribute structure is not used.

16.4.2 mq_close - Close a Message Queue

CALLING SEQUENCE:

```
#include <mqqueue.h>

int mq_close(
    mqd_t mqdes
);
```

STATUS CODES:

EINVAL - The descriptor does not represent a valid open message queue

DESCRIPTION:

The `mq_close` function removes the association between the message queue descriptor, `mqdes`, and its message queue. If `mq_close()` is successfully completed, the function returns a value of zero; otherwise, the function returns a value of -1 and sets `errno` to indicate the error.

NOTES:

If the process had successfully attached a notification request to the message queue via `mq_notify`, this attachment is removed, and the message queue is available for another process to attach for notification. `mq_close` has no effect on the contents of the message queue, all the messages that were in the queue remain in the queue.

16.4.3 mq_unlink - Remove a Message Queue

CALLING SEQUENCE:

```
#include <mqqueue.h>

int mq_unlink(
    const char *name
);
```

STATUS CODES:

EINVAL - The descriptor does not represent a valid message queue

DESCRIPTION:

The `mq_unlink()` function removes the named message queue. If the message queue is not open when `mq_unlink` is called, then the queue is immediately eliminated. Any messages that were on the queue are lost, and the queue can not be opened again. If processes have the queue open when `mq_unlink` is called, the removal of the queue is delayed until the last process using the queue has finished. However, the name of the message queue is removed so that no other process can open it. Upon successful completion, the function returns a value of zero. Otherwise, the named message queue is not changed by this function call, and the function returns a value of -1 and sets `errno` to indicate the error.

NOTES:

Calls to `mq_open()` to re-create the message queue may fail until the message queue is actually removed. However, the `mq_unlink()` call need not block until all references have been closed; it may return immediately.

16.4.4 mq_send - Send a Message to a Message Queue

CALLING SEQUENCE:

```
#include<mqqueue.h>
int mq_send(
    mqd_t      mqdes,
    const char *msg_ptr,
    size_t     msg_len,
    unsigned int msg_prio
);
```

STATUS CODES:

EBADF - The descriptor does not represent a valid message queue, or the queue was opened for read only **O_RDONLY** **EINVAL** - The value of `msg_prio` was greater than the **MQ_PRIO_MAX**. **EMSGSIZE** - The `msg_len` is greater than the `mq_msgsize` attribute of the message queue **EAGAIN** - The message queue is non-blocking, and there is no room on the queue for another message as specified by the `mq_maxmsg`. **EINTR** - The message queue is blocking. While the process was waiting for free space on the queue, a signal arrived that interrupted the wait.

DESCRIPTION:

The `mq_send()` function adds the message pointed to by the argument `msg_ptr` to the message queue specified by `mqdes`. Each message is assigned a priority, from 0 to **MQ_PRIO_MAX**. **MQ_PRIO_MAX** is defined in `<limits.h>` and must be at least 32. Messages are added to the queue in order of their priority. The highest priority message is at the front of the queue.

The maximum number of messages that a message queue may accept is specified at creation by the `mq_maxmsg` field of the attribute structure. If this amount is exceeded, the behavior of the process is determined according to what `oflag` was used when the message queue was opened. If the queue was opened with **O_NONBLOCK** flag set, then the **EAGAIN** error is returned. If the **O_NONBLOCK** flag was not set, the process blocks and waits for space on the queue, unless it is interrupted by a signal.

Upon successful completion, the `mq_send()` function returns a value of zero. Otherwise, no message is enqueued, the function returns -1, and `errno` is set to indicate the error.

NOTES:

If the specified message queue is not full, `mq_send` inserts the message at the position indicated by the `msg_prio` argument.

16.4.5 mq_receive - Receive a Message from a Message Queue

CALLING SEQUENCE:

```
#include <mqqueue.h>

size_t mq_receive(
    mqd_t      mqdes,
    char       *msg_ptr,
    size_t     msg_len,
    unsigned int *msg_prio
);
```

STATUS CODES:

EBADF - The descriptor does not represent a valid message queue, or the queue was opened for write only **O_WRONLY** **EMSGSIZE** - The msg_len is less than the mq_msgsize attribute of the message queue **EAGAIN** - The message queue is non-blocking, and the queue is empty **EINTR** - The message queue is blocking. While the process was waiting for a message to arrive on the queue, a signal arrived that interrupted the wait.

DESCRIPTION:

The mq_receive function is used to receive the oldest of the highest priority message(s) from the message queue specified by mqdes. The messages are received in FIFO order within the priorities. The received message's priority is stored in the location referenced by the msg_prio. If the msg_prio is a NULL, the priority is discarded. The message is removed and stored in an area pointed to by msg_ptr whose length is of msg_len. The msg_len must be at least equal to the mq_msgsize attribute of the message queue.

The blocking behavior of the message queue is set by **O_NONBLOCK** at mq_open or by setting **O_NONBLOCK** in mq_flags in a call to mq_setattr. If this is a blocking queue, the process blocks and waits on an empty queue. If this a non-blocking queue, the process does not block.

Upon successful completion, mq_receive returns the length of the selected message in bytes and the message is removed from the queue. Otherwise, no message is removed from the queue, the function returns a value of -1, and sets errno to indicate the error.

NOTES:

If the size of the buffer in bytes, specified by the msg_len argument, is less than the mq_msgsize attribute of the message queue, the function fails and returns an error

16.4.6 mq_notify - Notify Process that a Message is Available

CALLING SEQUENCE:

```
#include <mqqueue.h>

int mq_notify(
    mqd_t          mqdes,
    const struct sigevent *notification
);
```

STATUS CODES:

EBADF - The descriptor does not refer to a valid message queue
EBUSY - A notification request is already attached to the queue

DESCRIPTION:

If the argument notification is not NULL, this function registers the calling process to be notified of message arrival at an empty message queue associated with the specified message queue descriptor, mqdes.

Every message queue has the ability to notify one (and only one) process whenever the queue's state changes from empty (0 messages) to nonempty. This means that the process does not have to block or constantly poll while it waits for a message. By calling mq_notify, a notification request is attached to a message queue. When a message is received by an empty queue, if there are no processes blocked and waiting for the message, then the queue notifies the requesting process of a message arrival. There is only one signal sent by the message queue, after that the notification request is de-registered and another process can attach its notification request. After receipt of a notification, a process must re-register if it wishes to be notified again.

If there is a process blocked and waiting for the message, that process gets the message, and notification is not be sent. Only one process can have a notification request attached to a message queue at any one time. If another process attempts to register a notification request, it fails. You can de-register for a message queue by passing a NULL to mq_notify; this removes any notification request attached to the queue. Whenever the message queue is closed, all notification attachments are removed.

Upon successful completion, mq_notify returns a value of zero; otherwise, the function returns a value of -1 and sets errno to indicate the error.

NOTES:

It is possible for another process to receive the message after the notification is sent but before the notified process has sent its receive request.

16.4.7 mq_setattr - Set Message Queue Attributes

CALLING SEQUENCE:

```
#include <mqqueue.h>

int mq_setattr(
    mqd_t          mqdes,
    const struct mq_attr *mqstat,
    struct mq_attr  *omqstat
);
```

STATUS CODES:

EBADF - The message queue descriptor does not refer to a valid, open queue. **EINVAL** - The `mq_flag` value is invalid.

DESCRIPTION:

The `mq_setattr` function is used to set attributes associated with the open message queue description referenced by the message queue descriptor specified by `mqdes`. The `*omqstat` represents the old or previous attributes. If `omqstat` is non-NULL, the function `mq_setattr()` stores, in the location referenced by `omqstat`, the previous message queue attributes and the current queue status. These values are the same as would be returned by a call to `mq_getattr()` at that point.

There is only one `mq_attr.mq_flag` which can be altered by this call. This is the flag that deals with the blocking and non-blocking behavior of the message queue. If the flag is set then the message queue is non-blocking, and requests to send or receive do not block while waiting for resources. If the flag is not set, then message send and receive may involve waiting for an empty queue or waiting for a message to arrive.

Upon successful completion, the function returns a value of zero and the attributes of the message queue have been changed as specified. Otherwise, the message queue attributes is unchanged, and the function returns a value of -1 and sets `errno` to indicate the error.

NOTES:

All other fields in the `mq_attr` are ignored by this call.

16.4.8 mq_getattr - Get Message Queue Attributes

CALLING SEQUENCE:

```
#include <mqqueue.h>
int mq_getattr(
    mqd_t mqdes,
    struct mq_attr *mqstat
);
```

STATUS CODES:

EBADF - The message queue descriptor does not refer to a valid, open message queue.

DESCRIPTION:

The mqdes argument specifies a message queue descriptor. The mq_getattr function is used to get status information and attributes of the message queue associated with the message queue descriptor. The results are returned in the mq_attr structure referenced by the mqstat argument. All of these attributes are set at create time, except the blocking/non-blocking behavior of the message queue which can be dynamically set by using mq_setattr. The attribute mq_curmsg is set to reflect the number of messages on the queue at the time that mq_getattr was called.

Upon successful completion, the mq_getattr function returns zero. Otherwise, the function returns -1 and sets errno to indicate the error.

NOTES:

17 Thread Manager

17.1 Introduction

The thread manager implements the functionality required of the thread manager as defined by POSIX 1003.1b-1996. This standard requires that a compliant operating system provide the facilities to manage multiple threads of control and defines the API that must be provided.

The services provided by the thread manager are:

- `pthread_attr_init` - Initialize a Thread Attribute Set
- `pthread_attr_destroy` - Destroy a Thread Attribute Set
- `pthread_attr_setdetachstate` - Set Detach State
- `pthread_attr_getdetachstate` - Get Detach State
- `pthread_attr_setstacksize` - Set Thread Stack Size
- `pthread_attr_getstacksize` - Get Thread Stack Size
- `pthread_attr_setstackaddr` - Set Thread Stack Address
- `pthread_attr_getstackaddr` - Get Thread Stack Address
- `pthread_attr_setscope` - Set Thread Scheduling Scope
- `pthread_attr_getscope` - Get Thread Scheduling Scope
- `pthread_attr_setinheritsched` - Set Inherit Scheduler Flag
- `pthread_attr_getinheritsched` - Get Inherit Scheduler Flag
- `pthread_attr_setschedpolicy` - Set Scheduling Policy
- `pthread_attr_getschedpolicy` - Get Scheduling Policy
- `pthread_attr_setschedparam` - Set Scheduling Parameters
- `pthread_attr_getschedparam` - Get Scheduling Parameters
- `pthread_create` - Create a Thread
- `pthread_exit` - Terminate the Current Thread
- `pthread_detach` - Detach a Thread
- `pthread_join` - Wait for Thread Termination
- `pthread_self` - Get Thread ID
- `pthread_equal` - Compare Thread IDs
- `pthread_once` - Dynamic Package Initialization
- `pthread_setschedparam` - Set Thread Scheduling Parameters
- `pthread_getschedparam` - Get Thread Scheduling Parameters

17.2 Background

17.2.1 Thread Attributes

Thread attributes are utilized only at thread creation time. A thread attribute structure may be initialized and passed as an argument to the `pthread_create` routine.

stack address is the address of the optionally user specified stack area for this thread. If this value is `NULL`, then RTEMS allocates the memory for the thread stack from the RTEMS Workspace Area. Otherwise, this is the user specified address for the memory to be used for the thread's stack. Each thread must have a distinct stack area. Each processor family has different alignment rules which should be followed.

stack size is the minimum desired size for this thread's stack area. If the size of this area as specified by the stack size attribute is smaller than the minimum for this processor family and the stack is not user specified, then RTEMS will automatically allocate a stack of the minimum size for this processor family.

contention scope specifies the scheduling contention scope. RTEMS only supports the `PTHREAD_SCOPE_PROCESS` scheduling contention scope.

scheduling inheritance specifies whether a user specified or the scheduling policy and parameters of the currently executing thread are to be used. When this is `PTHREAD_INHERIT_SCHED`, then the scheduling policy and parameters of the currently executing thread are inherited by the newly created thread.

scheduling policy and parameters specify the manner in which the thread will contend for the processor. The scheduling parameters are interpreted based on the specified policy. All policies utilize the thread priority parameter.

17.3 Operations

There is currently no text in this section.

17.4 Services

This section details the thread manager's services. A subsection is dedicated to each of this manager's services and describes the calling sequence, related constants, usage, and status codes.

17.4.1 pthread_attr_init - Initialize a Thread Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_init(
    pthread_attr_t *attr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_init` routine initializes the thread attributes object specified by `attr` with the default value for all of the individual attributes.

NOTES:

The settings in the default attributes are implementation defined. For RTEMS, the default attributes are as follows:

- `stackaddr` is not set to indicate that RTEMS is to allocate the stack memory.
- `stacksize` is set to `PTHREAD_MINIMUM_STACK_SIZE`.
- `contentionscope` is set to `PTHREAD_SCOPE_PROCESS`.
- `inheritsched` is set to `PTHREAD_INHERIT_SCHED` to indicate that the created thread inherits its scheduling attributes from its parent.
- `detachstate` is set to `PTHREAD_CREATE_JOINABLE`.

17.4.2 pthread_attr_destroy - Destroy a Thread Attribute Set

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_destroy(
    pthread_attr_t *attr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.
EINVAL The attribute set is not initialized.

DESCRIPTION:

The `pthread_attr_destroy` routine is used to destroy a thread attributes object. The behavior of using an attributes object after it is destroyed is implementation dependent.

NOTES:

NONE

17.4.3 pthread_attr_setdetachstate - Set Detach State

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setdetachstate(
    pthread_attr_t *attr,
    int             detachstate
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The detachstate argument is invalid.

DESCRIPTION:

The `pthread_attr_setdetachstate` routine is used to value of the `detachstate` attribute. This attribute controls whether the thread is created in a detached state.

The `detachstate` can be either `PTHREAD_CREATE_DETACHED` or `PTHREAD_CREATE_JOINABLE`. The default value for all threads is `PTHREAD_CREATE_JOINABLE`.

NOTES:

If a thread is in a detached state, then the use of the ID with the `pthread_detach` or `pthread_join` routines is an error.

17.4.4 pthread_attr_getdetachstate - Get Detach State

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getdetachstate(
    const pthread_attr_t *attr,
    int *detachstate
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The detachstate pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_getdetachstate` routine is used to obtain the current value of the `detachstate` attribute as specified by the `attr` thread attribute object.

NOTES:

NONE

17.4.5 pthread_attr_setstacksize - Set Thread Stack Size

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setstacksize(
    pthread_attr_t *attr,
    size_t          stacksize
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.
EINVAL The attribute set is not initialized.

DESCRIPTION:

The `pthread_attr_setstacksize` routine is used to set the `stacksize` attribute in the `attr` thread attribute object.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_ATTR_STACKSIZE` to indicate that this routine is supported.

If the specified `stacksize` is below the minimum required for this CPU (`PTHREAD_STACK_MIN`), then the `stacksize` will be set to the minimum for this CPU.

17.4.6 pthread_attr_getstacksize - Get Thread Stack Size

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getstacksize(
    const pthread_attr_t *attr,
    size_t                *stacksize
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The stacksize pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_getstacksize` routine is used to obtain the `stacksize` attribute in the `attr` thread attribute object.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_ATTR_STACKSIZE` to indicate that this routine is supported.

17.4.7 pthread_attr_setstackaddr - Set Thread Stack Address

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setstackaddr(
    pthread_attr_t *attr,
    void          *stackaddr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.
EINVAL The attribute set is not initialized.

DESCRIPTION:

The `pthread_attr_setstackaddr` routine is used to set the `stackaddr` attribute in the `attr` thread attribute object.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_ATTR_STACKADDR` to indicate that this routine is supported.

It is imperative to the proper operation of the system that each thread have sufficient stack space.

17.4.8 pthread_attr_getstackaddr - Get Thread Stack Address

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getstackaddr(
    const pthread_attr_t *attr,
    void **stackaddr
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The stackaddr pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_getstackaddr` routine is used to obtain the `stackaddr` attribute in the `attr` thread attribute object.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_ATTR_STACKADDR` to indicate that this routine is supported.

17.4.9 pthread_attr_setscope - Set Thread Scheduling Scope

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setscope(
    pthread_attr_t *attr,
    int             contentionscope
);
```

STATUS CODES:

EINVAL	The attribute pointer argument is invalid.
EINVAL	The attribute set is not initialized.
EINVAL	The contention scope specified is not valid.
ENOTSUP	The contention scope specified (PTHREAD_SCOPE_SYSTEM) is not supported.

DESCRIPTION:

The `pthread_attr_setscope` routine is used to set the contention scope field in the thread attribute object `attr` to the value specified by `contentionscope`.

The `contentionscope` must be either `PTHREAD_SCOPE_SYSTEM` to indicate that the thread is to be within system scheduling contention or `PTHREAD_SCOPE_PROCESS` indicating that the thread is to be within the process scheduling contention scope.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.10 pthread_attr_getscope - Get Thread Scheduling Scope

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getscope(
    const pthread_attr_t *attr,
    int *contentionscope
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The contentionscope pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_getscope` routine is used to obtain the value of the contention scope field in the thread attributes object `attr`. The current value is returned in `contentionscope`.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.11 pthread_attr_setinheritsched - Set Inherit Scheduler Flag

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setinheritsched(
    pthread_attr_t *attr,
    int             inheritsched
);
```

STATUS CODES:

- EINVAL** The attribute pointer argument is invalid.
- EINVAL** The attribute set is not initialized.
- EINVAL** The specified scheduler inheritance argument is invalid.

DESCRIPTION:

The `pthread_attr_setinheritsched` routine is used to set the inherit scheduler field in the thread attribute object `attr` to the value specified by `inheritsched`.

The `contentionscope` must be either `PTHREAD_INHERIT_SCHED` to indicate that the thread is to inherit the scheduling policy and parameters from the creating thread, or `PTHREAD_EXPLICIT_SCHED` to indicate that the scheduling policy and parameters for this thread are to be set from the corresponding values in the attributes object. If `contentionscope` is `PTHREAD_INHERIT_SCHED`, then the scheduling attributes in the `attr` structure will be ignored at thread creation time.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.12 pthread_attr_getinheritsched - Get Inherit Scheduler Flag

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getinheritsched(
    const pthread_attr_t *attr,
    int *inheritsched
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The inheritsched pointer argument is invalid.

DESCRIPTION:

The `pthread_attr_getinheritsched` routine is used to object the current value of the inherit scheduler field in the thread attribute object `attr`.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.13 pthread_attr_setschedpolicy - Set Scheduling Policy

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setschedpolicy(
    pthread_attr_t *attr,
    int             policy
);
```

STATUS CODES:

EINVAL	The attribute pointer argument is invalid.
EINVAL	The attribute set is not initialized.
ENOTSUP	The specified scheduler policy argument is invalid.

DESCRIPTION:

The `pthread_attr_setschedpolicy` routine is used to set the scheduler policy field in the thread attribute object `attr` to the value specified by `policy`.

Scheduling policies may be one of the following:

- `SCHED_DEFAULT`
- `SCHED_FIFO`
- `SCHED_RR`
- `SCHED_SPORADIC`
- `SCHED_OTHER`

The precise meaning of each of these is discussed elsewhere in this manual.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.14 pthread_attr_getschedpolicy - Get Scheduling Policy

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getschedpolicy(
    const pthread_attr_t *attr,
    int *policy
);
```

STATUS CODES:

- EINVAL** The attribute pointer argument is invalid.
- EINVAL** The attribute set is not initialized.
- EINVAL** The specified scheduler policy argument pointer is invalid.

DESCRIPTION:

The `pthread_attr_getschedpolicy` routine is used to obtain the scheduler policy field from the thread attribute object `attr`. The value of this field is returned in `policy`.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.15 pthread_attr_setschedparam - Set Scheduling Parameters

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_setschedparam(
    pthread_attr_t      *attr,
    const struct sched_param param
);
```

STATUS CODES:

EINVAL The attribute pointer argument is invalid.

EINVAL The attribute set is not initialized.

EINVAL The specified scheduler parameter argument is invalid.

DESCRIPTION:

The `pthread_attr_setschedparam` routine is used to set the scheduler parameters field in the thread attribute object `attr` to the value specified by `param`.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.16 pthread_attr_getschedparam - Get Scheduling Parameters

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_attr_getschedparam(
    const pthread_attr_t *attr,
    struct sched_param  *param
);
```

STATUS CODES:

- EINVAL** The attribute pointer argument is invalid.
- EINVAL** The attribute set is not initialized.
- EINVAL** The specified scheduler parameter argument pointer is invalid.

DESCRIPTION:

The `pthread_attr_getschedparam` routine is used to obtain the scheduler parameters field from the thread attribute object `attr`. The value of this field is returned in `param`.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.17 pthread_create - Create a Thread

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_create(
    pthread_t          *thread,
    const pthread_attr_t *attr,
    void               (*start_routine)( void *),
    void               *arg
);
```

STATUS CODES:

EINVAL	The attribute set is not initialized.
EINVAL	The user specified a stack address and the size of the area was not large enough to meet this processor's minimum stack requirements.
EINVAL	The specified scheduler inheritance policy was invalid.
ENOTSUP	The specified contention scope was PTHREAD_SCOPE_PROCESS.
EINVAL	The specified thread priority was invalid.
EINVAL	The specified scheduling policy was invalid.
EINVAL	The scheduling policy was SCHED_SPORADIC and the specified replenishment period is less than the initial budget.
EINVAL	The scheduling policy was SCHED_SPORADIC and the specified low priority is invalid.
EAGAIN	The system lacked the necessary resources to create another thread, or the self imposed limit on the total number of threads in a process PTHREAD_THREAD_MAX would be exceeded.
EINVAL	Invalid argument passed.

DESCRIPTION:

The `pthread_create` routine is used to create a new thread with the attributes specified by `attr`. If the `attr` argument is `NULL`, then the default attribute set will be used. Modification of the contents of `attr` after this thread is created does not have an impact on this thread.

The thread begins execution at the address specified by `start_routine` with `arg` as its only argument. If `start_routine` returns, then it is functionally equivalent to the thread executing the `pthread_exit` service.

Upon successful completion, the ID of the created thread is returned in the `thread` argument.

NOTES:

There is no concept of a single main thread in RTEMS as there is in a tradition UNIX system. POSIX requires that the implicit return of the main thread results in the same effects as if there were a call to `exit`. This does not occur in RTEMS.

The signal mask of the newly created thread is inherited from its creator and the set of pending signals for this thread is empty.

17.4.18 pthread_exit - Terminate the Current Thread

CALLING SEQUENCE:

```
#include <pthread.h>

void pthread_exit(
    void *status
);
```

STATUS CODES:

NONE

DESCRIPTION:

The `pthread_exit` routine is used to terminate the calling thread. The `status` is made available to any successful join with the terminating thread.

When a thread returns from its start routine, it results in an implicit call to the `pthread_exit` routine with the return value of the function serving as the argument to `pthread_exit`.

NOTES:

Any cancellation cleanup handlers that have been pushed and not yet popped shall be popped in reverse of the order that they were pushed. After all cancellation cleanup handlers have been executed, if the thread has any thread-specific data, destructors for that data will be invoked.

Thread termination does not release or free any application visible resources including but not limited to mutexes, file descriptors, allocated memory, etc.. Similarly, exiting a thread does not result in any process-oriented cleanup activity.

There is no concept of a single main thread in RTEMS as there is in a traditional UNIX system. POSIX requires that the implicit return of the main thread results in the same effects as if there were a call to `exit`. This does not occur in RTEMS.

All access to any automatic variables allocated by the threads is lost when the thread exits. Thus references (i.e. pointers) to local variables of a thread should not be used in a global manner without care. As a specific example, a pointer to a local variable should NOT be used as the return value.

17.4.19 pthread_detach - Detach a Thread

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_detach(
    pthread_t thread
);
```

STATUS CODES:

ESRCH The thread specified is invalid.

EINVAL The thread specified is not a joinable thread.

DESCRIPTION:

The `pthread_detach` routine is used to indicate that storage for `thread` can be reclaimed when the thread terminates without another thread joining with it.

NOTES:

If any threads have previously joined with the specified thread, then they will remain joined with that thread. Any subsequent calls to `pthread_join` on the specified thread will fail.

17.4.20 pthread_join - Wait for Thread Termination

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_join(
    pthread_t    thread,
    void         **value_ptr
);
```

STATUS CODES:

ESRCH	The thread specified is invalid.
EINVAL	The thread specified is not a joinable thread.
EDEADLK	A deadlock was detected or thread is the calling thread.

DESCRIPTION:

The `pthread_join` routine suspends execution of the calling thread until `thread` terminates. If `thread` has already terminated, then this routine returns immediately. The value returned by `thread` (i.e. passed to `pthread_exit` is returned in `value_ptr`.

When this routine returns, then `thread` has been terminated.

NOTES:

The results of multiple simultaneous joins on the same thread is undefined.

If any threads have previously joined with the specified thread, then they will remain joined with that thread. Any subsequent calls to `pthread_join` on the specified thread will fail.

If `value_ptr` is `NULL`, then no value is returned.

17.4.21 pthread_self - Get Thread ID

CALLING SEQUENCE:

```
#include <pthread.h>
```

```
pthread_t pthread_self( void );
```

STATUS CODES:

The value returned is the ID of the calling thread.

DESCRIPTION:

This routine returns the ID of the calling thread.

NOTES:

NONE

17.4.22 pthread_equal - Compare Thread IDs

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_equal(
    pthread_t t1,
    pthread_t t2
);
```

STATUS CODES:

zero The thread ids are not equal.
non-zero The thread ids are equal.

DESCRIPTION:

The `pthread_equal` routine is used to compare two thread IDs and determine if they are equal.

NOTES:

The behavior is undefined if the thread IDs are not valid.

17.4.23 pthread_once - Dynamic Package Initialization

CALLING SEQUENCE:

```
#include <pthread.h>

pthread_once_t once_control = PTHREAD_ONCE_INIT;

int pthread_once(
    pthread_once_t  *once_control,
    void            (*init_routine)(void)
);
```

STATUS CODES:

NONE

DESCRIPTION:

The `pthread_once` routine is used to provide controlled initialization of variables. The first call to `pthread_once` by any thread with the same `once_control` will result in the `init_routine` being invoked with no arguments. Subsequent calls to `pthread_once` with the same `once_control` will have no effect.

The `init_routine` is guaranteed to have run to completion when this routine returns to the caller.

NOTES:

The behavior of `pthread_once` is undefined if `once_control` is automatic storage (i.e. on a task stack) or is not initialized using `PTHREAD_ONCE_INIT`.

17.4.24 `pthread_setschedparam` - Set Thread Scheduling Parameters

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_setschedparam(
    pthread_t      thread,
    int            policy,
    struct sched_param *param
);
```

STATUS CODES:

EINVAL	The scheduling parameters indicated by the parameter <code>param</code> is invalid.
EINVAL	The value specified by <code>policy</code> is invalid.
EINVAL	The scheduling policy was <code>SCHED_SPORADIC</code> and the specified replenishment period is less than the initial budget.
EINVAL	The scheduling policy was <code>SCHED_SPORADIC</code> and the specified low priority is invalid.
ESRCH	The thread indicated was invalid.

DESCRIPTION:

The `pthread_setschedparam` routine is used to set the scheduler parameters currently associated with the thread specified by `thread` to the policy specified by `policy`. The contents of `param` are interpreted based upon the `policy` argument.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

17.4.25 pthread_getschedparam - Get Thread Scheduling Parameters

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_getschedparam(
    pthread_t      thread,
    int            *policy,
    struct sched_param *param
);
```

STATUS CODES:

- EINVAL** The policy pointer argument is invalid.
- EINVAL** The scheduling parameters pointer argument is invalid.
- ESRCH** The thread indicated by the parameter thread is invalid.

DESCRIPTION:

The `pthread_getschedparam` routine is used to obtain the scheduler policy and parameters associated with `thread`. The current policy and associated parameters values returned in `policy` and `param`, respectively.

NOTES:

As required by POSIX, RTEMS defines the feature symbol `_POSIX_THREAD_PRIORITY_SCHEDULING` to indicate that the family of routines to which this routine belongs is supported.

18 Key Manager

18.1 Introduction

The key manager ...

The directives provided by the key manager are:

- `pthread_key_create` - Create Thread Specific Data Key
- `pthread_key_delete` - Delete Thread Specific Data Key
- `pthread_setspecific` - Set Thread Specific Key Value
- `pthread_getspecific` - Get Thread Specific Key Value

18.2 Background

There is currently no text in this section.

18.3 Operations

There is currently no text in this section.

18.4 Directives

This section details the key manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

18.4.1 pthread_key_create - Create Thread Specific Data Key

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_key_create(
    pthread_key_t *key,
    void (*destructor)( void )
);
```

STATUS CODES:

EAGAIN	There were not enough resources available to create another key.
ENOMEM	Insufficient memory exists to create the key.

18.4.2 pthread_key_delete - Delete Thread Specific Data Key

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_key_delete(
pthread_key_t key,
);
```

STATUS CODES:

EINVAL The key was invalid

DESCRIPTION:

NOTES:

18.4.3 pthread_setspecific - Set Thread Specific Key Value

CALLING SEQUENCE:

```
#include <pthread.h>

int pthread_setspecific(
pthread_key_t key,
const void *value
);
```

STATUS CODES:

EINVAL The specified key is invalid.

DESCRIPTION:

NOTES:

18.4.4 pthread_getspecific - Get Thread Specific Key Value

CALLING SEQUENCE:

```
#include <pthread.h>

void *pthread_getspecific(
pthread_key_t key
);
```

STATUS CODES:

NULL There is no thread-specific data associated with the specified key.
non-NULL The data associated with the specified key.

DESCRIPTION:

NOTES:

19 Thread Cancellation Manager

19.1 Introduction

The thread cancellation manager is ...

The directives provided by the thread cancellation manager are:

- `pthread_cancel` - Cancel Execution of a Thread
- `pthread_setcancelstate` - Set Cancelability State
- `pthread_setcanceltype` - Set Cancelability Type
- `pthread_testcancel` - Create Cancellation Point
- `pthread_cleanup_push` - Establish Cancellation Handler
- `pthread_cleanup_pop` - Remove Cancellation Handler

19.2 Background

There is currently no text in this section.

19.3 Operations

There is currently no text in this section.

19.4 Directives

This section details the thread cancellation manager's directives. A subsection is dedicated to each of this manager's directives and describes the calling sequence, related constants, usage, and status codes.

19.4.1 pthread_cancel - Cancel Execution of a Thread

CALLING SEQUENCE:

```
int pthread_cancel(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

19.4.2 pthread_setcancelstate - Set Cancelability State

CALLING SEQUENCE:

```
int pthread_setcancelstate(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

19.4.3 pthread_setcanceltype - Set Cancelability Type

CALLING SEQUENCE:

```
int pthread_setcanceltype(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

19.4.4 pthread_testcancel - Create Cancellation Point

CALLING SEQUENCE:

```
int pthread_testcancel(  
    );
```

STATUS CODES:

E The

DESCRIPTION:**NOTES:**

19.4.5 pthread_cleanup_push - Establish Cancellation Handler

CALLING SEQUENCE:

```
int pthread_cleanup_push(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

19.4.6 pthread_cleanup_pop - Remove Cancellation Handler

CALLING SEQUENCE:

```
int pthread_cleanup_pop(  
    );
```

STATUS CODES:

E The

DESCRIPTION:

NOTES:

20 Services Provided by C Library (libc)

20.1 Introduction

This section lists the routines that provided by the Newlib C Library.

20.2 Standard Utility Functions (stdlib.h)

- `abort` - Abnormal termination of a program
- `abs` - Integer absolute value (magnitude)
- `assert` - Macro for Debugging Diagnostics
- `atexit` - Request execution of functions at program exit
- `atof` - String to double or float
- `atoi` - String to integer
- `bsearch` - Binary search
- `calloc` - Allocate space for arrays
- `div` - Divide two integers
- `ecvtbuf` - Double or float to string of digits
- `ecvt` - Double or float to string of digits (malloc result)
- `__env_lock` - Lock environment list for `getenv` and `setenv`
- `gvcvt` - Format double or float as string
- `exit` - End program execution
- `getenv` - Look up environment variable
- `labs` - Long integer absolute value (magnitude)
- `ldiv` - Divide two long integers
- `malloc` - Allocate memory
- `realloc` - Reallocate memory
- `free` - Free previously allocated memory
- `mallinfo` - Get information about allocated memory
- `__malloc_lock` - Lock memory pool for `malloc` and `free`
- `mbstowcs` - Minimal multibyte string to wide string converter
- `mblen` - Minimal multibyte length
- `mbtowc` - Minimal multibyte to wide character converter
- `qsort` - Sort an array
- `rand` - Pseudo-random numbers
- `strtod` - String to double or float
- `strtol` - String to long
- `strtoul` - String to unsigned long
- `system` - Execute command string
- `wcstombs` - Minimal wide string to multibyte string converter
- `wctomb` - Minimal wide character to multibyte converter

20.3 Character Type Macros and Functions (ctype.h)

- `isalnum` - Alphanumeric character predicate
- `isalpha` - Alphabetic character predicate
- `isascii` - ASCII character predicate
- `iscntrl` - Control character predicate
- `isdigit` - Decimal digit predicate
- `islower` - Lower-case character predicate
- `isprint` - Printable character predicates (`isprint`, `isgraph`)
- `ispunct` - Punctuation character predicate
- `isspace` - Whitespace character predicate
- `isupper` - Uppercase character predicate
- `isxdigit` - Hexadecimal digit predicate
- `toascii` - Force integers to ASCII range
- `tolower` - Translate characters to lower case
- `toupper` - Translate characters to upper case

20.4 Input and Output (stdio.h)

- `clearerr` - Clear file or stream error indicator
- `fclose` - Close a file
- `feof` - Test for end of file
- `ferror` - Test whether read/write error has occurred
- `fflush` - Flush buffered file output
- `fgetc` - Get a character from a file or stream
- `fgetpos` - Record position in a stream or file
- `fgets` - Get character string from a file or stream
- `fiprintf` - Write formatted output to file (integer only)
- `fopen` - Open a file
- `fdopen` - Turn an open file into a stream
- `fputc` - Write a character on a stream or file
- `fputs` - Write a character string in a file or stream
- `fread` - Read array elements from a file
- `freopen` - Open a file using an existing file descriptor
- `fseek` - Set file position
- `fsetpos` - Restore position of a stream or file
- `ftell` - Return position in a stream or file
- `fwrite` - Write array elements from memory to a file or stream
- `getc` - Get a character from a file or stream (macro)
- `getchar` - Get a character from standard input (macro)

- `gets` - Get character string from standard input (obsolete)
- `iprintf` - Write formatted output (integer only)
- `mktemp` - Generate unused file name
- `perror` - Print an error message on standard error
- `putc` - Write a character on a stream or file (macro)
- `putchar` - Write a character on standard output (macro)
- `puts` - Write a character string on standard output
- `remove` - Delete a file's name
- `rename` - Rename a file
- `rewind` - Reinitialize a file or stream
- `setbuf` - Specify full buffering for a file or stream
- `setvbuf` - Specify buffering for a file or stream
- `siprintf` - Write formatted output (integer only)
- `printf` - Write formatted output
- `scanf` - Scan and format input
- `tmpfile` - Create a temporary file
- `tmpnam` - Generate name for a temporary file
- `vprintf` - Format variable argument list

20.5 Strings and Memory (`string.h`)

- `bcmp` - Compare two memory areas
- `bcopy` - Copy memory regions
- `bzero` - Initialize memory to zero
- `index` - Search for character in string
- `memchr` - Find character in memory
- `memcmp` - Compare two memory areas
- `memcpy` - Copy memory regions
- `memmove` - Move possibly overlapping memory
- `memset` - Set an area of memory
- `rindex` - Reverse search for character in string
- `strcasecmp` - Compare strings ignoring case
- `strcat` - Concatenate strings
- `strchr` - Search for character in string
- `strcmp` - Character string compare
- `strcoll` - Locale specific character string compare
- `strcpy` - Copy string
- `strcspn` - Count chars not in string
- `strerror` - Convert error number to string
- `strlen` - Character string length

- `strlwr` - Convert string to lower case
- `strncasecmp` - Compare strings ignoring case
- `strncat` - Concatenate strings
- `strncmp` - Character string compare
- `strncpy` - Counted copy string
- `strpbrk` - Find chars in string
- `strrchr` - Reverse search for character in string
- `strspn` - Find initial match
- `strstr` - Find string segment
- `strtok` - Get next token from a string
- `strupr` - Convert string to upper case
- `strxfrm` - Transform string

20.6 Signal Handling (`signal.h`)

- `raise` - Send a signal
- `signal` - Specify handler subroutine for a signal

20.7 Time Functions (`time.h`)

- `asctime` - Format time as string
- `clock` - Cumulative processor time
- `ctime` - Convert time to local and format as string
- `difftime` - Subtract two times
- `gmtime` - Convert time to UTC (GMT) traditional representation
- `localtime` - Convert time to local representation
- `mktime` - Convert time to arithmetic representation
- `strftime` - Flexible calendar time formatter
- `time` - Get current calendar time (as single number)

20.8 Locale (`locale.h`)

- `setlocale` - Select or query locale

20.9 Reentrant Versions of Functions

- Equivalent for `errno` variable:
 - `errno_r` - XXX
- Locale functions:
 - `localeconv_r` - XXX
 - `setlocale_r` - XXX
- Equivalents for `stdio` variables:
 - `stdin_r` - XXX

- `stdout_r` - XXX
- `stderr_r` - XXX
- Stdio functions:
 - `fdopen_r` - XXX
 - `perror_r` - XXX
 - `tempnam_r` - XXX
 - `fopen_r` - XXX
 - `putchar_r` - XXX
 - `tmpnam_r` - XXX
 - `getchar_r` - XXX
 - `puts_r` - XXX
 - `tmpfile_r` - XXX
 - `gets_r` - XXX
 - `remove_r` - XXX
 - `vfprintf_r` - XXX
 - `iprntf_r` - XXX
 - `rename_r` - XXX
 - `vsnprintf_r` - XXX
 - `mkstemp_r` - XXX
 - `snprintf_r` - XXX
 - `vsprintf_r` - XXX
 - `mktemp_t` - XXX
 - `sprintf_r` - XXX
- Signal functions:
 - `init_signal_r` - XXX
 - `signal_r` - XXX
 - `kill_r` - XXX
 - `_sigtramp_r` - XXX
 - `raise_r` - XXX
- Stdlib functions:
 - `calloc_r` - XXX
 - `mblen_r` - XXX
 - `srand_r` - XXX
 - `dtoa_r` - XXX
 - `mbstowcs_r` - XXX
 - `strtod_r` - XXX
 - `free_r` - XXX
 - `mbtowc_r` - XXX
 - `strtol_r` - XXX

- `getenv_r` - XXX
- `memalign_r` - XXX
- `strtoul_r` - XXX
- `mallinfo_r` - XXX
- `mstats_r` - XXX
- `system_r` - XXX
- `malloc_r` - XXX
- `rand_r` - XXX
- `wcstombs_r` - XXX
- `malloc_r` - XXX
- `realloc_r` - XXX
- `wctomb_r` - XXX
- `malloc_stats_r` - XXX
- `setenv_r` - XXX
- String functions:
 - `strtok_r` - XXX
- System functions:
 - `close_r` - XXX
 - `link_r` - XXX
 - `unlink_r` - XXX
 - `execve_r` - XXX
 - `lseek_r` - XXX
 - `wait_r` - XXX
 - `fcntl_r` - XXX
 - `open_r` - XXX
 - `write_r` - XXX
 - `fork_r` - XXX
 - `read_r` - XXX
 - `fstat_r` - XXX
 - `sbrk_r` - XXX
 - `gettimeofday_r` - XXX
 - `stat_r` - XXX
 - `getpid_r` - XXX
 - `times_r` - XXX
- Time function:
 - `asctime_r` - XXX

20.10 Miscellaneous Macros and Functions

- `unctrl` - Return printable representation of a character

20.11 Variable Argument Lists

- Stdarg (stdarg.h):
 - va_start - XXX
 - va_arg - XXX
 - va_end - XXX
- Vararg (varargs.h):
 - va_alist - XXX
 - va_start-trad - XXX
 - va_arg-trad - XXX
 - va_end-trad - XXX

20.12 Reentrant System Calls

- open_r - XXX
- close_r - XXX
- lseek_r - XXX
- read_r - XXX
- write_r - XXX
- fork_r - XXX
- wait_r - XXX
- stat_r - XXX
- fstat_r - XXX
- link_r - XXX
- unlink_r - XXX
- sbrk_r - XXX

21 Services Provided by the Math Library (libm)

21.1 Introduction

This section lists the routines that provided by the Newlib Math Library (libm).

21.2 Standard Math Functions (math.h)

- `acos` - Arccosine
- `acosh` - Inverse hyperbolic cosine
- `asin` - Arcsine
- `asinh` - Inverse hyperbolic sine
- `atan` - Arctangent
- `atan2` - Arctangent of y/x
- `atanh` - Inverse hyperbolic tangent
- `jN` - Bessel functions (j_N and y_N)
- `cbrt` - Cube root
- `copysign` - Sign of Y and magnitude of X
- `cosh` - Hyperbolic cosine
- `erf` - Error function (erf and erfc)
- `exp` - Exponential
- `expm1` - Exponential of x and - 1
- `fabs` - Absolute value (magnitude)
- `floor` - Floor and ceiling (floor and ceil)
- `fmod` - Floating-point remainder (modulo)
- `frexp` - Split floating-point number
- `gamma` - Logarithmic gamma function
- `hypot` - Distance from origin
- `ilogb` - Get exponent
- `infinity` - Floating infinity
- `isnan` - Check type of number
- `ldexp` - Load exponent
- `log` - Natural logarithms
- `log10` - Base 10 logarithms
- `log1p` - Log of $1 + X$
- `matherr` - Modifiable math error handler
- `modf` - Split fractional and integer parts
- `nan` - Floating Not a Number
- `nextafter` - Get next representable number
- `pow` - X to the power Y

- `remainder` - remainder of X divided by Y
- `scalbn` - `scalbn`
- `sin` - Sine or cosine (sin and cos)
- `sinh` - Hyperbolic sine
- `sqrt` - Positive square root
- `tan` - Tangent
- `tanh` - Hyperbolic tangent

22 Status of Implementation

This chapter provides an overview of the status of the implementation of the POSIX API for RTEMS. The *POSIX 1003.1b Compliance Guide* provides more detailed information regarding the implementation of each of the numerous functions, constants, and macros specified by the POSIX 1003.1b standard.

RTEMS supports many of the process and user/group oriented services in a "single user/single process" manner. This means that although these services may be of limited usefulness or functionality, they are provided and do work in a coherent manner. This is significant when porting existing code from UNIX to RTEMS.

- Implementation
 - The current implementation of `dup()` is insufficient.
 - FIFOs `mkfifo()` are not currently implemented.
 - Asynchronous IO is not implemented.
 - The `flockfile()` family is not implemented
 - `getc/putc` unlocked family is not implemented
 - Shared Memory is not implemented
 - Mapped Memory is not implemented
 - NOTES:
 - For Shared Memory and Mapped Memory services, it is unclear what level of support is appropriate and possible for RTEMS.
- Functional Testing
 - Tests for unimplemented services
- Performance Testing
 - There are no POSIX Performance Tests.
- Documentation
 - Many of the service description pages are not complete in this manual. These need to be completed and information added to the background and operations sections.
 - Example programs (not just tests) would be very nice.

Command and Variable Index

-		fdopen.....	146
_exit.....	14	fileno.....	145
A		flockfile.....	147
access.....	89	fork.....	4
aio_cancel.....	126	fpathconf.....	99
aio_error.....	124	fstat.....	87
aio_fsync.....	128	fsync.....	114
aio_read.....	121	ftruncate.....	95
aio_return.....	125	ftrylockfile.....	148
aio_suspend.....	127	funlockfile.....	149
aio_write.....	122	G	
alarm.....	34, 35	getc_unlocked.....	150
asctime_r.....	160	getchar_unlocked.....	151
C		getcwd.....	73
cfgetispeed.....	130	getdents.....	92
cfgetospeed.....	131	getegid.....	44
cfsetispeed.....	132	getenv.....	55
cfsetospeed.....	133	geteuid.....	42
chdir.....	71	getgid.....	43
chmod.....	90	getgrgid.....	166
chown.....	93	getgrgid_r.....	167
clock_getres.....	238	getgrnam.....	168
clock_gettime.....	235	getgrnam_r.....	169
clock_settime.....	237	getgroups.....	47
close.....	107	getlogin.....	48
closedir.....	70	getlogin_r.....	49
creat.....	76	getpgrp.....	50
ctermid.....	57	getpid.....	39
ctime_r.....	161	getppid.....	40
		getpwnam.....	172
		getpwnam_r.....	173
		getpwuid.....	170
		getpwuid_r.....	171
		gettimeofday.....	242
		getuid.....	41
		gmtime_r.....	162
D		I	
dup.....	105	isatty.....	60
dup2.....	106	K	
E		kill.....	26
execl.....	5	L	
execle.....	7	link.....	78
execlp.....	9	lio_listio.....	123
execv.....	6	localtime_r.....	163
execve.....	8	longjmp.....	155
execvp.....	10	lseek.....	113
F			
fchdir.....	72		
fchmod.....	91		
fcntl.....	111		
fdatasync.....	115		

lstat 88

M

microseconds alarm 35
 mkdir 81
 mkfifo 82
 mknod 101
 mlock 220
 mlockall 218
 mmap 222
 mount 117
 mprotect 224
 mq_attr 252
 mq_close 258
 mq_getattr 264
 mq_notify 262
 mq_open 256
 mq_receive 261
 mq_send 260
 mq_setattr 263
 mq_unlink 259
 mqd_t 251
 msync 225
 munlock 221
 munlockall 219
 munmap 223

N

nanosleep 241

O

open 74
 opendir 65

P

pathconf 97
 pause 29
 pipe 104
 pthread_atfork 11
 pthread_attr_destroy 268
 pthread_attr_getdetachstate 270
 pthread_attr_getinheritsched 278
 pthread_attr_getschedparam 282
 pthread_attr_getschedpolicy 280
 pthread_attr_getscope 276
 pthread_attr_getstackaddr 274
 pthread_attr_getstacksize 272
 pthread_attr_init 267
 pthread_attr_setdetachstate 269
 pthread_attr_setinheritsched 277
 pthread_attr_setschedparam 281
 pthread_attr_setschedpolicy 279
 pthread_attr_setscope 275
 pthread_attr_setstackaddr 273

pthread_attr_setstacksize 271
 pthread_cancel 300
 pthread_cleanup_pop 305
 pthread_cleanup_push 304
 pthread_cond_broadcast 213
 pthread_cond_destroy 211
 pthread_cond_init 210
 pthread_cond_signal 212
 pthread_cond_timedwait 215
 pthread_cond_wait 214
 pthread_condattr_destroy 207
 pthread_condattr_getpshared 209
 pthread_condattr_init 206
 pthread_condattr_setpshared 208
 pthread_create 283
 pthread_detach 286
 pthread_equal 289
 pthread_exit 285
 pthread_getschedparam 292
 pthread_join 287
 pthread_kill 23
 pthread_mutex_destroy 198
 pthread_mutex_getprioceiling 204
 pthread_mutex_init 197
 pthread_mutex_lock 199
 pthread_mutex_setprioceiling 203
 pthread_mutex_timedlock 201
 pthread_mutex_trylock 200
 pthread_mutex_unlock 202
 pthread_mutexattr_destroy 190
 pthread_mutexattr_getprioceiling 194
 pthread_mutexattr_getprotocol 192
 pthread_mutexattr_getpshared 196
 pthread_mutexattr_init 189
 pthread_mutexattr_setprioceiling 193
 pthread_mutexattr_setprotocol 191
 pthread_mutexattr_setpshared 195
 pthread_once 290
 pthread_self 288
 pthread_setcancelstate 301
 pthread_setcanceltype 302
 pthread_setschedparam 291
 pthread_sigmask 25
 pthread_testcancel 303
 putc_unlocked 152
 putchar_unlocked 153

R

rand_r 164
 read 108
 readdir 66
 readlink 80
 readv 119
 rename 85
 rewinddir 67
 rmdir 84

S

scandir 68
 sched_get_priority_max 232
 sched_get_priority_min 231
 sched_rr_get_interval 233
 sched_yield 234
 sem_close 180
 sem_destroy 178
 sem_getvalue 186
 sem_init 177
 sem_open 179
 sem_post 185
 sem_t 175
 sem_timedwait 184
 sem_trywait 183
 sem_unlink 181
 sem_wait 182
 setenv 56
 setgid 46
 setjmp 154
 setlocale 144
 setpgid 52
 setsid 51
 setuid 45
 shm_open 226
 shm_unlink 227
 sigaction 22
 sigaddset 17
 sigdelset 18
 sigemptyset 21
 sigfillset 19
 sigismember 20
 siglongjmp 157
 sigpending 27
 sigprocmask 24
 sigqueue 33
 sigsetjmp 156
 sigsuspend 28
 sigtimedwait 32
 sigwait 30
 sigwaitinfo 31

sleep 239
 stat 86
 strtok_r 159
 symlink 79
 sync 116
 sysconf 61

T

tcdrain 137
 tcflow 139
 tcflush 138
 tcgetattr 134
 tcgetpgrp 140
 tcsendbreak 136
 tcsetattr 135
 tcsetpgrp 141
 telldir 69
 time 243
 times 54
 truncate 96
 ttyname 58
 ttyname_r 59
 tzset 158

U

umask 77
 uname 53
 unlink 83
 unmount 118
 usecs alarm 35
 usleep 240
 utime 94

W

wait 12
 waitpid 13
 write 110
 writev 120

Concept Index

A

acquire ownership of file stream 147
 add a signal to a signal set 17
 associate stream with file descriptor 146
 asynchronous file synchronization 128
 asynchronous read 121
 asynchronous write 122

B

broadcast a condition variable 213

C

cancel asynchronous i/o request 126
 cancel execution of a thread 300
 change access and/or modification times of an
 inode 94
 change memory protection 224
 changes file mode 90
 changes permissions of a file 91
 changes the current working directory 71, 72
 changes the owner and/or group of a file 93
 check permissions for a file 89
 close a message queue 258
 close a named semaphore 180
 closes a file 107
 compare thread ids 289
 create a directory 101
 create a new file or rewrite an existing one 76
 create a process 4
 create a thread 283
 create an inter 104
 create cancellation point 303
 create session and set process group id 51
 creates a link to a file 78
 creates a symbolic link to a file 79

D

delay process execution 239, 240
 delay with high resolution 241
 delete a directory 84
 delete a signal from a signal set 18
 destroy a condition variable 211
 destroy a condition variable attribute set 207
 destroy a mutex 198
 destroy a mutex attribute set 190
 destroy a thread attribute set 268
 destroy an unnamed semaphore 178
 detach a thread 286
 determine if file descriptor is terminal 60
 determine terminal device name 58
 discards terminal data 138

duplicates an open file descriptor 105, 106
 dynamic package initialization 290
 dynamically set the priority ceiling 203

E

empty a signal set 21
 ends directory read operation 70
 establish cancellation handler 304
 examine and change process blocked signals 24
 examine and change signal action 22
 examine and change thread blocked signals 25
 examine pending signals 27
 execute a file 5, 6, 7, 8, 9, 10

F

fill a signal set 19

G

generate terminal pathname 57
 get character from stdin without locking 151
 get character without locking 150
 get clock resolution 238
 get configurable system variables 61
 get detach state 270
 get directory entries 92
 get effective group id 44
 get effective user id 42
 get environment variables 55
 get group file entry for id 166
 get group file entry for name 168
 get inherit scheduler flag 278
 get maximum priority value 232
 get message queue attributes 264
 get minimum priority value 231
 get parent process id 40
 get password file entry for uid 170
 get process group id 50
 get process id 39
 get process shared attribute 209
 get process times 54
 get real group id 43
 get scheduling parameters 282
 get scheduling policy 280
 get supplementary group ids 47
 get system name 53
 get the blocking protocol 192
 get the current priority ceiling 204
 get the priority ceiling 194
 get the time of day 242
 get the value of a semaphore 186
 get the visibility 196

get thread id 288
 get thread scheduling parameters 292
 get thread scheduling scope 276
 get thread stack address 274
 get thread stack size 272
 get time in seconds 243
 get timeslicing quantum 233
 get user id 41
 get user name 48
 get user name, reentrant 49
 gets configuration values for files 97, 99
 gets current working directory 73
 gets file status 87, 88
 gets foreground process group id 140
 gets information about a file 86
 gets terminal attributes 134

I

initialize a condition variable 210
 initialize a condition variable attribute set 206
 initialize a mutex 197
 initialize a mutex attribute set 189
 initialize a thread attribute set 267
 initialize an unnamed semaphore 177
 initialize time conversion information 158
 is signal a member of a signal set 20

L

list directed i/o 123
 lock a mutex 199
 lock a mutex with timeout 201
 lock a range of the process address space 220
 lock the address space of a process 218

M

makes a directory 81
 makes a fifo special file 82
 manipulates an open file descriptor 111
 map process addresses to a memory object 222
 memory object synchronization 225
 microsecond delay process execution 240
 mount a file system 117

N

non 155, 157, 183
 notify process that a message is available 262

O

obtain file descriptor number for this file 145
 obtain the name of a symbolic link destination .. 80
 obtain time of day 235
 open a directory 65

open a message queue 256
 open a named semaphore 179
 open a shared memory object 226
 opens a file 74

P

password file entry for name 172
 poll to acquire ownership of file stream 148
 poll to lock a mutex 200
 put character to stdin without locking 153
 put character without locking 152

Q

queue a signal to a process 33

R

reads a directory 66
 reads from a file 108
 reads terminal input baud rate 130
 reads terminal output baud rate 131
 receive a message from a message queue 261
 reentrant determine terminal device name 59
 reentrant extract token from string 159
 reentrant get group file entry 167
 reentrant get group file entry for name 169
 reentrant get password file entry for name 173
 reentrant get password file entry for uid 171
 reentrant get user name 49
 reentrant local time conversion 163
 reentrant random number generation 164
 reentrant struct tm to ascii time conversion ... 160
 reentrant time_t to ascii time conversion 161
 reentrant utc time conversion 162
 register fork handlers 11
 release ownership of file stream 149
 remove a message queue 259
 remove a shared memory object 227
 remove cancellation handler 305
 removes a directory entry 83
 renames a file 85
 reposition read/write file offset 113
 resets the readdir() pointer 67
 retrieve error status of asynchronous i/o operation
 124
 retrieve return status asynchronous i/o operation
 125
 return current location in directory stream 69

S

save context for non 154
 save context with signal status for non 156
 scan a directory for matching entries 68
 schedule alarm 34
 schedule alarm in microseconds 35

send a message to a message queue 260
 send a signal to a process 26
 send a signal to a thread 23
 sends a break to a terminal 136
 set cancelability state 301
 set cancelability type 302
 set detach state 269
 set environment variables 56
 set group id 46
 set inherit scheduler flag 277
 set message queue attributes 263
 set process group id for job control 52
 set process shared attribute 208
 set scheduling parameters 281
 set scheduling policy 279
 set terminal attributes 135
 set the blocking protocol 191
 set the current locale 144
 set the priority ceiling 193
 set the visibility 195
 set thread scheduling parameters 291
 set thread scheduling scope 275
 set thread stack address 273
 set thread stack size 271
 set time of day 237
 set user id 45
 sets a file creation mask 77
 sets foreground process group id 141
 sets terminal input baud rate 132
 sets terminal output baud rate 133
 signal a condition variable 212
 suspend process execution 29
 suspends/restarts terminal output 139
 synchronize file complete in 114
 synchronize file in 115
 synchronize file systems 116
 synchronously accept a signal 30, 31
 synchronously accept a signal with timeout 32

T

terminate a process 14
 terminate the current thread 285
 truncate a file to a specified length 95, 96

U

unlink a semaphore 181
 unlock a mutex 202
 unlock a range of the process address space ... 221
 unlock a semaphore 185
 unlock the address space of a process 219
 unmap previously mapped addresses 223
 unmount file systems 118
 uses delay process execution 240

V

vectored read from a file 119
 vectored write to a file 120

W

wait for a signal 28
 wait for asynchronous i/o request 127
 wait for process termination 12, 13
 wait for thread termination 287
 wait on a condition variable 214
 wait on a semaphore 182
 wait on a semaphore for a specified time 184
 wait with timeout a condition variable 215
 waits for all output to be transmitted to the
 terminal 137
 writes to a file 110

Y

yield the processor 234

